11.Application Servers (aka.TPM onitors)

CSEP 545 Transaction Processing Philip A . Bernstein

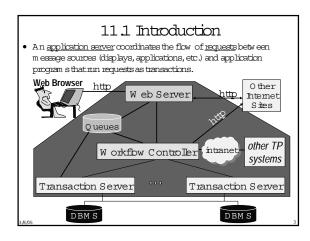
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Reading

Textbook: Chapter 2, and Chapter 3 sections 3.1 - 3.3

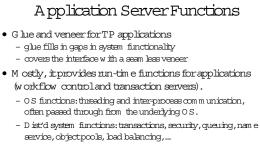
Outline

- 1. Introduction
- 2.Two-Tiervs.Three-Tier
- 3.W eb Servers
- 4. Transaction Bracketing
- 5. Processes and Threads
- 6.Remote Procedure Call



A pplication Server C on ponents Web Browser A smartdevice, with form s, menus, input validation Web server Performs s front-end work, e.g., security, data caching,.... "Calls" the web page associated with the URL, which in turn calls a workflow controller Workflow controller Calls Start, Commit, and Abort App logic that transform s request (autom atic ban payment,

- App logic that transform singuest (automatic ban payment, money transfer) into calls on basic objects (ban, account).
 Som etim es called business rules.
- Transaction server
 Business objects (custom er, account, ban, teller)
- DBM S Database M anagem ent System

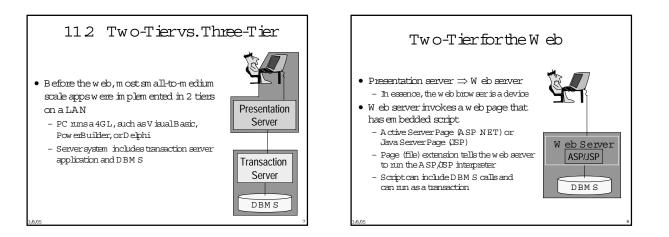


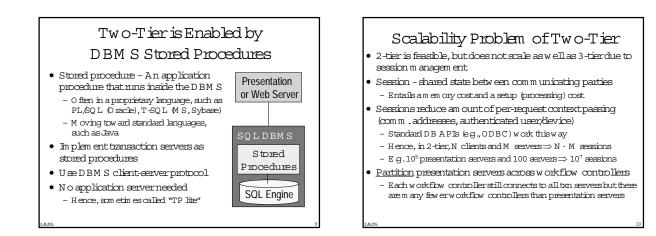
- Portal functions: shopping cart, catalog m gm t, personalization ...
- Provides som e application developm ent tools
- Provides system m gmt for the running application.

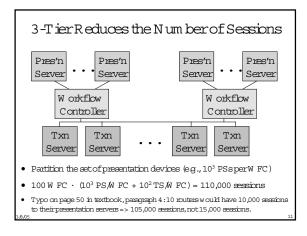
Application Server Products

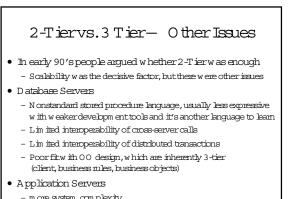
- Apple W ebO bjectsBEA Tuxedo
- BEA TUXEdo
- BEA Webbgic
- Macromedia ColdFusion
- HP (Tandem) Pathway
- HP (DEC) ACM S
- BM CICS
- IBM MS/DC
- IBM W ebsphere

- Iona iPortal App Server
- M icrosoft NET Enterprise Services (form enly COM + , M S Transaction Server (M TS))
- NovellSilverStream
- O racle Application Server
- Sun 0 ne Application Server
- Sybase EA Server
- Synder in Server
- And m any others. See serverw atch.internet.com









How the Web Changed Things

- Presentation server \Rightarrow W eb server
- All requests have to pass through a W eb server
 In 2-tier, each W eb serverneeds assions to all D B servers
 Session reduction by workflow control is less critical but still useful
- W orkflow control is still useful for request mgmt - Calling Start, Commit, and Abort
 - Encapsulating business rules that transform each request into calls on basic objects

113WebServers

- Presentation independence application is independent of the display device used
 - Today, this is via http and htm l
 - In the past, itw as via a display controllerorm iddle-tier m inicom puterwhose presentation functions insulated the rest of the back-end system from different device types
- W eb server perform s presentation functions:
 - Gathering input DB caching
 - Validating input Authentication
- They also do som e basic request routing
 Constructing requests Invoking applications
- Examples IIS (MS), Apache, Netscape Server

Gathering Input

 G athering input - Select transaction type (m enu item, etc.), and fill in a form (request's parameters)
 Today, W eb form s, m oving to XM L (X Form s, X SLT,...)

- 30 year evolution of presentation devices
 - Teletype, character-at-a-tin e term inal (async), block-m ode term inal (IBM 3270)
 - Specialized devices ATM s, barcode readers, gas pum ps, robots, credit card authorization, cash registers, ticket printers, etc.
 - 4GL on a PC ActiveX controls accessed from V isualBasic (VB), PowerBuilder, D elphi, etc.

