Greedy Algorithms

- Hard to define exactly but can give general properties
  - Solution is built in small steps
  - Decisions on how to build the solution are made to maximize some criterion without looking to the future
    - Want the ‘best’ current partial solution as if the current step were the last step
- May be more than one greedy algorithm using different criteria to solve a given problem
Greedy Algorithms

- Greedy algorithms
  - Easy to produce
  - Fast running times
  - Work only on certain classes of problems

- Three methods for proving that greedy algorithms do work
  - Greedy algorithm stays ahead
    - At each step any other algorithm will have a worse value for the criterion
  - Exchange argument
    - Can transform any other solution to the greedy solution at no loss in quality
  - Structural property
Interval Scheduling

- Interval Scheduling
  - Single resource
  - Reservation requests
    - Of form “Can I reserve it from start time $s$ to finish time $f$?”
    - $s < f$
  - **Find**: maximum number of requests that can be scheduled so that no two reservations have the resource at the same time
Interval Scheduling

- Interval scheduling.
  - Job $j$ starts at $s_j$ and finishes at $f_j$.
  - Two jobs compatible if they don't overlap.
  - Goal: find maximum subset of mutually compatible jobs.
Interval scheduling

Formally

Requests 1, 2, ..., n
- request i has start time $s_i$ and finish time $f_i > s_i$
- Requests i and j are compatible iff either
  - request i is for a time entirely before request j
    - $f_i \leq s_j$
  - or, request j is for a time entirely before request i
    - $f_j \leq s_i$
- Set A of requests is compatible iff every pair of requests i,j ∈ A, i≠j is compatible
- Goal: Find maximum size subset A of compatible requests
Greedy Algorithms for Interval Scheduling

What criterion should we try?
- Earliest start time $s_i$
- Shortest request time $f_i - s_i$
- Fewest conflicting jobs
- Earliest finish time $f_i$
Greedy Algorithms for Interval Scheduling

What criterion should we try?

- Earliest start time $s_i$
  - Doesn’t work

- Shortest request time $f_i - s_i$
  - Doesn’t work

- Even fewest conflicts doesn’t work
  
- Earliest finish time $f_i$
  - Works
Greedy Algorithm for Interval Scheduling

R ← set of all requests
A ← ∅
While R ≠ ∅ do
    Choose request i ∈ R with smallest finishing time f_i
    Add request i to A
    Delete all requests in R that are not compatible with request i

Return A
Claim: \( A \) is a compatible set of requests and these are added to \( A \) in order of finish time

- When we add a request to \( A \) we delete all incompatible ones from \( R \)

Claim: For any other set \( O \subseteq R \) of compatible requests then if we order requests in \( A \) and \( O \) by finish time then for each \( k \):

- If \( O \) contains a \( k \)th request then so does \( A \) and the finish time of the \( k \)th request in \( A \), is \( \leq \) the finish time of the \( k \)th request in \( O \), i.e. “\( a_k \leq o_k \)” where \( a_k \) and \( o_k \) are the respective finish times

Enough to prove that \( A \) is optimal
Inductive Proof of Claim: $a_k \leq o_k$

- **Base Case:** This is true for the first request in $A$ since that is the one with the smallest finish time.

- **Inductive Step:** Suppose $a_k \leq o_k$
  - By definition of compatibility
    - If $O$ contains a $k+1^{st}$ request $r$ then the start time of that request must be after $o_k$ and thus after $a_k$.
    - Thus $r$ is compatible with the first $k$ requests in $A$.
    - Therefore
      - $A$ has at least $k+1$ requests since a compatible one is available after the first $k$ are chosen.
      - $r$ was among those considered by the greedy algorithm for that $k+1^{st}$ request in $A$.
    - Therefore by the greedy choice the finish time of $r$ which is $o_{k+1}$ is at least the finish time of that $k+1^{st}$ request in $A$ which is $a_{k+1}$.
Implementing the Greedy Algorithm

- Sort the requests by finish time
  - $O(n \log n)$ time
- Maintain current latest finish time scheduled
- Keep array of start times indexed by request number
- Only eliminate incompatible requests as needed
  - Walk along array of requests sorted by finish times skipping those whose start time is before current latest finish time scheduled
  - $O(n)$ additional time for greedy algorithm
Sort jobs by finish times so that $0 \leq f_1 \leq f_2 \leq \ldots \leq f_n$.

