Anonymous Credentials: How to show credentials without compromising privacy

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Credentials: Motivation

- ID cards
 - Sometimes used for other uses



- E.g. prove you're over 21, or verify your address
- Don't necessarily need to reveal all of your information
- Don't necessarily want issuer of ID to track all of it's uses
- How can we get the functionality/verifiability of an physical id in electronic form without extra privacy loss

Credentials: Motivation

- The goal
 - Users should be able to
 - obtain credentials
 - Show some properties
 - Without
 - Revealing additional information
 - Allowing tracking

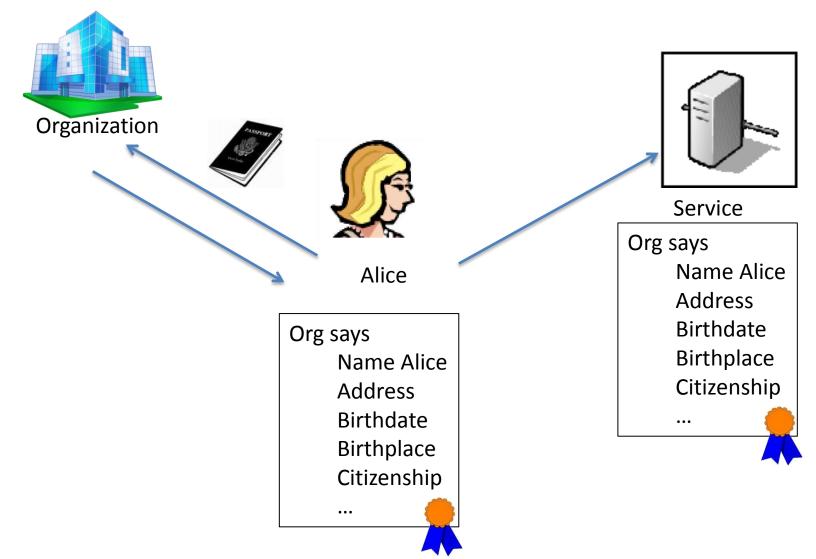
Credentials: Motivation

- Other applications
 - Transit tokens/passes
 - Electronic currency
 - Online polling

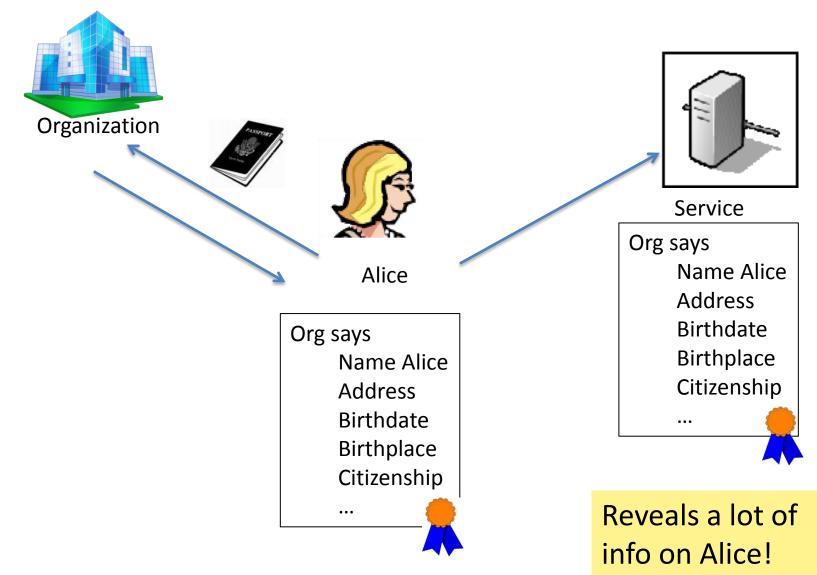
• Implementations

- Idemix (IBM), UProve (Microsoft)

