CSE P 590 / CSE M 590 (Spring 2010)

# Computer Security and Privacy

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Thanks to Dan Boneh, Dieter Gollmann, John Manferdelli, John Mitchell, Vitaly Shmatikov, Bennet Yee, and many others for sample slides and materials ...

# Goals for Today

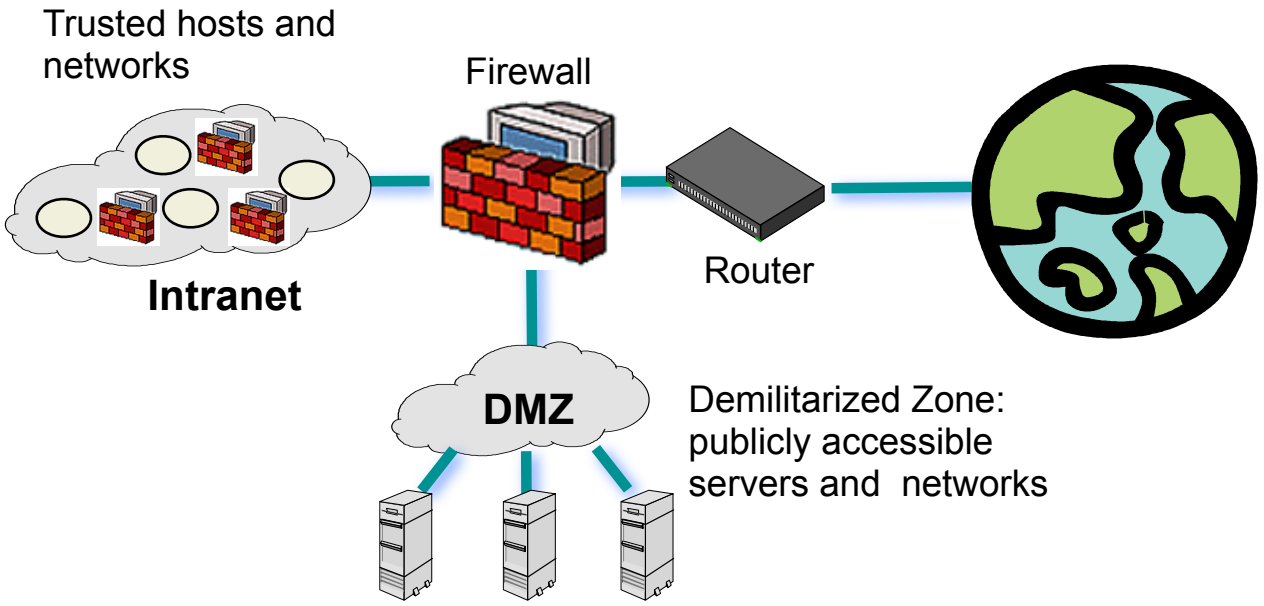
- Lab 2 discussion / HW 3 discussion

- Network security
- Hardware security (based on requests)
- Research reading

# Network Security

# Firewalls

- Idea: separate local network from the Internet



# Castle and Moat Analogy

- More like the moat around a castle than a firewall
  - Restricts access from the outside
  - Restricts outbound connections, too
    - Important: filter out undesirable activity from internal hosts!

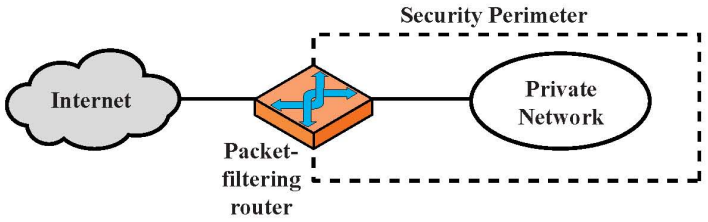
# Firewall Locations in the Network

- Between internal LAN and external network
- At the gateways of sensitive subnetworks within the organizational LAN
  - Payroll's network must be protected separately within the corporate network
- On end-user machines
  - "Personal firewall"
  - Microsoft's Internet Connection Firewall (ICF) comes standard with Windows XP

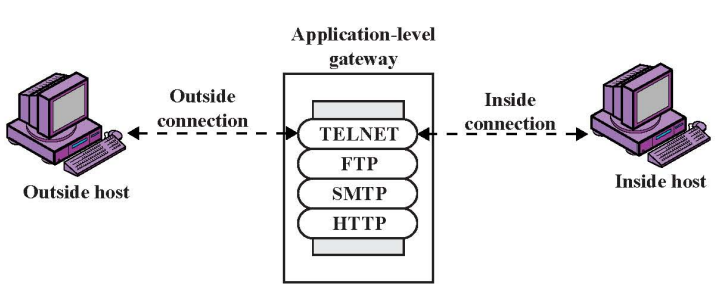
# Firewall Types

- Packet- or session-filtering router (filter)
- Proxy gateway
  - All incoming traffic is directed to firewall, all outgoing traffic appears to come from firewall
  - Application-level: separate proxy for each application
    - Different proxies for SMTP (email), HTTP, FTP, etc.
    - Filtering rules are application-specific
  - Circuit-level: application-independent, "transparent"
    - Only generic IP traffic filtering (example: SOCKS)
- Personal firewall with application-specific rules
  - E.g., no outbound telnet connections from email client