\begin{align*}
A &\leftarrow \emptyset \\
\text{last} &\leftarrow 0 \\
\text{for } j = 1 \text{ to } n \{ \\
&\quad \text{if } (\text{last} \leq s_j) \\
&\quad \quad A \leftarrow A \cup \{j\} \\
&\quad \quad \text{last} \leftarrow f_j \\
\} \\
\text{return } A
\end{align*}
Scheduling All Intervals: Interval Partitioning

- Interval partitioning.
  - Lecture \( j \) starts at \( s_j \) and finishes at \( f_j \).
  - Goal: find minimum number of classrooms to schedule all lectures so that no two occur at the same time in the same room.

- Example: This schedule uses 4 classrooms to schedule 10 lectures.
Interval Partitioning

- Interval partitioning.
  - Lecture $j$ starts at $s_j$ and finishes at $f_j$.
  - **Goal:** find minimum number of classrooms to schedule all lectures so that no two occur at the same time in the same room.

- **Example:** This schedule uses only 3 classrooms
**Scheduling all intervals**

- **Interval Partitioning Problem**: We have resources to serve more than one request at once and want to schedule all the intervals using as few of our resources as possible.

- Obvious requirement: At least the depth of the set of requests.
**Interval Partitioning: Lower Bound on Optimal Solution**

- **Definition.** The depth of a set of open intervals is the maximum number that contain any given time.

- **Key observation.** Number of classrooms needed \( \geq \) depth.

- **Ex:** Depth of schedule below = 3 \( \Rightarrow \) schedule below is optimal.

- **Q.** Does there always exist a schedule equal to depth of intervals?
A simple greedy algorithm

Sort requests in increasing order of start times \((s_1, f_1), \ldots, (s_n, f_n)\)

For \(i=1\) to \(n\)

\[ j \leftarrow 1 \]

While (request \(i\) not scheduled)

\[ \text{last}_j \leftarrow \text{finish time of the last request currently scheduled on resource } j \]

if \(s_i \geq \text{last}_j\) then schedule request \(i\) on resource \(j\)

\[ j \leftarrow j + 1 \]

End While

End For
Interval Partitioning: Greedy Analysis

- **Observation.** Greedy algorithm never schedules two incompatible lectures in the same classroom.

- **Theorem.** Greedy algorithm is optimal.

- **Proof.**
  - Let $d =$ number of classrooms that the greedy algorithm allocates.
  - Classroom $d$ is opened because we needed to schedule a job, say $i$, that is incompatible with all $d-1$ other classrooms.
  - Since we sorted by start time, all these incompatibilities are caused by lectures that start no later than $s_i$.
  - Thus, we have $d$ lectures overlapping at time $s_i + \epsilon$.
  - Key observation $\Rightarrow$ all schedules use $\geq d$ classrooms. □
A simple greedy algorithm

Sort requests in increasing order of start times \((s_1,f_1),\ldots,(s_n,f_n)\)

For \(i=1\) to \(n\)
\[ j \leftarrow 1 \]
While (request \(i\) not scheduled)
\[ \text{last}_j \leftarrow \text{finish time of the last request currently scheduled on resource } j \]
if \(s_i \geq \text{last}_j\) then schedule request \(i\) on resource \(j\)
\[ j \leftarrow j+1 \]
End While
End For

\(O(n \log n)\) time

May be slow \(O(nd)\) which may be \(\Omega(n^2)\)
A more efficient implementation

Sort requests in increasing order of start times \((s_1, f_1), \ldots, (s_n, f_n)\)

\[d \leftarrow 1\]
Schedule request 1 on resource 1
\[\text{last}_1 \leftarrow f_1\]
Insert 1 into priority queue \(Q\) with key = \text{last}_1

For \(i=2\) to \(n\)

\[j \leftarrow \text{findmin}(Q)\]
if \(s_i \geq \text{last}_j\) then
  schedule request \(i\) on resource \(j\)
  \[\text{last}_j \leftarrow f_i\]
  Increasekey(\(j, Q\)) to \text{last}_j
else
  \[d \leftarrow d+1\]
  schedule request \(i\) on resource \(d\)
  \[\text{last}_d \leftarrow f_i\]
  Insert \(d\) into priority queue \(Q\) with key = \text{last}_d

End For

\(O(n \log n)\) time
Greedy Analysis Strategies

- **Greedy algorithm stays ahead.** Show that after each step of the greedy algorithm, its solution is at least as good as any other algorithm's.