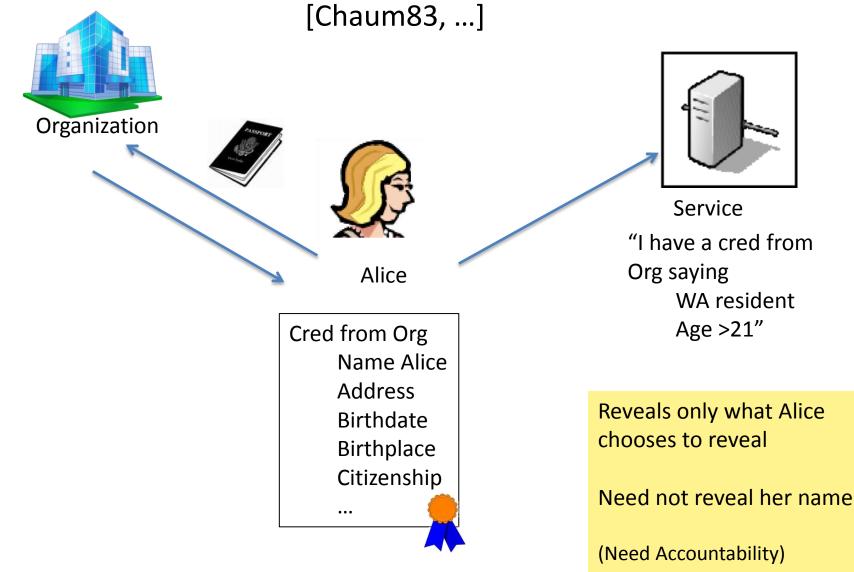
Credentials



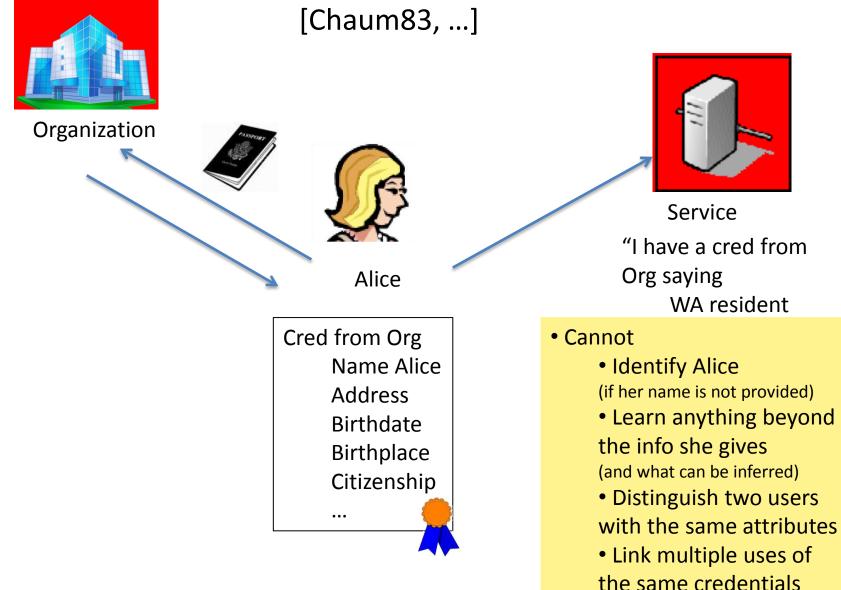
Credentials



A new model



A new model



How can we do this?

Signatures/Certs?
 – No privacy!

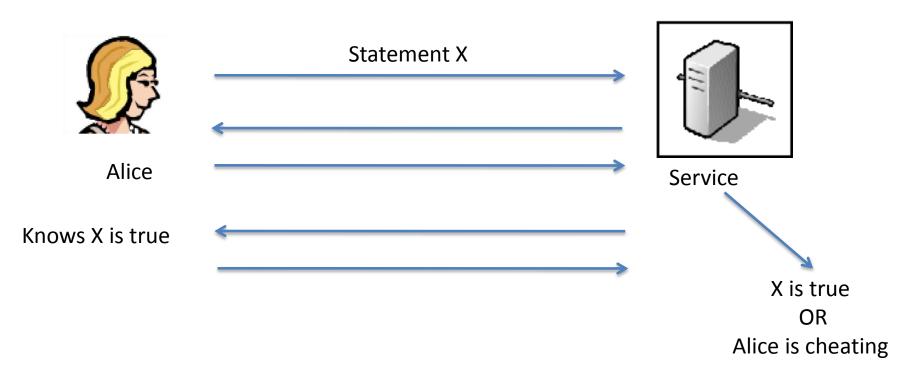


- What about other crypto tools?
- We will use
 - Zero Knowledge Proof of knowledge
 - (interactive or Fiat-Shamir)
 - Commitments
 - Blind signatures