Caching

- Every process-to-process callhas a cost - Adds to response time and consum es resources
- Use a cache in W eb server to avoid calling workflow controller or DB system
- Cache popular read-only data that need not be refreshed frequently
- E g., catalog item s, sale item s, coverpage at an auction site, recentnew s, etc.
- A lso, data required for input validation info

Input V alidation

- V alidate input against <u>locally cached</u> tables
 E g., product types, department num bers
- A voids wasting communications and server resources for obvious input enors
 - Few er round-trips to the DBM S
 - And faster feedback to the end user
- "Cache" is part of the web page
 - Listboxes, script
 - Cache size is a factor (it affects page access tin e)

A uthentication

- <u>A uthentication</u> determ ining the identity of a user and/or display device
 - Clientsystem (eg., PC) m ay do authentication, but the server usually does it too (doesn't trust clients)
 - Encrypt the w ize to avoid w izetapping and spoofing (on the W eb, use https = http over secure socket layer)
- <u>Geographical entitlem ent</u> check that a particular device is allow ed access (e.g., security trading room)
- Need system m gm t functions to create accounts, initialize passwords, bracket hours of access (sim plify it using a <u>role</u> abstraction)

Constructing Requests

- A request includes
 - U serid for authorization and personalization
 - Device id where to send a reply
 - Device type -whatmessage types can it understand?
 - ObjectID in a OO setting
 - RequestID to ask later about request status & to link a reply
 - Request type name of transaction type requested
 - Request-specific param eters
- Can be com bined with protocol header (e.g., http header)

Application Invocation

- Request arrives as an http m essage.
- N eed to calla program (i.e. a W FC), to perform the request
- Common Gateway Interface
 - Write a script, store it as a file in ogi-bin
 - W eb server creates a process to execute the request (S bw !!)
- ISAPI (Microsoft) and NSAPI (Netscape)
 - W eb server calls an in-proc dll instead of creating a process
 - Webserver can cache the dll
 - M ore com plex program m ing m odel, butm uch faster
- A ctive Server Pages and Java Server Pages
- Offers the perform ance of ISA PIw ith program m ability of CG I

Load Balancing

• W eb servers enable <u>scale out</u>, so you can just add m ore serverboxes to handle m ore load.

- To sim plify this problem
 - Ensure all web servers are identical (no server-specific state)
 don'tretain client state on web servers (hard to avoid ...)
 - Random ly assign requests to servers (e.g., use an IP sprayer)
 - A void sending requests to a failed server

11.4 Transaction Bracketing

- For the mostpart, W orkflow Controllers (W FC) and Transaction Servers are just plain old server program s
- The main W FC differentiating features
 - Brackets transactions (issues Start, Comm it, and Abort)
 - Handles Aborts (returns cause of the Abort)
 - Does not access the DBM S

Nested Transaction Calls

- W hat does Start do, when executed within a txn?
 - 1. it starts an independent transaction, or
 - 2.itdoesnothing,or
 - 3. it increments a nested transaction count (which is decremented by each commit and abort), or
 - 4. it starts a sub-transaction
- (2) and (3) are common.
 - Enables a transaction-bracketed program to be called by another transaction
- (1) im plies Be Careful!

Transparent Transaction B racketing

- Transaction-hood is a property of the app component.
- In COM +, a class is declared:
 - requires new callee always starts a new transaction
 - required if caller is in a transaction, then nun callee in caller's transaction, else start a new transaction
 - supported if caller is in a transaction, then nun caller in caller's transaction, else nun outside of any transaction
 not supported don't nun in a transaction
- Caller can create a transaction context, which supports Committand A bort (chained model).
 - Callee issues SetComplete when it's done and willing to commit, or SetA bort to abort.

TransparentTxn Bracketing (cont'd)

- EJB and J2EE
 - Im plem ents COM + technology in Java:RequiredNew , Required, Supported, NotSupported
 - It cam e later, so there are two additions.
 - M andatory If caller is in a transaction, then run the callee in that transaction, else raise an exception
 - N ever If caller is in a transaction, then raise an exception

Runtim e Library Support

TP services require runtim e library support
 M ay orm ay not be language-specific

• Language-specific

- Java 2 Enterprise Edition (J2EE , form enly Enterprise Java Beans)
- Encapsulates runtim e library as a container object.
- BEA Weblogic, IBM Websphere,... .
- O blerexam ples are Tandem Pathway (Screen COBOL) and Digital's ACM Sxp (Structured Txn Defn Lang)
- Language-independent runtim e library
- MSCOM+, IBM 'sCICS, Oracle App Server,...

Savepoints

- Savepoint a point in a program where an application saves all its recoverable state
- Can <u>restore</u> a savepoint within the transaction that issued the savepoint. (It's a partial rollback.)
- SQLDBM Ssue them to support atom ic SQL statem ents.
 Start:
 - get-request; Savepoint("B"); . . .; if (error) {Restore("B"); ...; Commit;}
 - Commit:
- Savepoints are not recoverable. If the system fails or the transaction aborts, the txn is completely undone.