- **Exchange argument.** Gradually transform any solution to the one found by the greedy algorithm without hurting its quality.

- **Structural.** Discover a simple "structural" bound asserting that every possible solution must have a certain value. Then show that your algorithm always achieves this bound.
Scheduling to Minimize Lateness

- Scheduling to minimize lateness
  - Single resource as in interval scheduling but instead of start and finish times request $i$ has
    - Time requirement $t_i$ which must be scheduled in a contiguous block
    - Target deadline $d_i$ by which time the request would like to be finished
    - Overall start time $s$
  - Requests are scheduled by the algorithm into time intervals $[s_i,f_i]$ such that $t_i = f_i - s_i$
  - Lateness of schedule for request $i$ is
    - If $d_i < f_i$ then request $i$ is late by $L_i = f_i - d_i$ otherwise its lateness $L_i = 0$
  - Maximum lateness $L = \max_i L_i$
- **Goal:** Find a schedule for all requests (values of $s_i$ and $f_i$ for each request $i$) to minimize the maximum lateness, $L$
### Scheduling to Minimizing Lateness

**Example:**

<table>
<thead>
<tr>
<th>$t_j$</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
<th>6</th>
</tr>
</thead>
<tbody>
<tr>
<td>$d_j$</td>
<td>3</td>
<td>2</td>
<td>1</td>
<td>4</td>
<td>3</td>
<td>2</td>
</tr>
</tbody>
</table>

**Lateness Calculation:**

- $d_1 = 6$ result in $t_1 = 6$, which is on time, so $lateness = 0$.
- $d_2 = 8$ result in $t_2 = 3$, so $lateness = 2$.
- $d_3 = 9$, $d_4 = 9$, $d_5 = 14$, and $d_6 = 15$ result in $max\,lateness = 6$.

**Graphical Representation:**

- The timeline shows the due dates and the processing times.
- The lateness is marked at the points where the job was completed after the due date.
- The arrow indicates the job completion time relative to its due date.
Minimizing Lateness: Greedy Algorithms

- Greedy template. Consider jobs in some order.
  
  - [Shortest processing time first] Consider jobs in ascending order of processing time \( t_j \).
  
  - [Earliest deadline first] Consider jobs in ascending order of deadline \( d_j \).
  
  - [Smallest slack] Consider jobs in ascending order of slack \( d_j - t_j \).
Minimizing Lateness: Greedy Algorithms

- Greedy template. Consider jobs in some order.

  - [Shortest processing time first] Consider jobs in ascending order of processing time $t_j$.

    |   |   |   |   |
    |---|---|---|---|
    | t | j | t | 1 | 1 |
    |   |   | d | 100 | 10 |

  - [Smallest slack] Consider jobs in ascending order of slack $d_j - t_j$.

    |   |   |   |   |
    |---|---|---|---|
    | t | j | t | 1 | 1 |
    |   |   | d | 2 | 10 |
Greedy Algorithm: Earliest Deadline First

- Order requests in increasing order of deadlines
- Schedule the request with the earliest deadline as soon as the resource becomes available
Minimizing Lateness: Greedy Algorithm

- Greedy algorithm. Earliest deadline first.

Sort deadlines in increasing order ($d_1 \leq d_2 \leq \ldots \leq d_n$)

\[
\begin{align*}
f & \leftarrow s \\
\text{for } i & \leftarrow 1 \text{ to } n \text{ to} \\
& \quad s_i \leftarrow f \\
& \quad f_i \leftarrow s_i + t_i \\
& \quad f \leftarrow f_i \\
\text{end for}
\end{align*}
\]

Max lateness = 1

<table>
<thead>
<tr>
<th>$d_1 = 6$</th>
<th>$d_2 = 8$</th>
<th>$d_3 = 9$</th>
<th>$d_4 = 9$</th>
<th>$d_5 = 14$</th>
<th>$d_6 = 15$</th>
</tr>
</thead>
<tbody>
<tr>
<td>0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
Proof for Greedy Algorithm: Exchange Argument

- We will show that if there is another schedule O (think optimal schedule) then we can gradually change O so that
  - at each step the maximum lateness in O never gets worse
  - it eventually becomes the same cost as A
Minimizing Lateness: No Idle Time

Observation. There exists an optimal schedule with no idle time.