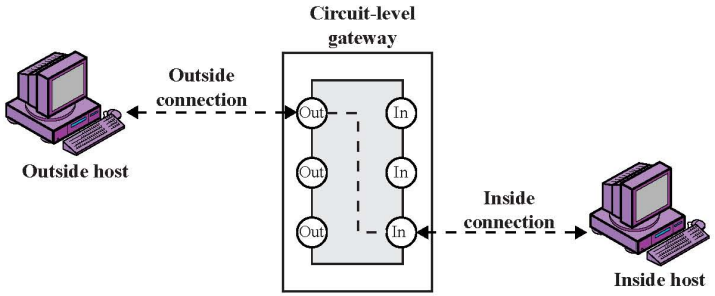
# Firewall Types: Illustration




(a) Packet-filtering router




(b) Application-level gateway




(c) Circuit-level gateway

# Packet Filtering

- For each packet, firewall decides whether to allow it to proceed
  - Decision must be made on per-packet basis
    - Stateless; cannot examine packet's context (TCP connection, application to which it belongs, etc.)
- To decide, use information available in the packet
  - IP source and destination addresses, ports
  - Protocol identifier (TCP, UDP, ICMP, etc.)
  - TCP flags (SYN, ACK, RST, PSH, FIN)
  - ICMP message type
- Filtering rules are based on pattern-matching

# Packet Filtering Examples

**A**

| action | ourhost | port | theirhost | port | comment |
|---|---|---|---|---|---|
| block | * | * | SPIGOT | * | we don't trust these people |
| allow | OUR-GW | 25 | * | * | connection to our SMTP port |

**B**

| action | ourhost | port | theirhost | port | comment |
|---|---|---|---|---|---|
| block | * | * | * | * | default |

**C**

| action | ourhost | port | theirhost | port | comment |
|---|---|---|---|---|---|
| allow | * | * | * | 25 | connection to their SMTP port |