Roadmap

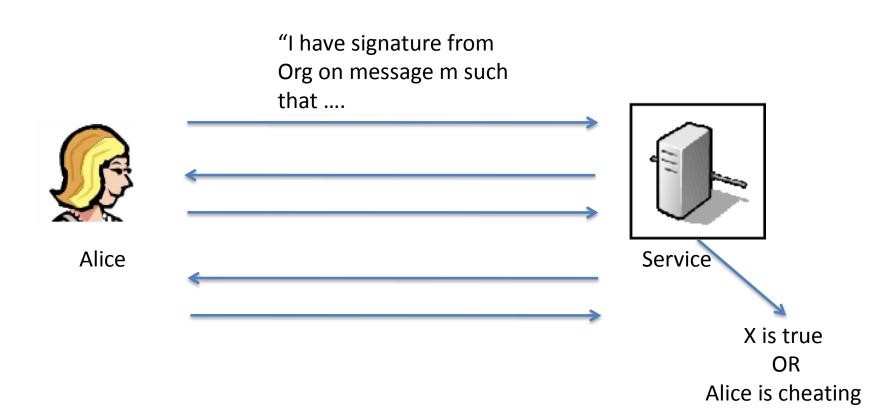
- Review crypto tools
- Construct basic credential systems
- Additional issues
 - Revocation
 - Deciding who to revoke
- Additional features
 - Non-interactive credentials/signatures
 - Delegation
- Conclusion

Zero Knowledge Proofs



Alice wants to convince service that statement X is true, Without revealing any other information

Zero Knowledge Proofs



Alice wants to convince service that she has such a signature Without revealing any other information

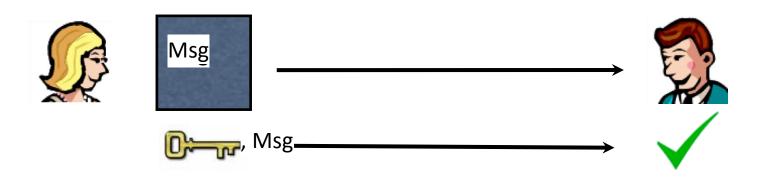
Fiat Shamir: get challenge from hash function

Commitments

Like locked box or safe

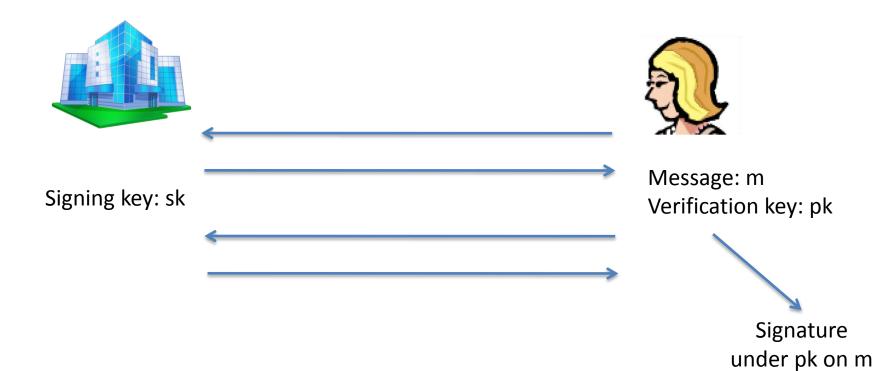


- Hiding hard to tell which message is committed to
- Binding there is a unique message corresponding to each commitment