Observation. The greedy schedule has no idle time.
Minimizing Lateness: Inversions

Definition. An inversion in schedule $S$ is a pair of jobs $i$ and $j$ such that $d_i < d_j$ but $j$ scheduled before $i$.

Observation. Greedy schedule has no inversions.

Observation. If a schedule (with no idle time) has an inversion, it has one with a pair of inverted jobs scheduled consecutively (by transitivity of $<$).
Minimizing Lateness: Inversions

- **Definition.** An inversion in schedule $S$ is a pair of jobs $i$ and $j$ such that $d_i < d_j$ but $j$ scheduled before $i$.

- **Claim.** Swapping two adjacent, inverted jobs reduces the number of inversions by one and does not increase the max lateness.
Minimizing Lateness: Inversions

- If \( d_j > d_i \) but \( j \) is scheduled in \( O \) immediately before \( i \) then swapping requests \( i \) and \( j \) to get schedule \( O' \) does not increase the maximum lateness
  - Lateness \( L_i' \leq L_i \) since \( i \) is scheduled earlier in \( O' \) than in \( O \)
  - Requests \( i \) and \( j \) together occupy the same total time slot in both schedules
    - All other requests \( k \neq i, j \) have \( L_k' = L_k \)
    - \( f_j' = f_i \) so \( L_j' = f_j' - d_j = f_i - d_j < f_j - d_j = L_j \)
    - Maximum lateness has not increased!
Claim: There is an optimal schedule with no idle time and no inversions

Proof:
- By previous argument there is an optimal schedule \( O \) with no idle time
- If \( O \) has an inversion then it has a consecutive pair of requests in its schedule that are inverted and can be swapped without increasing lateness
Optimal schedules and inversions

- Eventually these swaps will produce an optimal schedule with no inversions
  - Each swap decreases the number of inversions by 1
  - There are a bounded number of (at most $\frac{n(n-1)}{2}$) inversions (we only care that this is finite.)

QED
Idleness and Inversions are the only issue

■ **Claim:** All schedules with no inversions and no idle time have the same maximum lateness

■ **Proof**
  ■ Schedules can differ only in how they order requests with equal deadlines
  ■ Consider all requests having some common deadline \( d \)

  ■ Maximum lateness of these jobs is based only on the finish time of the last of these jobs but the set of these requests occupies the same time segment in both schedules
    ■ Last of these requests finishes at the same time in any such schedule.
Earliest Deadline First is optimal

- We know that
  - There is an optimal schedule with no idle time or inversions
  - All schedules with no idle time or inversions have the same maximum lateness
  - EDF produces a schedule with no idle time or inversions

- Therefore
  - EDF produces an optimal schedule
Optimal Caching/Paging

- Memory systems
  - many levels of storage with different access times
  - smaller storage has shorter access time
  - to access an item it must be brought to the lowest level of the memory system

- Consider the management problem between adjacent levels
  - Main memory with \( n \) data items from a set \( U \)
  - Cache can hold \( k < n \) items
  - Simplest version with no direct-mapping or other restrictions about where items can be
  - Suppose cache is full initially
    - Holds \( k \) data items to start with
Optimal Caching/Paging

- Given a memory request \( d \) from \( U \)
  - If \( d \) is stored in the cache we can access it quickly
  - If not then we call it a **cache miss** and (since the cache is full)
    - we must bring it into cache and **evict** some other data item from the cache
  - which one to evict?

- **Given** a sequence \( D = d_1, d_2, \ldots, d_m \) of elements from \( U \) corresponding to memory requests
- **Find** a sequence of evictions (an eviction schedule) that has as few cache misses as possible
Optimal Offline Caching

- **Goal.** Eviction schedule that minimizes number of cache misses (actually, # of evictions).

- **Example:** $k = 2$, initial cache = $ab$, requests: $a, b, c, b, c, a, a, b$.

