CSE P 501 – Compilers

x86-64, Running MiniJava,
Basic Code Generation and Bootstrapping
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Adminstrivia

• Codegen part of project posted now
  – Final push deadline Sun. 3/13, 11pm – no extensions
  – Project report due next night (will post info later)
  – Project meetings Tue. & Wed. 3/15-16 – scheduling later
  – Grading of previous parts – long story...

• HW4 out (now | soon) [pick one]
  – Review problems on analysis/opt/IRs/SSA
  – Due Monday night, 2/29 – no extensions; solutions posted next day

• Exam next Thur. 3/3 @UW and @Microsoft, 6:30-8:00
  – Topic list & old exams on web now
  – OK to use textbook + slides + topic list + 2 handouts (x86-64, MiniJava grammar) during exam; laptops ok if only used to access items linked from exam resources or topics pages
Codegen/Execution for MiniJava

• Build on Codeshape lectures

• More specifics here about execution environment and compiler code gen routines
Running MiniJava Programs

• To run a MiniJava program
  – Space needs to be allocated for a stack and a heap
  – %rsp and other registers need to have sensible initial values
  – We need some way to allocate storage (new) and communicate with the outside world
Bootstraping from C

• Idea: take advantage of the existing C runtime library
• Use a small C main program to call the MiniJava main method as if it were a C function
• C’s standard library provides the execution environment and we can call C functions from compiled code for I/O, malloc, etc.
Assembler File Format

• Compiler output is an assembly-language source file (.s)
• GNU syntax is roughly this (asm samples included in starter src/runtime directory, links to gcc/as refs. on project page)

```
# comments start with # and go to end of the line
.text # code segment
.globl asm_main # start of compiled static main
<generated code> # repeat .text/.data as needed
asm_main: # start of compiled “main”
...
.data
<generated method tables>
# repeat .text/.data as needed
...
end
```
External Names

• In a Linux environment, an external symbol is used as-is (xyzzy)
• In Windows and OS X, an external symbol xyzzy is written in asm code as _xyzzy (leading underscore)
• Adapt to whatever environment you’re using
  – but what you turn in should run on CSE linux using the Linux conventions
  – Exceptions? Check with course staff first
Generating .asm Code

• Suggestion: isolate the actual compiler output operations in a handful of routines
  – Modularity & saves some typing
  – Possibilities
    // write code string s to .asm output
    void gen(String s) { ... }
    // write “op src,dst” to .asm output
    void genbin(String op, String src, String dst) { ... }
    // write label L to .asm output as “L:”
    void genLabel(String L) { ... }
  – A handful of these methods should do it
A Simple Code Generation Strategy

• Goal: quick ‘n dirty correct code, optimize later if time
• Traverse AST primarily in execution order and emit code in visitor methods
  – Visitor may traverse the tree in ad-hoc ways depending on the sequence in which parts need to be written to the code assembler file
• Treat the x86 as a 1-register machine with a stack for additional intermediate values
  – Ugly code, but will work – better later if there’s time
(The?) Simplifying Assumption

• Store all values (reference, int, boolean) in 64-bit quadwords
  – Natural size for 64-bit pointers, e.g., object references (variables of class types)
  – C’s “long” size for integers
    • Better to use int64_t or uint64_t in C code to guarantee size (declared in <stdint.h>)
x86 as a Stack Machine

• Idea: Use x86-64 stack for expression evaluation with %rax as the “top” of the stack
• Invariant: Whenever an expression (or part of one) is evaluated at runtime, the generated code leaves the result in %rax
• If a value needs to be preserved while another expression is evaluated, push %rax, evaluate, then pop when first value is needed
  – Remember: *always pop what you push*
  – Will produce lots of redundant, but correct, code
• Examples below follow code shape examples, but with some more details about code generation
Example: Generate Code for Constants and Identifiers

Integer constants, say 17

\[
\text{gen}(\text{movq} \ \$17,\%rax)
\]

• leaves value in %rax

Local variables (any type – int, bool, reference)

\[
\text{gen}(\text{movq} \ \text{varoffset}(\%rbp),\%rax)
\]
Example: Generate Code for exp1 + exp1

Visit exp1
  – generate code to evaluate exp1 with result in %rax
  gen(pushq %rax)
  – push exp1 onto stack

Visit exp2
  – generate code for exp2; result in %rax
  gen(popq %rdx)
  – pop left argument into %rdx; clean up stack
  gen(addq %rdx,%rax)
  – perform the addition; result in %rax
Example: var = exp; (1)

Assuming that var is a local variable:

Visit node for exp

• Generates code to eval exp and leave result in %rax

gen(movq %rax,offset_of_variable(%rbp))
Example: var = exp; (2)

If var is a more complex expression (object or array reference, for example)

visit var

  gen(pushq %rax)
    • push reference to (i.e., address of) variable or object containing variable onto stack

visit exp

  • leaves rhs value in %rax

  gen(popq %rdx)

  gen(movq %rax,appropriate_offset(%rdx))
Example: Generate Code for obj.f(e1,e2,...en)

In principal the code should work like this:

Visit obj
  • leaves reference to object in %rax
  gen(movq %rax,%rdi)
    • “this” pointer is first argument

Visit e1, e2, ..., en. For each argument,
  • gen(movq %rax,correct_argument_register)
generate code to load method table pointer located at
0(%rdi) into some register, probably %rax
generate call instruction with indirect jump
Method Call Complications

• Big one: code to evaluate any argument might clobber argument registers (i.e., method call in some parameter value)
  – Possible strategy to cope on next slides, but feel free to do something better

• Other one: what if a method has too many parameters?
  – OK for P 501 to assume that all methods have ≤ 5 parameters plus “this” – do better if you want
Method Calls in Parameters

• Suggestion to avoid trouble:
  – Evaluate parameters and push them on the stack
  – Right before the call instruction, pop the parameters into the correct registers

• But....
Stack Alignment (1)

• Above idea hack works provided we don’t call a method while an odd number of parameter values are pushed on the stack!
  – (violates 16-byte alignment on method call…)

• We have a similar problem if an odd number of intermediate values are pushed on the stack when we call a function while evaluating an expression
  – (We might get away with it if it only involves calls to our own generated, not library, code, but it would be wrong* to do that)
    *i.e., might “work”, but not the right way to solve the problem
Stack Alignment (2)

• Workable solution: keep a counter in the code generator of how much has been pushed on the stack. If needed, emit extra `gen(pushq %rax)` to align stack before generating a call instruction
  – Be sure to pop it after!!