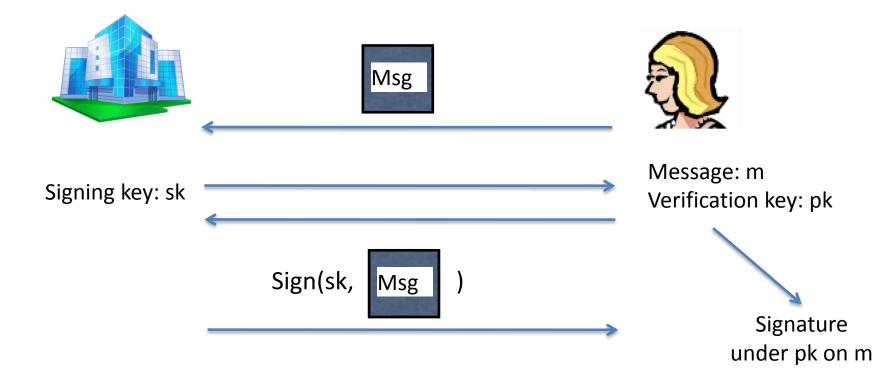
E.g. Pederson Commitment: $C = g^m h^r$

Blind signatures



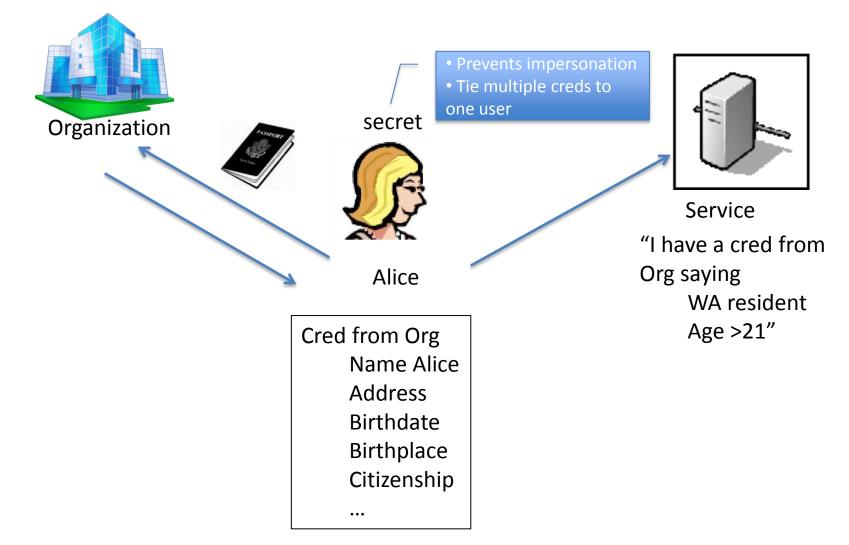
Alice learns only signature on her message. Signer learns nothing.

Blind signatures

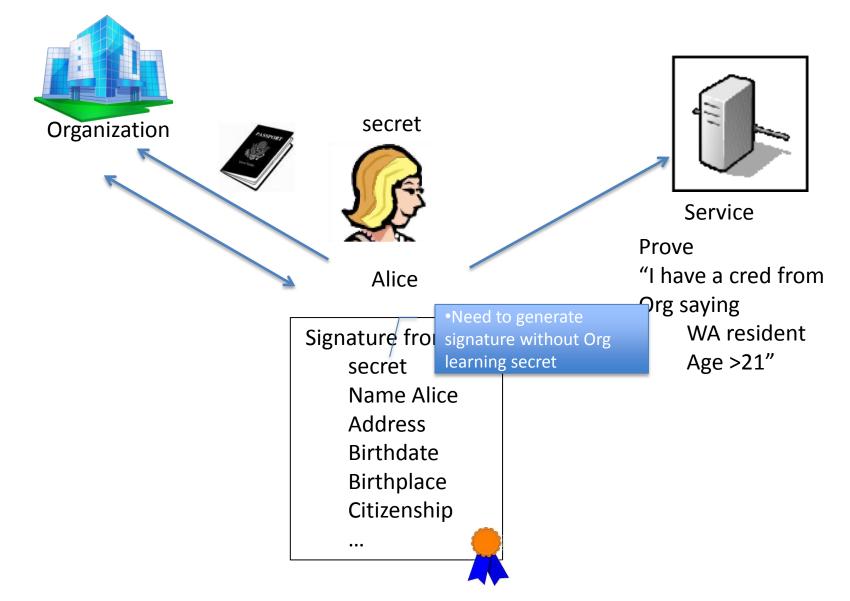


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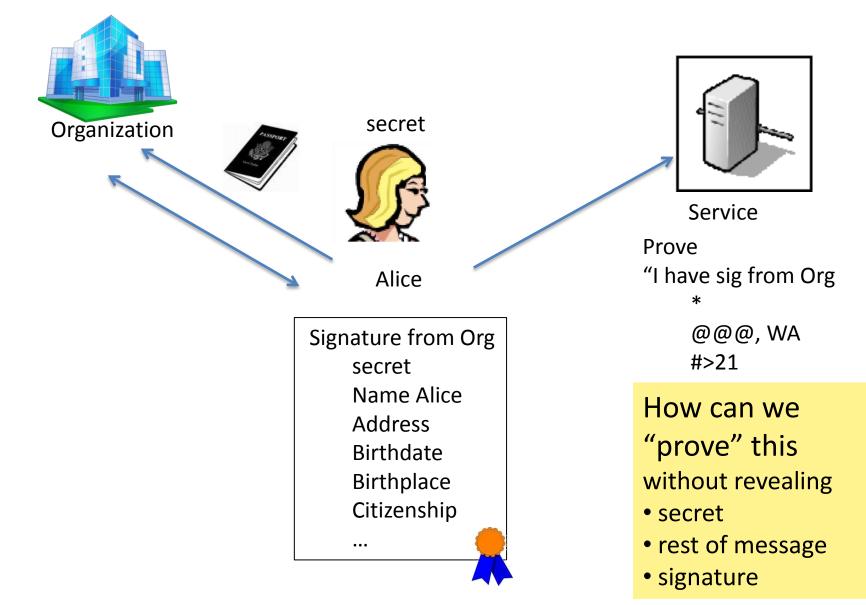
How it works (abstractly) Anonymous Credentials/Minimal Disclosure Tokens