- Optimal eviction schedule: 2 cache misses.
A Note on Optimal Caching

- In real operating conditions one typically needs an on-line algorithm
  - make the eviction decisions as each memory request arrives
- However to design and analyze these algorithms it is also important to understand how the best possible decisions can be made if one did know the future
  - Field of on-line algorithms compares the quality of on-line decisions to that of the optimal schedule
- What does an optimal schedule look like?
Belady’s Greedy Algorithm: Farthest-In-Future

- Given sequence $D = d_1, d_2, \ldots, d_m$

- When $d_i$ needs to be brought into the cache evict the item that is needed farthest in the future
  - Let $\text{NextAccess}_i(d) = \min \{ j \geq i : d_j = d \}$ be the next point in $D$ that item $d$ will be requested
  - Evict $d$ such that $\text{NextAccess}_i(d)$ is largest
Optimal Offline Caching: Farthest-In-Future

- **Farthest-in-future.** Evict item in the cache that is not requested until farthest in the future.

  current cache:  
  ![current cache: a b c d e f]

  future queries:  
  ![future queries: g a b c e d a b b a c d e a f a d e f g h ...]

  cache miss  
  eject this one

- **Theorem.** [Bellady, 1960s] FIF is an optimal eviction schedule.
- **Proof.** Algorithm and theorem are intuitive; proof is subtle.
Other Algorithms

- Often there is flexibility, e.g.
  - \( k=3, C=\{a,b,c\} \)
  - \( D= a \ b \ c \ d \ a \ d \ e \ a \ d \ b \ c \)
  - \( S_{\text{FIF}}= c \ b \ e \ d \)
  - \( S = b \ c \ d \ e \)

- Why aren’t other algorithms better?
  - Least-Frequently-Used-In-Future?

- Exchange Argument
  - We can swap choices to convert other schedules to Farthest-In-Future without losing quality
Reduced Eviction Schedules

- **Definition.** A reduced schedule is a schedule that only inserts an item into the cache in a step in which that item is requested.

- **Intuition.** Can transform an unreduced schedule into a reduced one with no more cache misses.

```
unreduced schedule
 a a b c
 a a x c
 c a d c
d a d b
 a a c b
 b a x b
c a c b
 a a b c
 a a b c

reduced schedule
 a a b c
 a a b c
 a a d c
d a d c
 a a d c
 b a d b
 c a c b
 a a c b
 a a c b
```
Claim. Given any unreduced schedule $S$, can transform it into a reduced schedule $S'$ with no more cache misses.

Proof. (by induction on number of unreduced items)
- Suppose $S$ brings $d$ into the cache at time $t$, without a request.
- Let $c$ be the item $S$ evicts when it brings $d$ into the cache.
- Case 1: $d$ evicted at time $t'$, before next request for $d$.
- Case 2: $d$ requested at time $t'$ before $d$ is evicted.

\[\begin{align*}
\text{Case 1} &: S & S' \\
& c & c \\
& t & t \\
& d & d \\
& t' & t' \\
& e & e \\
& d \text{ evicted at time } t', \text{ before next request}
\end{align*}\]

\[\begin{align*}
\text{Case 2} &: S & S' \\
& c & c \\
& t & t \\
& d & d \\
& t' & t' \\
& e & e \\
& d \text{ requested at time } t'
\end{align*}\]
Farthest-In-Future: Analysis

- **Theorem.** FIF is optimal eviction algorithm.
- **Proof.** (by induction on number or requests $j$)

   **Invariant:** There exists an optimal reduced schedule $S$ that makes the same eviction schedule as $S_{\text{FIF}}$ through the first $j$ requests.

Let $S$ be reduced schedule that satisfies invariant through $j$ requests. We produce $S'$ that satisfies invariant after $j+1$ requests.

- Consider $(j+1)^{\text{st}}$ request $d = d_{j+1}$.
- Since $S$ and $S_{\text{FIF}}$ have agreed up until now, they have the same cache contents before request $j+1$.
- **Case 1:** ($d$ is already in the cache). $S' = S$ satisfies invariant.
- **Case 2:** ($d$ is not in the cache and $S$ and $S_{\text{FIF}}$ evict the same element). $S' = S$ satisfies invariant.
Proof. (continued)

- **Case 3**: \(d\) is not in the cache; \(S_{\text{FIF}}\) evicts \(e\); \(S\) evicts \(f \neq e\).
  - begin construction of \(S'\) from \(S\) by evicting \(e\) instead of \(f\)

\[
\begin{array}{cccc|cccc}
  j & \text{same} & e & f & \text{same} & e & f \\
  \hline 
  S & & & & S' & & \\
  j+1 & \text{same} & e & d & \text{same} & d & f \\
  S & & & & S' & & 
\end{array}
\]

- now \(S'\) agrees with \(S_{\text{FIF}}\) on first \(j+1\) requests; we show that having element \(f\) in cache is no worse than having element \(e\)
**Farthest-In-Future: Analysis**

Let \( j' \) be the first time after \( j+1 \) that \( S \) and \( S' \) take a different action, and let \( g \) be item requested at time \( j' \).