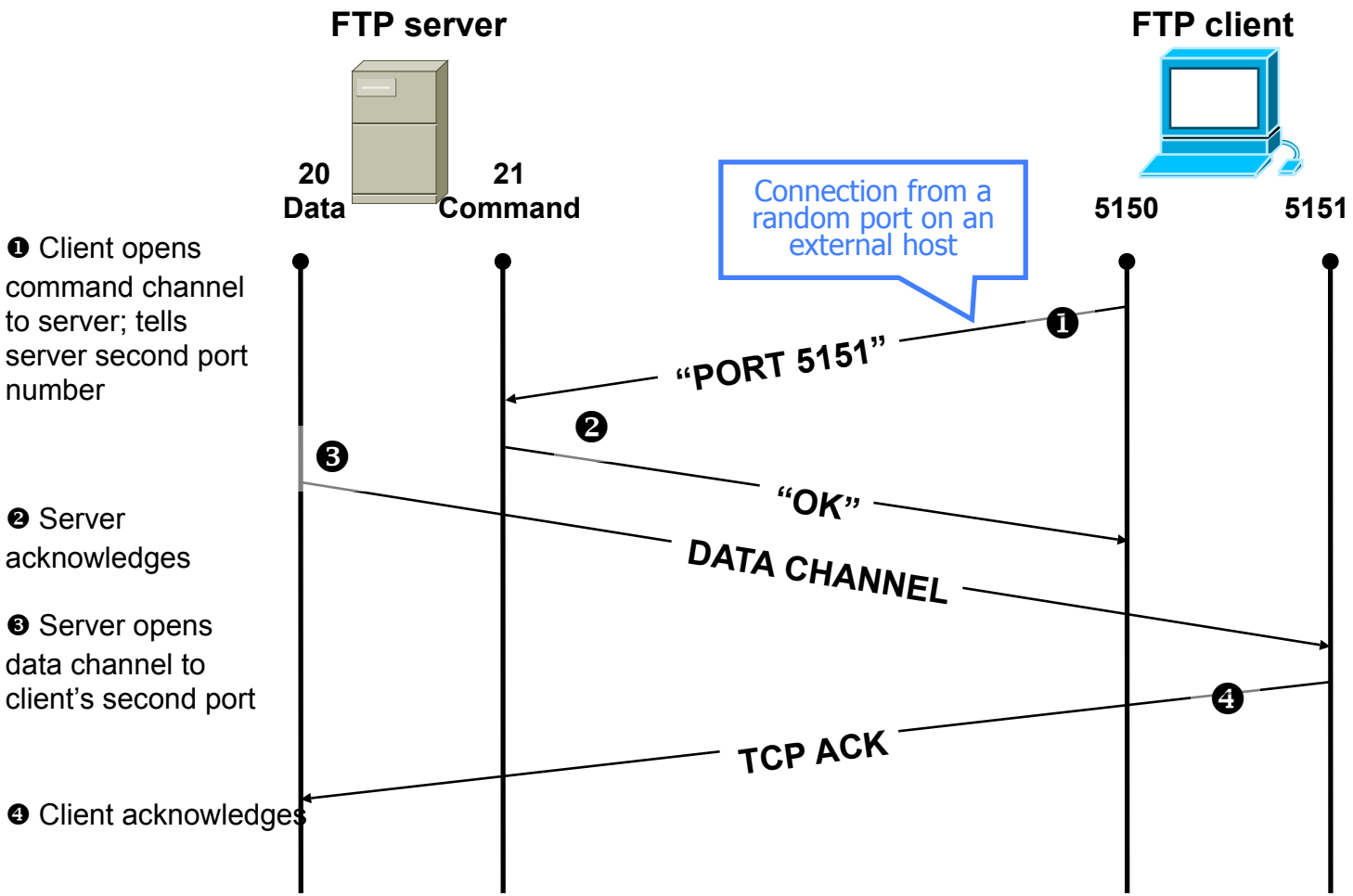
**D**

| action | src | port | dest | port | flags | comment |
|---|---|---|---|---|---|---|
| allow | {our hosts} | * | * | 25 | | our packets to their SMTP port |
| allow | * | 25 | * | * | ACK | their replies |

**E**

| action | src | port | dest | port | flags | comment |
|---|---|---|---|---|---|---|
| allow | {our hosts} | * | * | * | | our outgoing calls |
| allow | * | * | * | * | ACK | replies to our calls |
| allow | * | * | * | >1024 | | traffic to nonservers |

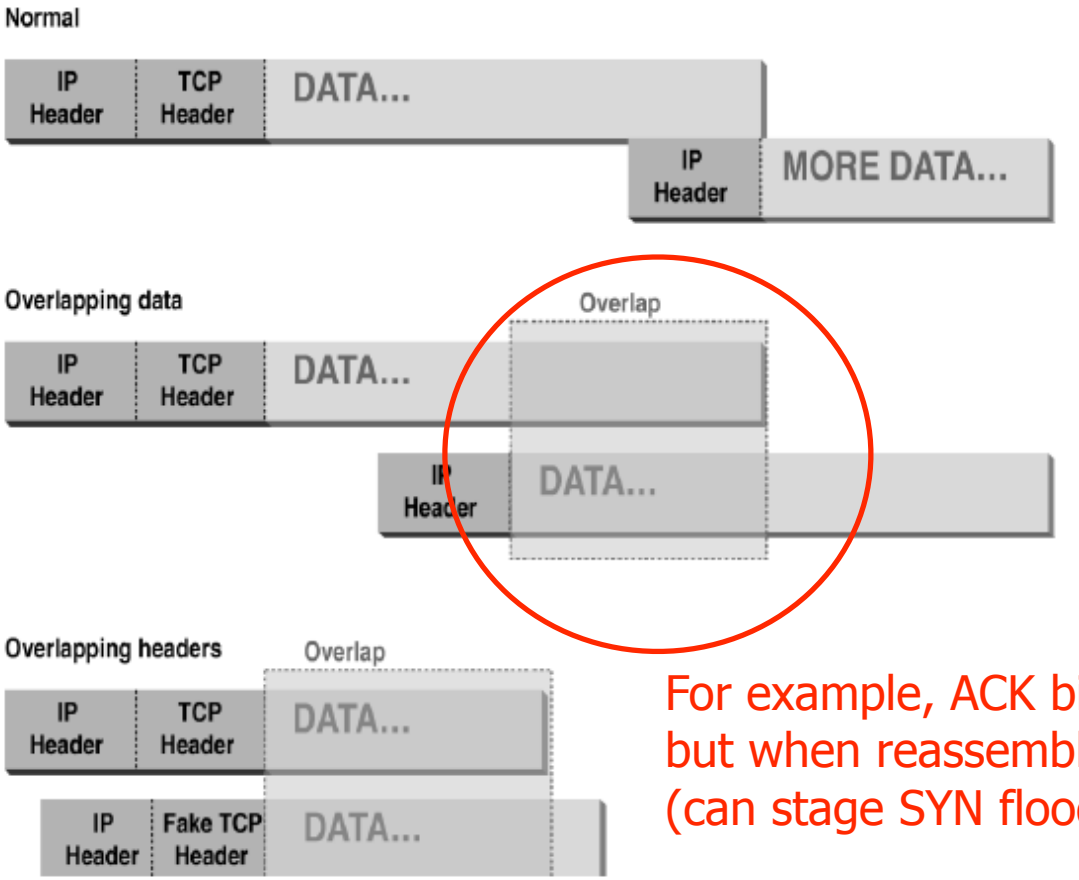
# Example: FTP (borrowed from Wenke Lee)