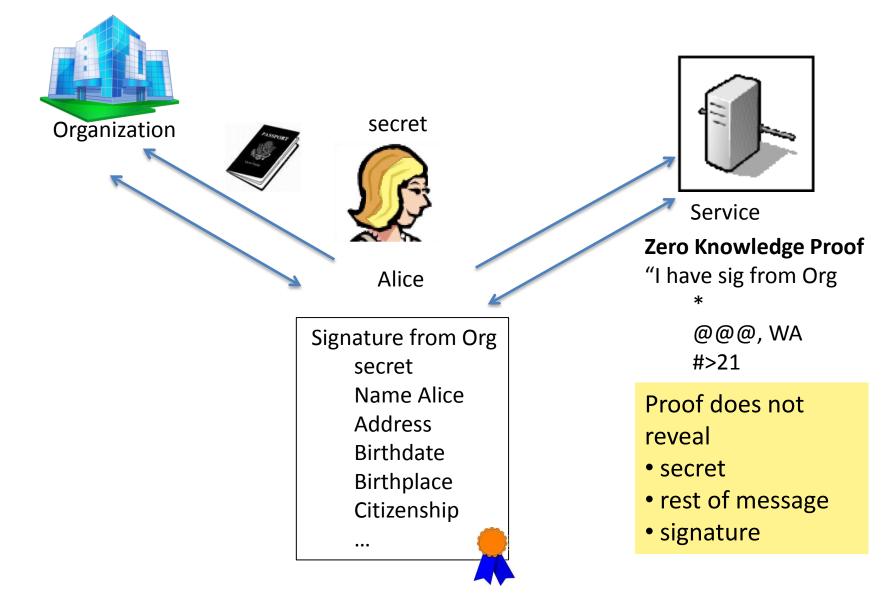
How it works



How it works

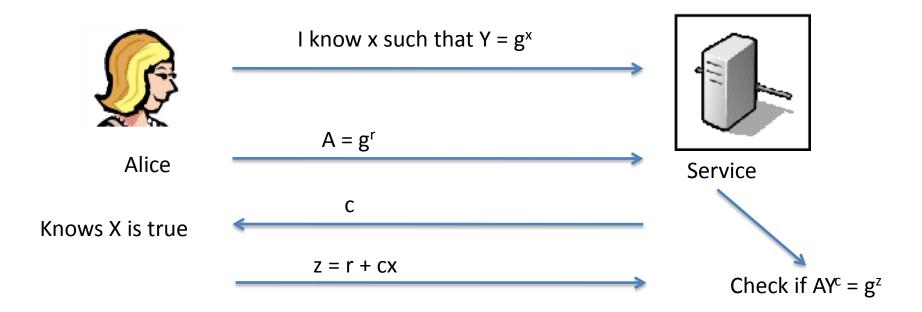


How it works



- Depends on how we implement proofs and blind signatures
- Two main approaches:
 - RSA type signatures [CL02]
 - Based on strong version of RSA assumption
 - Idemix (IBM)
 - DSA type signatures [Brands 99]
 - Based on discrete logarithm problem (more or less)
 - UProve (Microsoft)
 - Also third type based on elliptic curves with pairings [BCKLO8]
 - Less efficient
 - Allows for extra features

- Key tool: Proof of knowledge of discrete log
 - Given Y, g, prove "I know x such that $Y = g^{x}$ "
 - Generalized:
 - Given Y, g, h, prove "I know x, z such that $Y = g^x h^{z^{"}}$
 - Given Y, W, g, h, prove "I know x such that $Y = g^x$ and $Z = h^{x^n}$
 - Prove arithmetic relationships
 - Prove that values are not equal
 - •
 - Prove statements about commitments, signatures, encryptions, etc.



Alice wants to convince service that she knows x, Without revealing any other information

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Roadmap

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Credentials

• Now we have an anonymous credential system. What other issues come up?