<table>
<thead>
<tr>
<th>( j' )</th>
<th>( S )</th>
<th>( S' )</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>same</td>
<td>( f )</td>
</tr>
<tr>
<td></td>
<td>( e )</td>
<td></td>
</tr>
</tbody>
</table>

**Case 3a:** \( g = e \). Can't happen with Farthest-In-Future since there must be a request for \( f \) before \( e \).

**Case 3b:** \( g = f \). Element \( f \) can't be in cache of \( S \), so let \( e' \) be the element that \( S \) evicts.

- if \( e' = e \), \( S' \) accesses \( f \) from cache; now \( S \) and \( S' \) have same cache
- if \( e' \neq e \), \( S' \) evicts \( e' \) and brings \( e \) into the cache; now \( S \) and \( S' \) have the same cache

Note: \( S' \) is no longer reduced, but can be transformed into a reduced schedule that agrees with \( S_{FIF} \) through step \( j+1 \).
Farthest-In-Future: Analysis

- Let \( j' \) be the first time after \( j + 1 \) that \( S \) and \( S' \) take a different action, and let \( g \) be item requested at time \( j' \).

<table>
<thead>
<tr>
<th>( j' )</th>
<th>same</th>
<th>( e )</th>
<th>( S )</th>
<th>same</th>
<th>( f )</th>
<th>( S' )</th>
</tr>
</thead>
</table>

must involve \( e \) or \( f \) (or both)

- **Case 3c**: \( g \neq e, f \). \( S \) must evict \( e \).
  Make \( S' \) evict \( f \); now \( S \) and \( S' \) have the same cache. ▪

| \( j' \) | same | \( g \) | \( S \) | same | \( g \) | \( S' \) |
Caching Perspective

- Online vs. offline algorithms.
  - Offline: full sequence of requests is known a priori.
  - Online (reality): requests are not known in advance.
  - Caching is among most fundamental online problems in CS.

- LIFO. Evict page brought in most recently.
- LRU. Evict page whose most recent access was earliest.

- Theorem. FF is optimal offline eviction algorithm.
  - Provides basis for understanding and analyzing online algorithms.
  - LRU is k-competitive. [Section 13.8]
  - LIFO is arbitrarily bad.

[FF with direction of time reversed!]

Section 13.8
Single-source shortest paths

- Given an (un)directed graph $G=(V,E)$ with each edge $e$ having a non-negative weight $w(e)$ and a start vertex $s$

- Find length of shortest paths from $s$ to each vertex in $G$
A greedy algorithm

- **Dijkstra’s Algorithm:**
  - Maintain a set $S$ of vertices whose shortest paths are known
    - initially $S=\{s\}$
  - Maintaining current best lengths of paths that only go through $S$ to each of the vertices in $G$
    - path-lengths to elements of $S$ will be right, to $V-S$ they might not be right
  - Repeatedly add vertex $v$ to $S$ that has the shortest path-length of any vertex in $V-S$
    - update path lengths based on new paths through $v$
Dijkstra’s Algorithm

Dijkstra(G, w, s)

S ← {s}
d[s] ← 0

while S ≠ V do

of all edges e = (u, v) s.t. v ∉ S and u ∈ S select* one with the minimum value of d[u] + w(e)

S ← S ∪ {v}
d[v] ← d[u] + w(e)
pred[v] ← u

*For each v ∉ S maintain d'[v] = minimum value of d[u] + w(e) over all vertices u ∈ S s.t. e = (u, v) is in of G
Dijkstra’s Algorithm Correctness

Suppose all distances to vertices in $S$ are correct and $u$ has smallest current value in $V-S$

$\therefore$ distance value of vertex in $V-S = \text{length of shortest path from } s \text{ with only last edge leaving } S$