# Weaknesses of Packet Filters

- Do not prevent application-specific attacks
  - For example, if there is a buffer overflow in URL decoding routine, firewall will not block an attack string
- No user authentication mechanisms
  - … except (spoofable) address-based authentication
  - Firewalls don't have any upper-level functionality
- Vulnerable to TCP/IP attacks such as spoofing
  - Solution: list of addresses for each interface (packets with internal addresses shouldn't come from outside)
- Security breaches due to misconfiguration

# Abnormal Fragmentation

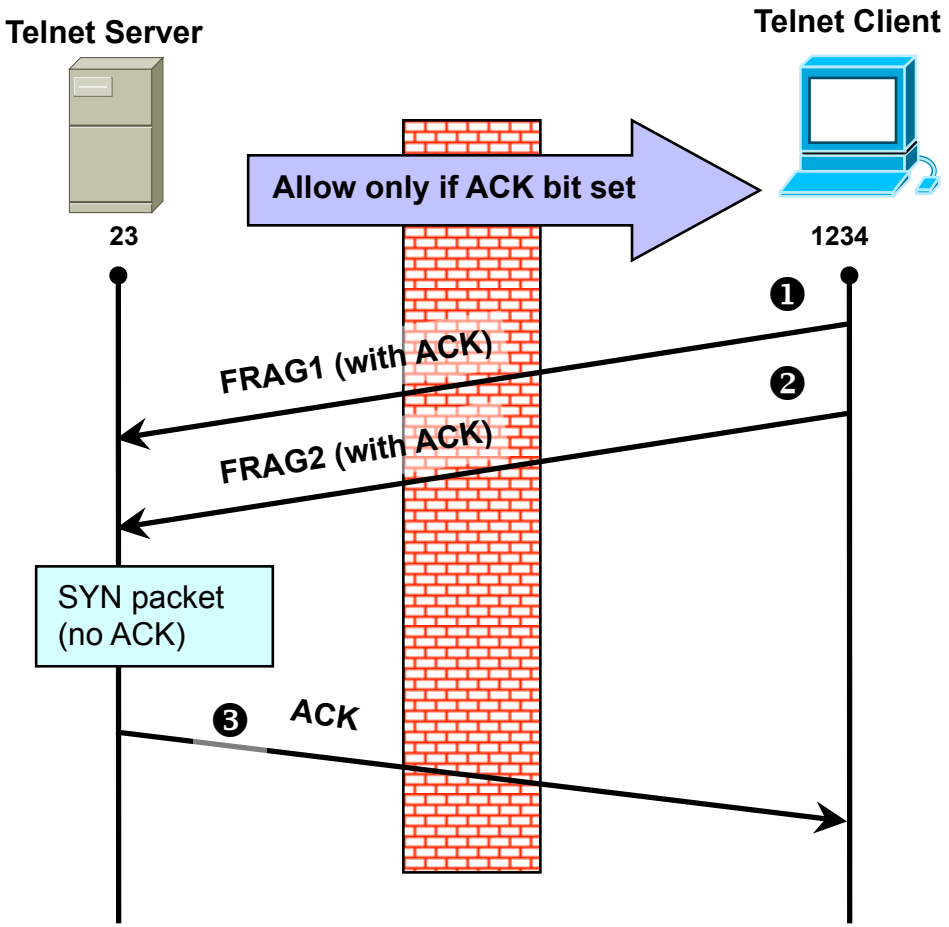



For example, ACK bit is set in both fragments, but when reassembled, SYN bit is set (can stage SYN flooding through firewall)

# Fragmentation Attack (borrowed from Wenke Lee)

❶,❷ Send 2 fragments with the ACK bit set; fragment offsets are chosen so that the full datagram re-assembled by server forms a packet with the SYN bit set (the fragment offset of the second packet overlaps into the space of the first packet)