- What about misuse of credentials?
 - If everyone is completely anonymous, how do we deal with misuse of privileges?
 - Can we revoke credentials?
 - Can we even tell whose credential to revoke?

Credential Revocation

- Expiration dates
 - Can be embedded in anonymous credentials prove that expiration date > current date
- CRL (Certificate Revocation List)
 - List of all revoked certificates
 - Verifier can check that presented cert is not on list
 - Anonymous CRLs? : How to check that the credential is not on the revoked list without compromising privacy?

Anonymous CRLs

- Option 1:
 - Verifier gives Alice CRL
 - Alice proves that her credential is not on the list (for each value on the list, prove that her value is different)
- Option 2:
 - We can do this more concisely using *accumulators*
 - Issuer publishes accumulator single value that encapsulates all revoked credentials (or all good credentials)
 - Users, given updates to CRL (or list of all good credentials), can give short proof they are not on CRL (or they are on whitelist).

How do we deal with misuse of privileges? (How do we tell who to revoke?)

- Depends how we define misuse:
 - Simple type: reused one-use token
 - Tried to vote twice in a poll
 - Tried to spend transit token twice
 - More complex scenarios
 - Trust a judge to determine misuse

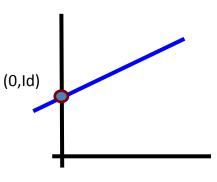
- Credentials meant to be used only once (or fixed number of times)
 - Subway tokens
 - Electronic currency (e-cash)
 - Movie tickets
 - Access passes for online service





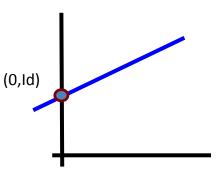
- Service records "serial number" on every token used
- As long as each token is only used once
 - user is anonymous
 - multiple tokens used by the same user are unlinkable
- If token is used twice, identity of user is revealed.
- Previous work [Chaum83, CFN90,... CHL05,... BCKL09]

- Anything digital can be copied!
- Why can't Alice just copy her credential, and give one copy to Bob and the other to Carol?
 - Efficient Solution: offline e-cash [CFN90]
 - Cred includes (T, Id) unknown to Org
 - Id: the identifying info for the user
 - T: the slope of a line with f(0)=Id
 - When cred is used it includes (R, D):
 - R: transaction information (station name, timestamp, etc)
 - D: Doublespending tag (f(R)).
 - » (R,D) and (R',D') gives Id



(R, D)

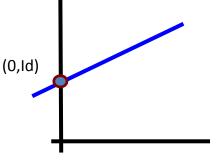
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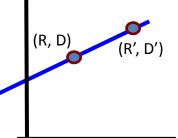


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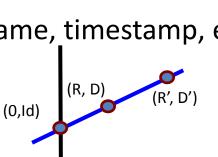
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(0,Id)

How do we deal with misuse of privileges? More complex scenarios

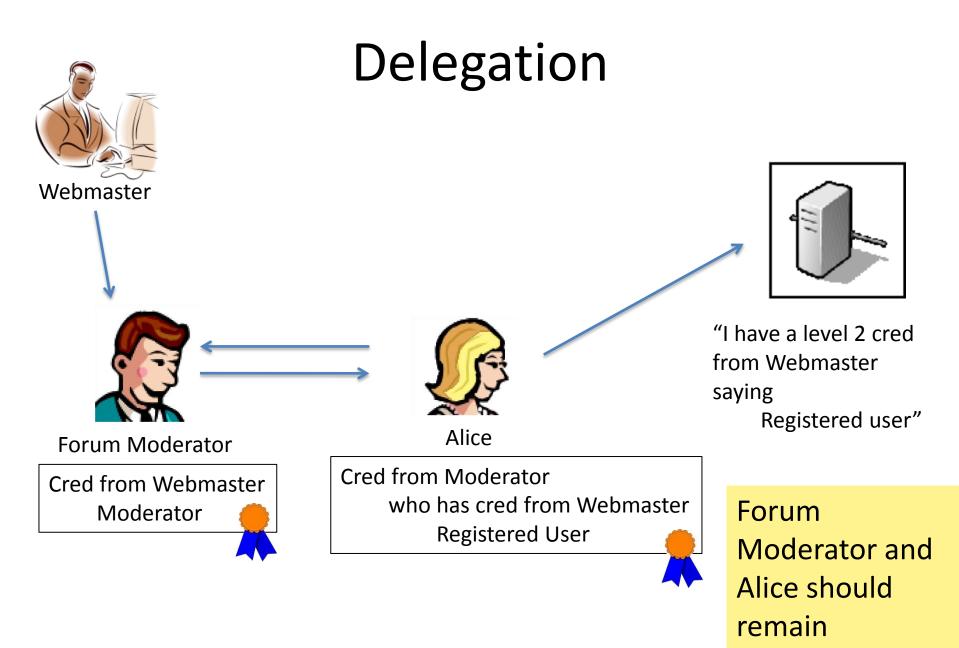
- Trusted judge (anonymity revocation authority)
 - Alice also sends encryption of her identity under judge's public key (*Identity escrow*)
 - In case of misuse,
 - Service gives encryption to judge
 - If judge agrees credential was misused, it can decrypt and find Alice's identity
- Disadvantage: users have no anonymity w.r.t. revocation authority
 - Judge must be trusted
- Advantage: very flexible
- Techniques: Verifiable encryption