Suppose some other path to $v$ and $x$ = first vertex on this path not in $S$

$d'(v) \leq d'(x)$

$x-v$ path length $\geq 0$

$\therefore$ other path is longer

Therefore adding $v$ to $S$ keeps correct distances
Dijkstra’s Algorithm

- Algorithm also produces a tree of shortest paths to $v$ following $\text{pred}$ links
  - From $w$ follow its ancestors in the tree back to $v$

- If all you care about is the shortest path from $v$ to $w$ simply stop the algorithm when $w$ is added to $S$
Implementing Dijkstra’s Algorithm

- Need to
  - keep current distance values for nodes in $V-S$
  - find minimum current distance value
  - reduce distances when vertex moved to $S$
Data Structure Review

- **Priority Queue:**
  - Elements each with an associated **key**
  - Operations
    - **Insert**
    - **Find-min**
      - Return the element with the smallest key
    - **Delete-min**
      - Return the element with the smallest key and delete it from the data structure
    - **Decrease-key**
      - Decrease the key value of some element

- **Implementations**
  - Arrays: $O(n)$ time find/delete-min, $O(1)$ time insert/decrease-key
  - Heaps: $O(\log n)$ time insert/decrease-key/delete-min, $O(1)$ time find-min
Dijkstra’s Algorithm with Priority Queues

- For each vertex u not in tree maintain cost of current cheapest path through tree to u
  - Store u in priority queue with key = length of this path

- Operations:
  - n-1 insertions (each vertex added once)
  - n-1 delete-mins (each vertex deleted once)
    - pick the vertex of smallest key, remove it from the priority queue and add its edge to the graph
  - <m decrease-keys (each edge updates one vertex)
Dijskstra’s Algorithm with Priority Queues

- Priority queue implementations
  - Array
    - insert $O(1)$, delete-min $O(n)$, decrease-key $O(1)$
    - total $O(n+n^2+m)=O(n^2)$
  - Heap
    - insert, delete-min, decrease-key all $O(\log n)$
    - total $O(m \log n)$
  - $d$-Heap ($d=m/n$)
    - insert, decrease-key $O(\log_{m/n} n)$
    - delete-min $O((m/n) \log_{m/n} n)$
    - total $O(m \log_{m/n} n)$
Minimum Spanning Trees (Forests)

- Given an undirected graph $G=(V,E)$ with each edge $e$ having a weight $w(e)$

- Find a subgraph $T$ of $G$ of minimum total weight s.t. every pair of vertices connected in $G$ are also connected in $T$
  - if $G$ is connected then $T$ is a tree otherwise it is a forest
Weighted Undirected Graph
Greedy Algorithm

- **Prim’s Algorithm:**
  - start at a vertex $s$
  - add the cheapest edge adjacent to $s$
  - repeatedly add the cheapest edge that joins the vertices explored so far to the rest of the graph
  - Exactly like Dijsktra’s Algorithm but with a different metric
**Dijsktra’s Algorithm**

Dijkstra(G,w,s)

\[ S \leftarrow \{ s \} \]
\[ d[s] \leftarrow 0 \]

while \( S \neq V \) do

of all edges \( e=(u,v) \) s.t. \( v \notin S \) and \( u \in S \) select* one with the minimum value of \( d[u]+w(e) \)

\[ S \leftarrow S \cup \{ v \} \]
\[ d[v] \leftarrow d[u]+w(e) \]
\[ \text{pred}[v] \leftarrow u \]

*For each \( v \notin S \) maintain \( d'[v]=\text{minimum value of } d[u]+w(e) \) over all vertices \( u \in S \) s.t. \( e=(u,v) \) is in of \( G \)
Prim’s Algorithm

Prim(G, w, s)

\[ S \leftarrow \{s\} \]

while \( S \neq V \) do

of all edges \( e=(u, v) \) s.t. \( v \not\in S \) and \( u \in S \) select* one with the minimum value of \( w(e) \)

\[ S \leftarrow S \cup \{v\} \]

\[ \text{pred}[v] \leftarrow u \]

*For each \( v \not\in S \) maintain \( \text{small}[v] = \text{minimum value of } w(e) \) over all vertices \( u \in S \) s.t. \( e=(u, v) \) is in of \( G \)
Second Greedy Algorithm

- Kruskal’s Algorithm
  - Start with the vertices and no edges
  - Repeatedly add the cheapest edge that joins two different components. i.e. that doesn’t create a cycle
Why greed is good

- **Definition:** Given a graph $G=(V,E)$, a **cut** of $G$ is a partition of $V$ into two non-empty pieces, $S$ and $V-S$.