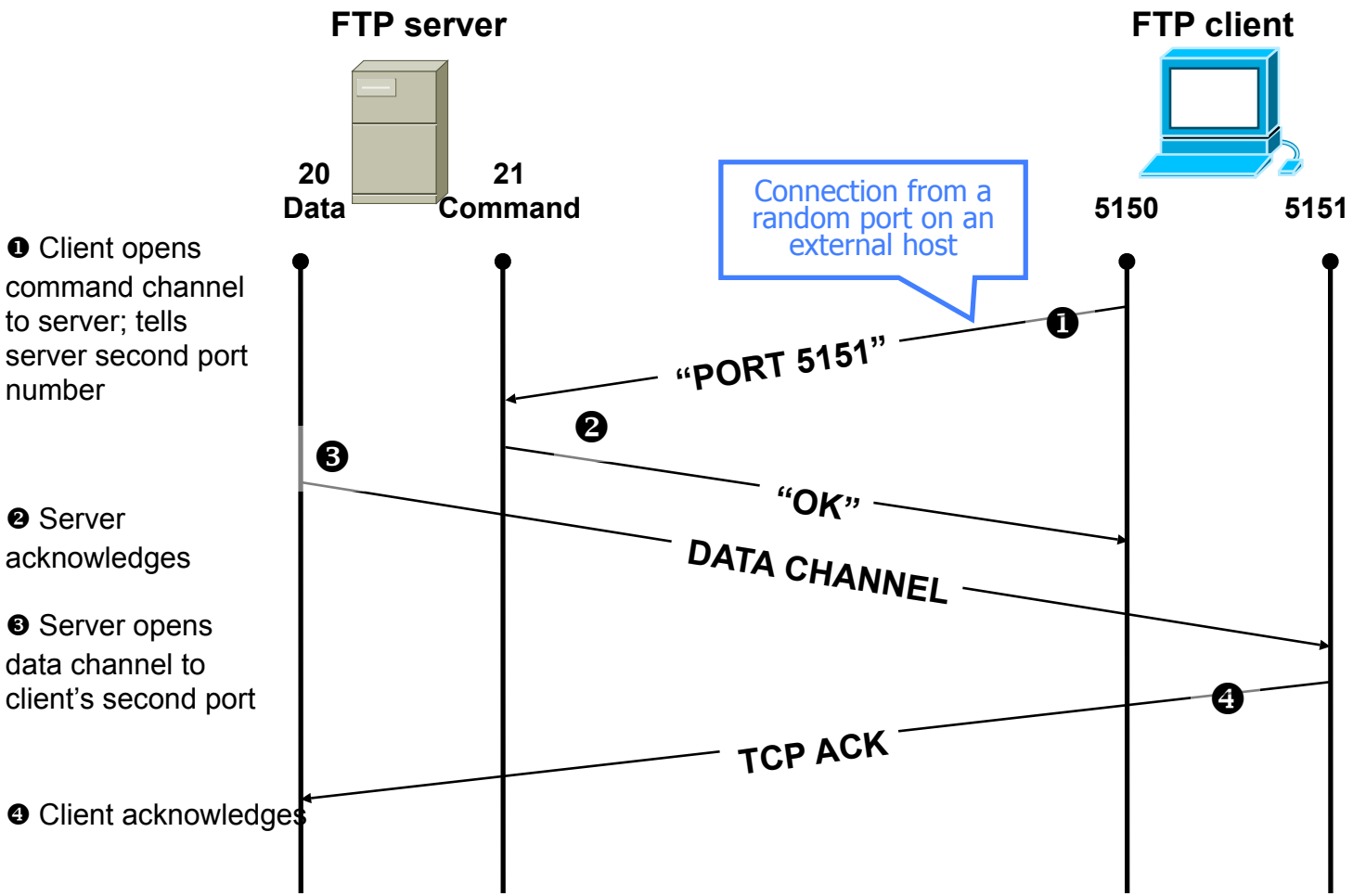
❸ All following packets will have the ACK bit set



# Stateless Filtering Is Not Enough

- In TCP connections, ports with numbers less than 1024 are permanently assigned to servers
  - 20,21 for FTP, 23 for telnet, 25 for SMTP, 80 for HTTP…
- Clients use ports numbered from 1024 to 16383
  - They must be available for clients to receive responses
- What should a firewall do if it sees, say, an incoming request to some client's port 5612?
  - It must allow it: this could be a server's response in a previously established connection…
  - …OR it could be malicious traffic
  - Can't tell without keeping state for each connection

# Example: FTP (borrowed from Wenke Lee)



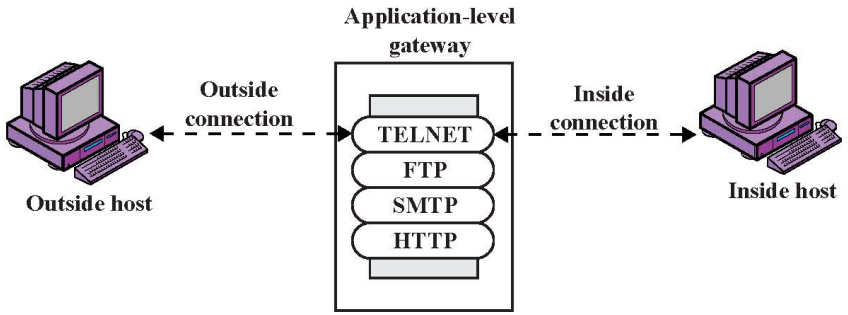
# Session Filtering

- Decision is still made separately for each packet, but in the context of a connection
  - If new connection, then check against security policy
  - If existing connection, then look it up in the table and update the table, if necessary
    - Only allow incoming traffic to a high-numbered port if there is an established connection to that port
- Hard to filter stateless protocols (UDP) and ICMP
- Typical filter: deny everything that's not allowed
  - Must be careful filtering out service traffic such as ICMP
- Filters can be bypassed with IP tunneling