Roadmap

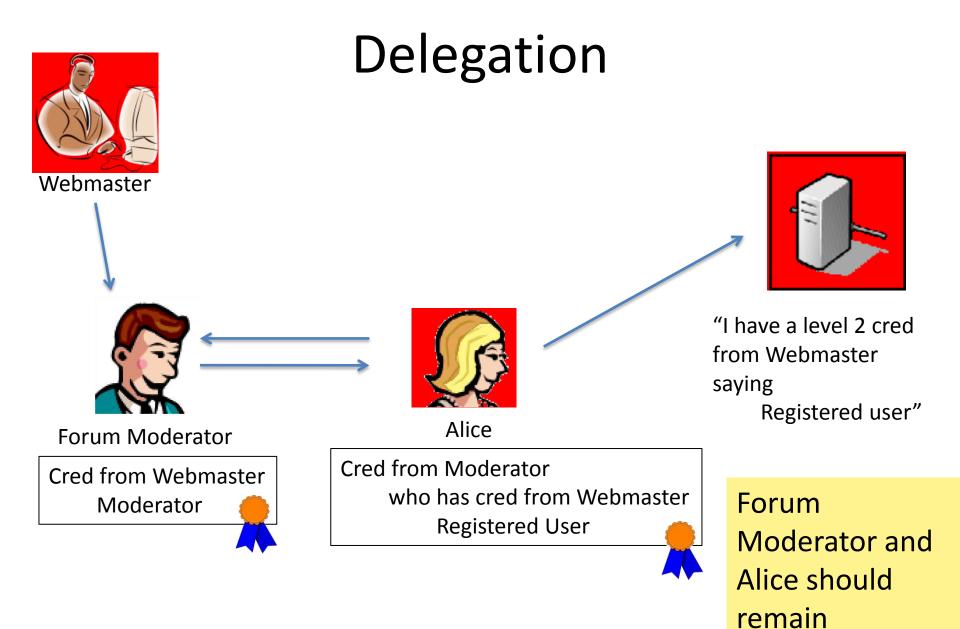
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Other Features

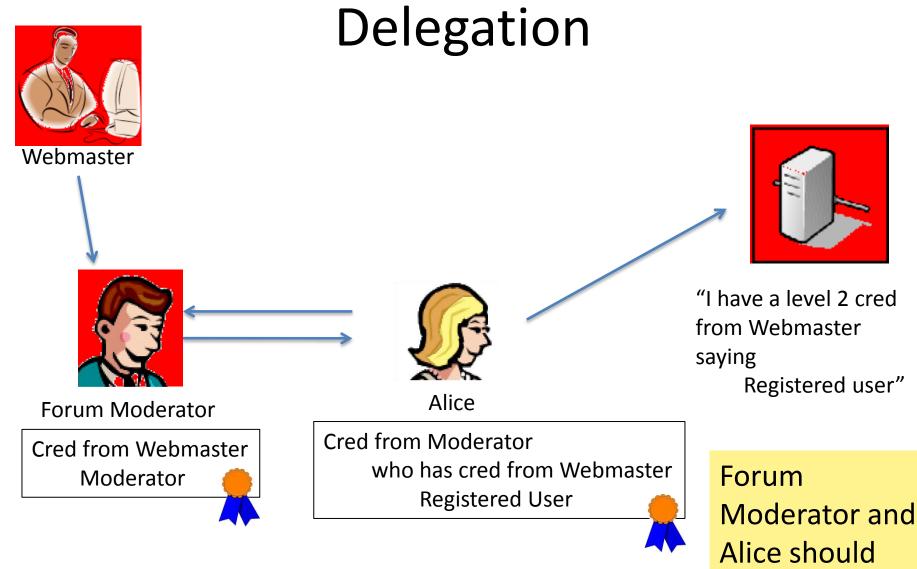
- Log of all valid users and their credentials?
- Post an anonymous message with proof of a credential?
- Non-interactive credentials (Signatures)
 - Challenge: proof needs to be one message
 - Non interactive Zero Knowledge proof
 - Fiat-Shamir (using hash as challenge)
 - Or recent proof techniques based on special elliptic curves