• Another solution: make stack frame big enough and use `movq` instead of `pushq` to store arguments and temporaries
  – Will need some extra bookkeeping to keep track of how much to allocate and how temps are used
Sigh...

- Multiple registers for method arguments is a big win compared to pushing on the stack, but complicates our life since we do not have a fancy register allocator.
- Feel free to do better than this simple push/pop scheme – but remember, simple and works wins over fancy and broken.
Code Gen for Method Definitions

• Generate label for method
  Classname$methodname:

• Generate method prologue
  Push %rbp, copy %rsp to %rbp, subtract frame size from %rsp

• Visit statements in order
  – Method epilogue is normally generated as part of each return statement (next)
  – In MiniJava the return is generated after visiting the method body to generate its code
Registers again...

• Method parameters are in registers
• But code generated for methods also will be using registers, even if there are no calls to other methods
• So how do we avoid clobbering parameters?
• Suggestion: Allocate space in the stack frame and save copies of all parameter registers on method entry. Use those copies as local variables when you need to reference a parameter.
Example: return exp;

- Visit exp; this leaves result in %rax where it should be
- Generate method epilogue (copy %rbp to %rsp, pop %rbp) to unwind the stack frame; follow with ret instruction
Control Flow: Unique Labels

• Needed in code generator: a String-valued method that returns a different label each time it is called (e.g., L1, L2, L3, ...)

  – Improvement: a set of methods that generate different kinds of labels for different constructs (can really help readability of the generated code)

    • (while1, while2, while3, ...; if1, if2, ...; else1, else2, ...; fi1, fi2, ... .)
Control Flow: Tests

• Recall that the context for compiling a boolean expression is:
  – Label or address of jump target
  – Whether to jump if true or false

• So the visitor for a boolean expression should receive this information from the parent node
Example: while(exp) body

• Assuming we want the test at the bottom of the generated loop...
  gen(jmp testLabel)
  gen(bodyLabel:)
  visit body
  gen(testLabel:)
  visit exp (condition) with target=bodyLabel and sense="jump if true"
Example: exp1 < exp2

• Similar to other binary operators
• Difference: context is a target label and whether to jump if true or false

• Code
  visit exp1
  gen(pushq %rax)
  visit exp2
  gen(popq %rdx)
  gen(cmpq %rdx, %rax)
  gen(condjump targetLabel)
    • appropriate conditional jump depending on sense of test
Boolean Operators

&& (and || if you add it)
   – Create label(s) needed to skip around parts of the expression
   – Generate subexpressions with appropriate target labels and conditions

!exp
   – Generate exp with same target label, but reverse the sense of the condition
Reality check

• Lots of projects in the past have evaluated all booleans to get 1 or 0, then tested that value for control flow

• Would be nice to do better (as above), but “simple and works…”
Join Points

- Loops and conditional statements have join points where execution paths merge
- Generated code must ensure that machine state will be consistent regardless of which path is taken to get there
  - i.e., the paths through an if-else statement must not leave a different number of words pushed onto the stack
  - If we want a particular value in a particular register at a join point, both paths must put it there, or we need to generate additional code to move the value to the correct register
- With a simple 1-accumulator model of code generation, this should usually be true without needing extra work; with better use of registers it becomes a bigger issue
Bootstrap Program

• The bootstrap is a tiny C program that calls your compiled code as if it were an ordinary C function

• It also contains some functions that compiled code can call as needed
  – Mini “runtime library”
  – Add to this if you like
    • Sometimes simpler to generate a call to a hand-written library routine instead of generating in-line code – implementer tradeoff
Bootstrap Program Sketch

```c
#include <stdio.h>
extern void asm_main( ); /* compiled code */
/* execute compiled program */
void main( ) { asm_main(); } /* write x to standard output */
void put(int64_t x) { ... }
/* return a pointer to a zeroed block of memory at least nBytes large (or null on failure) */
char* mjcalloc(size_t nBytes) { return calloc(1,nBytes); }
```
Main Program Label

• Compiler needs special handling for the public static void main method label
  – Label must be the same as the one declared extern in the C bootstrap program and declared .globl in the .s asm file
  – asm_main used above
     • Could be changed, but probably no point
     • Why not “main”? (Hint: where is the real main?)
Interfacing to “Library” code

- Trivial to call “library” functions
- Evaluate parameters using the regular calling conventions
- Generate a call instruction using the function label
  - (External names need leading _ in Windows, OS X)
  - Linker will hook everything up
System.out.println(exp)

MiniJava’s “print” statement
<compile exp; result in %rax>
movq %rax,%rdi   # load argument register
call put        # call external put routine

• If the stack is not kept 16-byte aligned, calls to external C or library code can cause a runtime error (will cause error halt on OS X)
And That’s It…

- We’ve now got enough on the table to complete the compiler project
- Coming Attractions
  - Back end (instruction selection and scheduling, register allocation)
  - and more…
    (including an exam)