# Example: Connection State Table

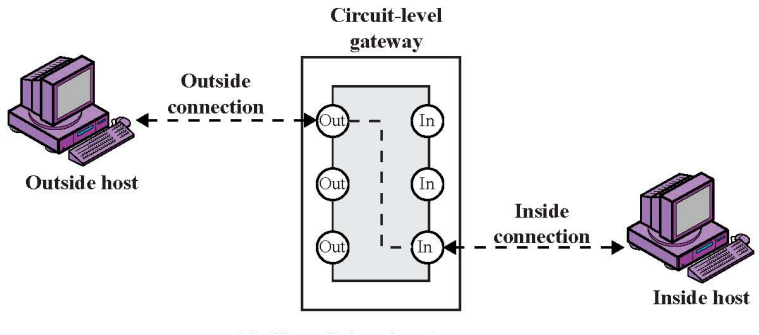
| Source Address | Source Port | Destination Address | Destination Port | Connection State |
|---|---|---|---|---|
| 192.168.1.100 | 1030 | 210.9.88.29 | 80 | Established |
| 192.168.1.102 | 1031 | 216.32.42.123 | 80 | Established |
| 192.168.1.101 | 1033 | 173.66.32.122 | 25 | Established |
| 192.168.1.106 | 1035 | 177.231.32.12 | 79 | Established |
| 223.43.21.231 | 1990 | 192.168.1.6 | 80 | Established |
| 219.22.123.32 | 2112 | 192.168.1.6 | 80 | Established |
| 210.99.212.18 | 3321 | 192.168.1.6 | 80 | Established |
| 24.102.32.23 | 1025 | 192.168.1.6 | 80 | Established |
| 223.212.212 | 1046 | 192.168.1.6 | 80 | Established |

# Application-Level Gateway



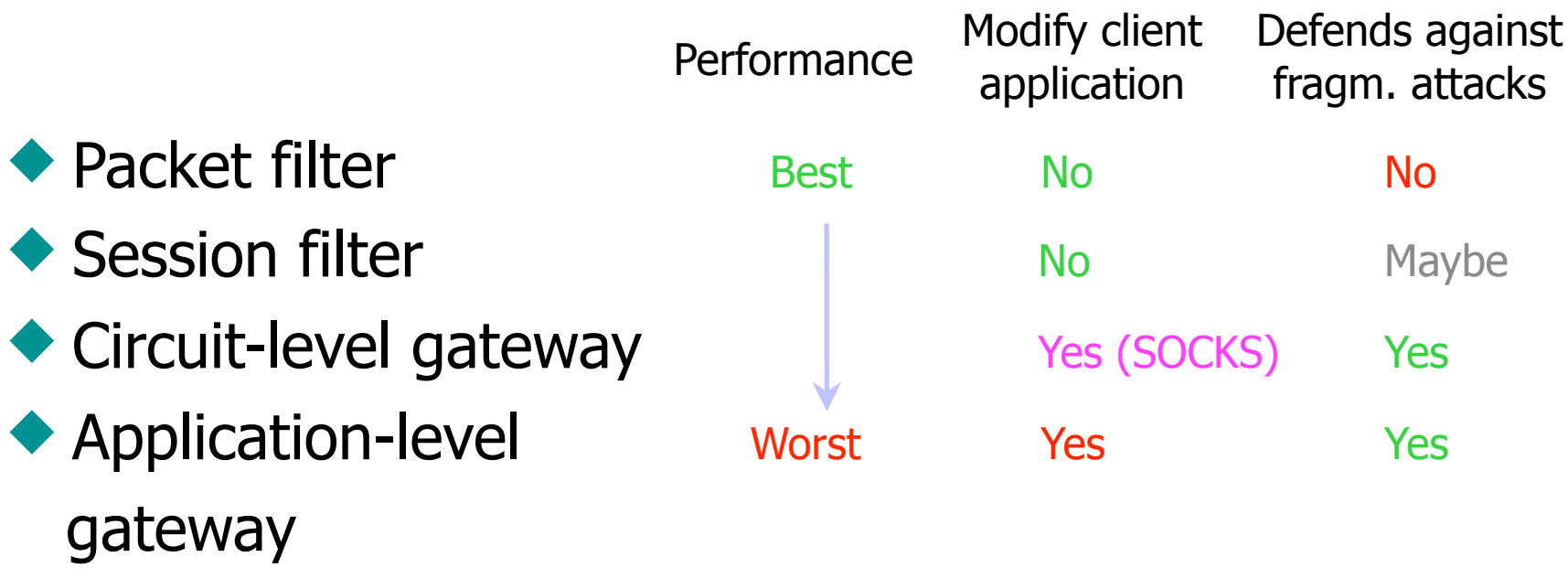

- Splices and relays two application-specific connections
  - Example: Web browser proxy
  - Daemon spawns proxy process when communication is detected
  - Big processing overhead, but can log and audit all activity
- Can support high-level user-to-gateway authentication
  - Log into the proxy server with name, password, etc
- Simpler filtering rules than for arbitrary TCP/IP traffic
- Each application requires implementing its own proxy

# Circuit-Level Gateway




- Splices two TCP connections, relays TCP segments
- Less control over data than application-level gateway
  - Does not examine the contents of TCP segment
- Client's TCP stack must be aware of the gateway
  - Client applications are often adapted to support SOCKS
- Often used when internal users are trusted
  - Application-level proxy on inbound connections, circuit-level proxy on outbound connections (lower overhead)