anonymous

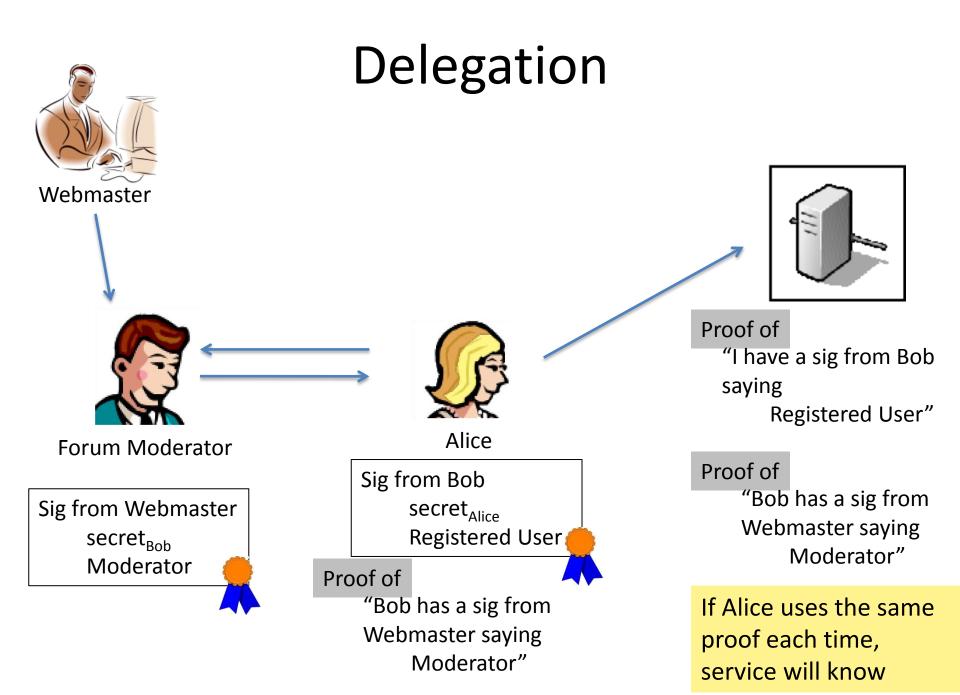


anonymous

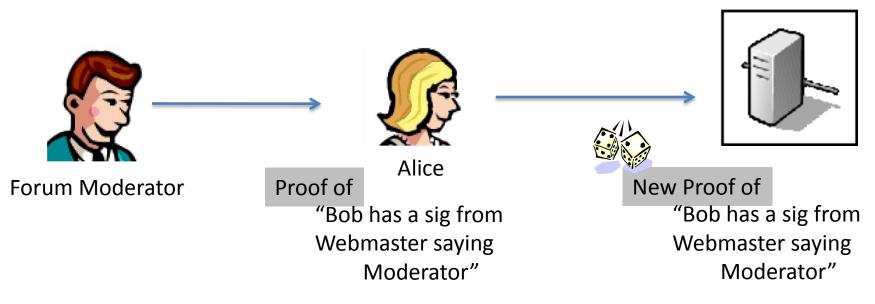


remain

anonymous



Randomizable proofs



- Can we do this?
 - Not clear with traditional techniques
 - Need proofs with special properties

Delegating Credentials

Randomizable proof system



- Elliptic curve with pairings based proofs
 [GOS06,GS08] satisfy this property
- Delegatable Anonymous Credentials [BCCKLS09]
 - Requires some additional techniques
- In progress: delegatable one-time credentials (i.e. transferrable e-cash) [CCKR]

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Other issues

- How do you tie a digital credential to a real world person/identity?
 - Harder when you add anonymity
 - Circular encryption, smart card, POK of credit card number
- Safety in numbers:
 - What if the issuer only ever issues one credential?
 - Even with anonymous credentials, if yours is the only credential issued, issuer will know when you show it
- Adoption will anyone ever use this?

– Do people care enough about privacy?

Questions