- **Lemma:** For every cut $(S,V-S)$ of $G$, there is a minimum spanning tree (or forest) containing any **cheapest edge crossing the cut**, i.e. connecting some node in $S$ with some node in $V-S$.
  - call such an edge **safe**
Cuts and Spanning Trees
The greedy algorithms always choose safe edges

- Prim’s Algorithm
  - Always chooses cheapest edge from current tree to rest of the graph
  - This is cheapest edge across a cut which has the vertices of that tree on one side.
Prim’s Algorithm
The greedy algorithms always choose safe edges

- Kruskal’s Algorithm
  - Always chooses cheapest edge connecting two pieces of the graph that aren’t yet connected
  - This is the cheapest edge across any cut which has those two pieces on different sides and doesn’t split any current pieces.
Kruskal’s Algorithm
Proof of Lemma: An Exchange Argument

Suppose you have an MST not using cheapest edge $e$

Endpoints of $e$, $u$ and $v$ must be connected in $T$
Proof of Lemma

Suppose you have an MST $T$ not using cheapest edge $e$

Endpoints of $e$, $u$ and $v$ must be connected in $T$
Proof of Lemma

Suppose you have an MST $T$ not using cheapest edge $e$

Endpoints of $e$, $u$ and $v$ must be connected in $T$
Proof of Lemma

Suppose you have an MST not using cheapest edge $e$

Endpoints of $e$, $u$ and $v$ must be connected in $T$

$w(e) \leq w(h)$
Proof of Lemma

Replacing $h$ by $e$ does not increase weight of $T$

\[ w(e) \leq w(h) \]

All the same points are connected by the new tree
Kruskal’s Algorithm Implementation & Analysis

- First sort the edges by weight \( O(m \log m) \)
- Go through edges from smallest to largest
  - if endpoints of edge \( e \) are currently in different components
    - then add to the graph
  - else skip
- Union-find data structure handles last part
- Total cost of last part: \( O(m \alpha(n)) \) where \( \alpha(n) \ll \log m \)
- Overall \( O(m \log n) \)
Union-find disjoint sets data structure

- Maintaining components
  - start with $n$ different components
    - one per vertex
  - find components of the two endpoints of $e$
    - $2m$ finds
  - union two components when edge connecting them is added
    - $n-1$ unions
Prim’s Algorithm with Priority Queues

- For each vertex \( u \) not in tree maintain current cheapest edge from tree to \( u \)
  - Store \( u \) in priority queue with key = weight of this edge

- Operations:
  - \( n-1 \) insertions (each vertex added once)
  - \( n-1 \) delete-mins (each vertex deleted once)
    - pick the vertex of smallest key, remove it from the p.q. and add its edge to the graph
  - \(<m\) decrease-keys (each edge updates one vertex)
Prim’s Algorithm with Priority Queues

- Priority queue implementations
  - Array
    - insert $O(1)$, delete-min $O(n)$, decrease-key $O(1)$
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Boruvka’s Algorithm (1927)

- A bit like Kruskal’s Algorithm
  - Start with $n$ components consisting of a single vertex each
  - At each step, each component chooses its cheapest outgoing edge to add to the spanning forest
    - Two components may choose to add the same edge
  - Useful for parallel algorithms since components may be processed (almost) independently
Many other minimum spanning tree algorithms, most of them greedy

- Cheriton & Tarjan
  - $O(m \log \log n)$ time using a queue of components
- Chazelle
  - $O(m \alpha(m))$ time
    - Incredibly hairy algorithm
- Pettie & Ramachandran
  - “Optimal” deterministic algorithm
  - Actual runtime unknown, best analysis is $O(m \alpha(m))$ time
- Karger, Klein & Tarjan
  - $O(m+n)$ time randomized algorithm that works most of the time
More on Cheriton-Tarjan algorithm

- Cheriton & Tarjan
  - \(O(m \log \log n)\) time using a queue of components
    - At each step we add the cheapest edge leaving the component at the head of the queue and add the new component to the end of the queue
    - On the \(j^{th}\) pass through the queue each component has at least \(2^j\) vertices
  - Each component has an associated heap of candidate edges
  - When components join, their heaps are melded