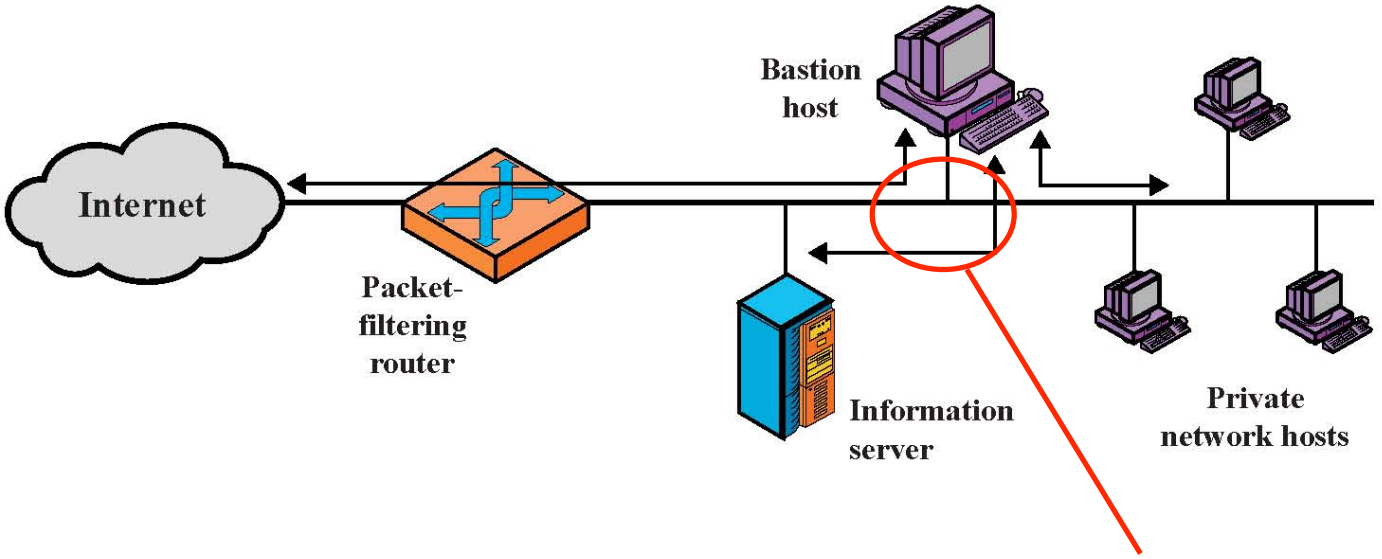
# Comparison


Performance
Modify client application
Defends against fragm. attacks
Packet filter
Session filter
Circuit-level gateway
Application-level gateway
Best
Worst
No
No
Maybe
Yes (SOCKS)
Yes
Yes
Yes

# Bastion Host

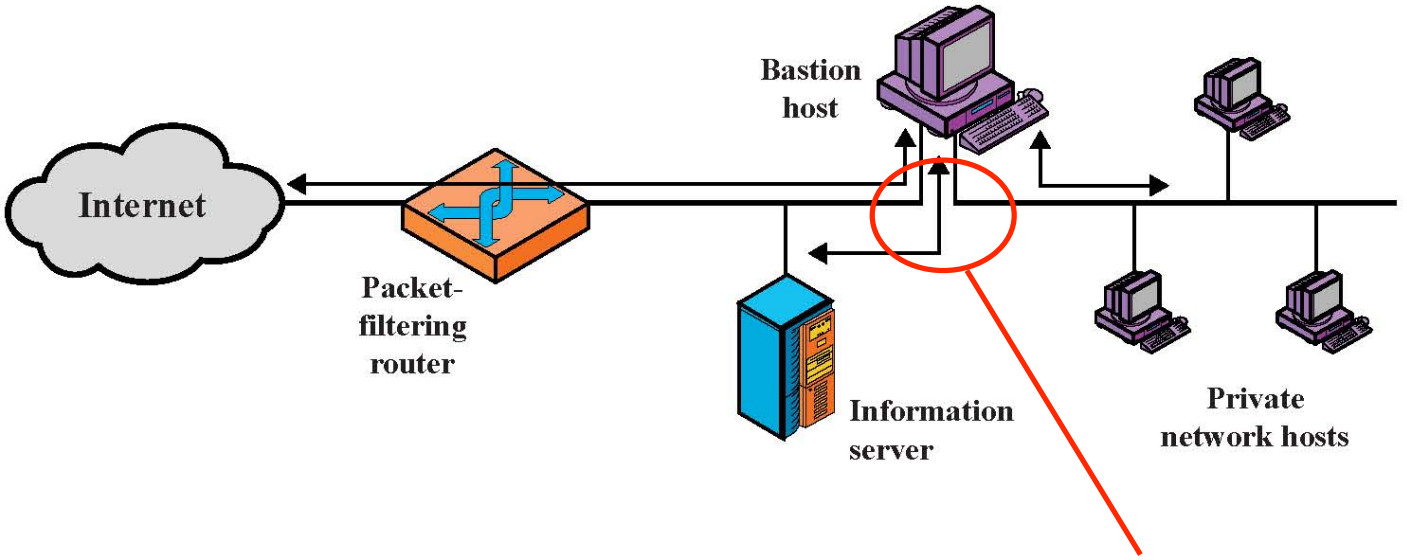
- Bastion host is a hardened system implementing application-level gateway behind packet filter
  - All non-essential services are turned off
  - Application-specific proxies for supported services
    - Each proxy supports only a subset of application's commands, is logged and audited, disk access restricted, runs as a non-privileged user in a separate directory (independent of others)
  - Support for user authentication
- All traffic flows through bastion host
  - Packet router allows external packets to enter only if their destination is bastion host, and internal packets to leave only if their origin is bastion host