11.5 Processes and Threads

- Application Server architecture is greatly affected by

 which components share an address space
 - how m any control threads per address space
- TP grew up in the days of batch processing, and reached maturity in the days of tim esharing.
- TP users learned early that a process-per-user fails:
 - Toomuch contextswitching
 - Toom uch fixed m em ory overhead perprocess
 - Process peruserperm achine, when distributed
 - Som e 0 S functions scan the list of processes
 Load control is hard
 - Load control is hard

M ultithreading

- Have multiple threads of control in an address space
- U sed to be a m ajor A pplication Server feature

 A pplication Server sw itches threads when app calls a
 A pplication Server function thatblocks
- Now, mostOS's support it natively
- Can run a process's threads on different processors (SM P)
- W hetherat the useror OS level,
 - multithreading has few erprocesses and less context switching
 - but little protection between threads and a server failure affects $\ensuremath{\mathfrak{m}}$ any transactions

M apping Servers to Processes

- Presentation/W eb servers, workflow controllers, and transaction servers are multithreaded servers
- Costs 1500 25,000 instructions per process call, vs. 50 instructions per local procedure call...
 - but it scales, with flexible configuration and control

11.6 Remote Procedure Call

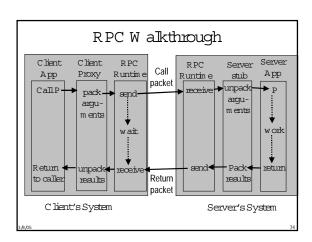
- Program calls remote procedure the same way it would call a local procedure
- Hides certain underlying com plexities
 - communications and message ordering errors
 - data representation differences between program s
- TransactionalRPC
 - Ideally, Start returns a transaction ID that's hidden from the caller
 - Procedures don't need to explicitly pass transaction id's.
 - Easier and avoids enors

Binding

- Interface definitions
 - From app or written in an interface definition language (ID L)
 - com piles into Proxy and Stub program s
 - Client calls the Proxy (representing the server)
 - Stub calls the Server (represents the client on the server)
- Marshaling
 - proxy <u>m arshals</u> (sequentially lays out) calling param eters in a packet and decodesm arshaled return values
- stub decodes m arshaled calling param s and m arshals return param s
- Communications binding
- C lient finds the server location via a directory service, based on servernam e and possibly a parameter value
- To load balance across identical servers, random ly choose a server

Binding (cont'd)

- The binding process has security guarantees
 - The clientm usthave privileges to bind to the serverThe clientm ustknow it's binding to an appropriate server to
 - avoid being spoofed - E.g. client and server authenticate each other during session
 - creation, and m aybe per-access too



Perform ance

- There are basically 3 costs
 - marshaling and unmarshaling
 - RPC runtin e and network protocol
 - physicalw ire transfer
- In a LAN , these are typically about equal
- Typical com m ercial num bers are 10-15K m achine instructions
- Can do much better in the local case by avoiding a full context sw itch

Stateful Applications

- Som etim es an application m aintains state on client's behalf, possibly across bansactions. E g.,
 - Server scans a file. Each tin e it hits a relevant record it returns it. N ext call picks up the scan where it left off.
 - Webservermaintains a shopping basket or itinerary, etc.

A pproach 1: client passes state to server on each call, and server returns it on each reply. Server retains no state.

- Doesn'twork well for TP, because there's too much state
- Note that transaction id context is handled this way.

StatefulServers (cont'd)

A pproach 2 : serverm aintains state, indexed by client id (ban id or cookie). Client's laterRPCsm ustgo to sam e server.

- If the client fails, serverm ustbe notified to release client's state or deallocate based on timeout
- Fortransaction RPC, encapsulate context as a (volatile) resource.
 D elete the state at com m it/abort. O rpossibly, m aintain state across transaction boundaries, but reconstruct it after system failure.
- Eg., COM +: Client can call a server objectm any times
- Client creates server object, which retains state across RPCs
- SetComplete (or SetA bort) by server app says that transaction can be comm itted (or aborted) and state can be deleted
- <u>EnableCommit</u> (or<u>DisableCommit</u>) by server app says transaction can (or cannot) be committed by client and don't delete server state

FaultTolerance

- If a client doesn't receive a reply within its tim coutperiod
- RPC runtime can send a "ping" for non-idem potent calls
- A flerm ultiple pings, it return an enor.
- For idem potent calls, RPC nuntim e can retry the call (server interface definition can say whether it's idem potent)

W eb Services

- A new round of distributed computing standards to enable interoperation on the Internet
- SOAP RPC with XM L as marshalling form at and W SDL as interface definition
- UDDI-directory for finding W eb Service descriptions
- W S-Transaction 2PC
- W S-Security, W S-Coordination, W S-Routing,...
- www.ws-iorg

Sum m ary

- Scalability 2 vs. 3 tier, sessions, stored procedures
- W eb Server gathering input, validating input, caching, authentication, constructing requests, invoking applications, load balancing
- Transaction bracketing transparency, nesting, exceptions, request integrity, savepoints
- Serverprocesses threads
- RPC binding, stateful servers