

Evolution of the Internet Internet started with best-effort point-to-point communication TCP for reliable delivery Extensions IP Multicast for many-to-many communication TCP-SACK to optimize TCP performance Mobile IP for seamless mobility

Why haven't they been deployed yet

- 1 Multicast need router upgrades
 - Router vendors must implement them
 - ISPs must deploy the new routers
- 1 Many technical and business concerns
 - Protocol complexity: keep the network simple
 - Network heterogeneity
 - Layering complex transport protocols

What about end-to-end protocols

- 1 TCP-SACK is technically superior
 - Yet not widely deployed
- 1 Backward compatibility
- 1 Upgrade hell
 - Windows 98/98/NT still account for 7.5% of deployed systems

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Examples of successful evolution

- 1 Domain Name System
 - Replacement for hosts.txt
- 1 TCP Congestion Control
 - [Jacobson '88]
- 1 Cell-phone networks
 - Upgrading from analog to digital systems

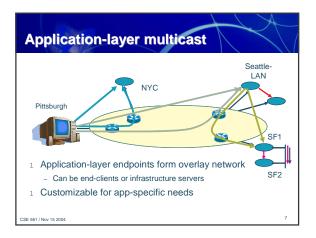
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An alternative approach

- 1 Keep the network unchanged
 - Layer new protocols on top
- 1 Application-layer multicast
 - An alternative evolutionary approach
 - Instead of IP-layer multicast, deploy it entirely at the end-points
 - E.g., End-system Multicast, Scattercast, Overcast

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Other overlay networks	
 RON: Resilient Overlay Network Move routing control toward end-systems Route using overlays; bypass ISP policies Akamai: Content delivery network Intelligent distribution of content 	
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Why are overlays attractive

- 1 Deploying new network services is hard
 - The standards challenge
 - Upgrade inertia
- 1 Get the service provider out of the loop
 - Makes deployment "trivial"
 - The overlay IS the infrastructure

i3: A general-purpose routing overlay

- 1 Generalized rendezvous/routing service
 - Sender sends packets to logical id
 - Receiver inserts "trigger" indicating interest in the logical id
- 1 Can be used to implement multicast, anycast, mobility, service composition and more

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Protocol description

- 1 IDs have m bits
- 1 Matching algorithm
 - Must match at least k bits
 - Pick longest prefix match from the IDs that pass above test
 - Forward to all matching triggers
 - Chain of forwarding using stacks of identifiers

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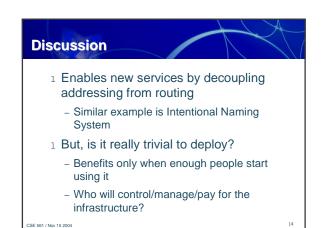
Usage Multicast All receivers subscribe to the same ID In practice, build a hierarchy of IDs Anycast All anycast receivers pick the same first k bits for their ID and random (m-k) bits

- Sender picks an ID whose first k bits match the anycast group ID
- 1 Mobility

- Trivial: receiver updates its trigger when it moves



- 1 Sender-driven or receiver-driven
- 1 Use stacks of identifiers
- 1 Can be used to implement heterogeneous multicast



i3-like overlays can cause problems

- 1 Service level agreements
 - Routing overlays violate SLAs
- 1 Traffic engineering
 - Becomes harder in the face of overlays
- 1 Churn

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- Potentially dynamic membership
- Management overhead may overwhelm the network

Should network providers care?

- 1 Providing just connectivity is a losing proposition
- ${\scriptstyle 1}$ Overlays enable new services \Rightarrow potential for more \$
- 1 Improved performance, reduced traffic in core ⇒ potential for saving \$\$\$
- Control over monitoring and engineering the overlay

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