# Single-Homed Bastion Host

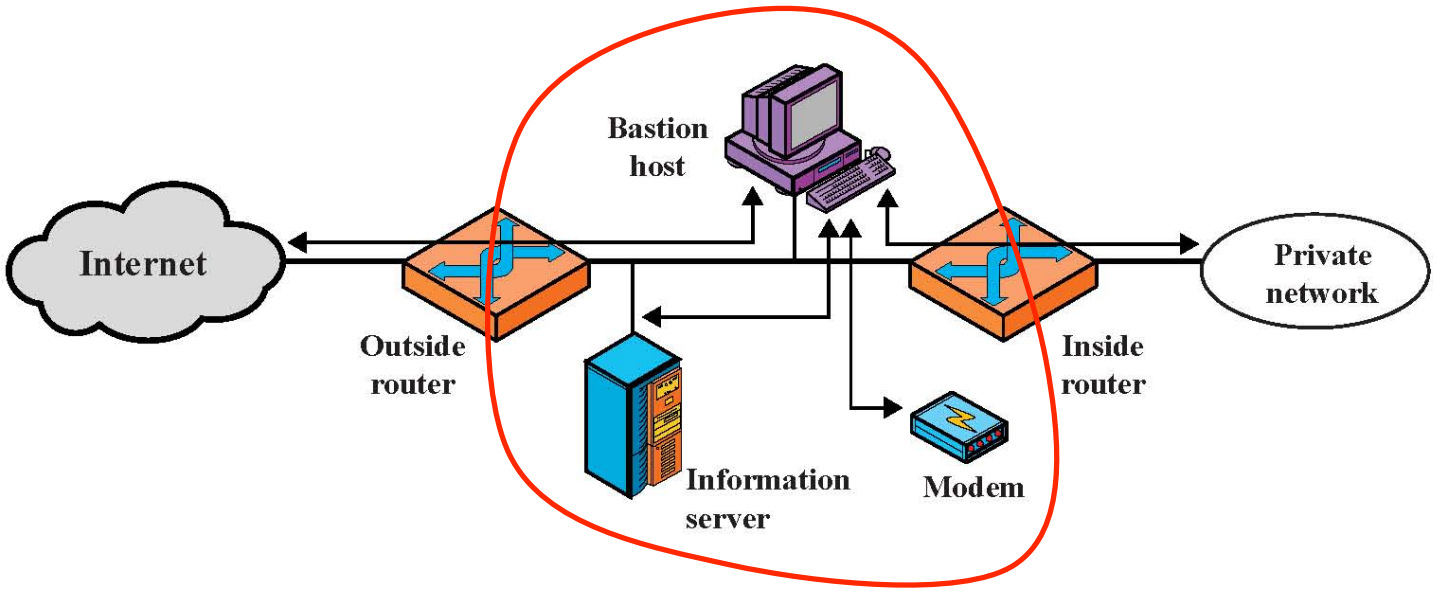



If packet filter is compromised, traffic can flow to internal network

# Dual-Homed Bastion Host



# Screened Subnet




Only the screened subnet is visible to the external network; internal network is invisible

# Protecting Addresses and Routes

- Hide IP addresses of hosts on internal network
  - Only services that are intended to be accessed from outside need to reveal their IP addresses
  - Keep other addresses secret to make spoofing harder
- Use NAT (network address translation) to map addresses in packet headers to internal addresses
  - 1-to-1 or N-to-1 mapping
- Filter route announcements
  - No need to advertise routes to internal hosts
  - Prevent attacker from advertising that the shortest route to an internal host lies through him

# General Problems with Firewalls

- Interfere with networked applications
- Doesn't solve all the problems
  - Buggy software (e.g., buffer overflow exploits)
  - Bad protocol design (e.g., WEP in 802.11b)
- Generally don't prevent denial of service
- Don't prevent insider attacks
- Increasing complexity and potential for misconfiguration

# Cold Boot Attacks

http://citp.princeton.edu/memory/

# Power Analysis

http://www.cc.gatech.edu/~traynor/f08/slides/lecture11-dpa.pdf

# More Hardware Security (Not Covered)

http://www.slideshare.net/guest3bd2a12/advanced-hardware-hacking-techniques-presentation