

# CSE 517

# Natural Language Processing

# Winter 2017

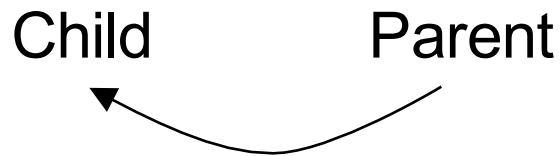
## Dependency Parsing

## And Other Grammar Formalisms

Yejin Choi - University of Washington

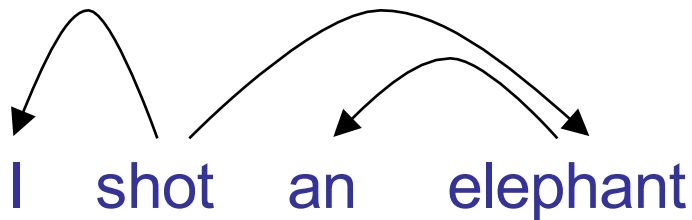
# Dependency Grammar

For each word, find one parent.

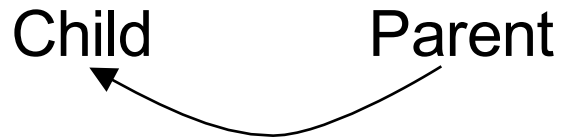


A child is **dependent** on the parent.

- A child is an **argument** of the parent.
- A child **modifies** the parent.

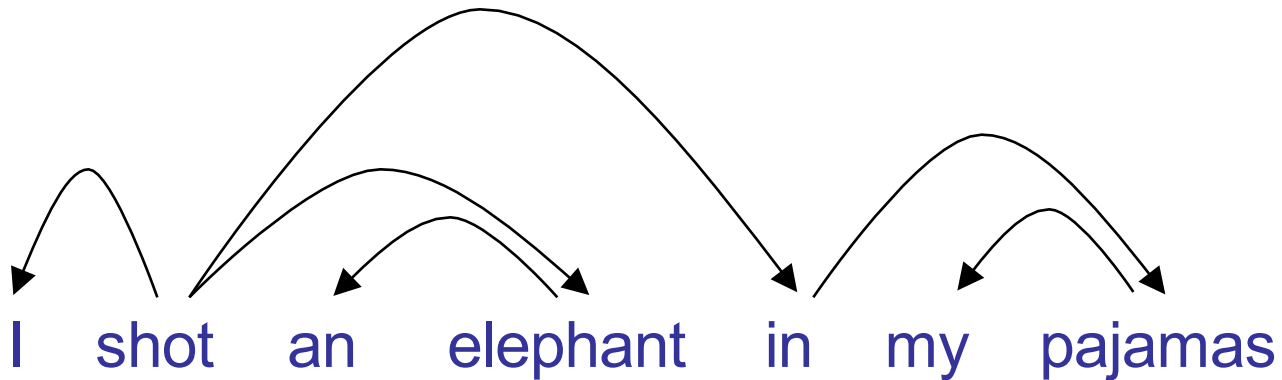


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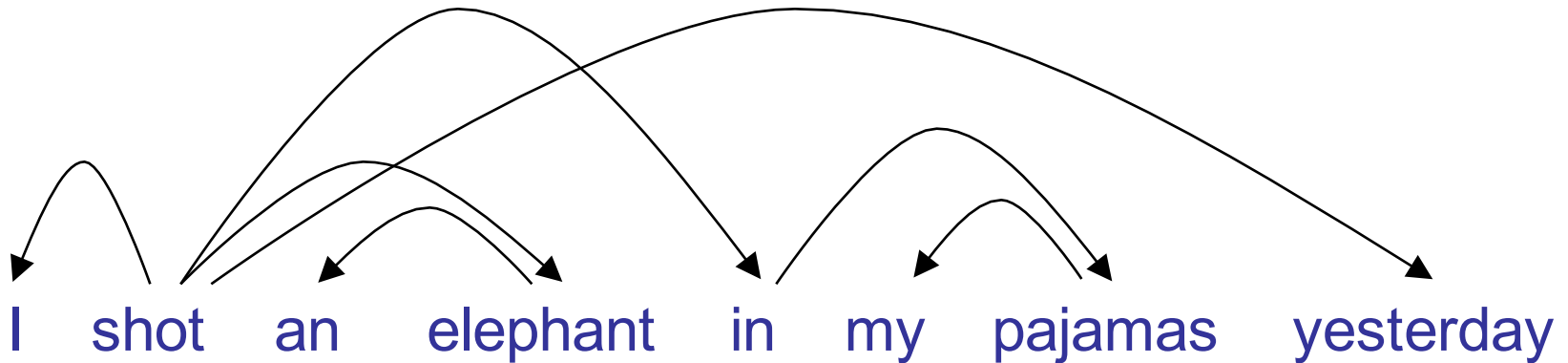


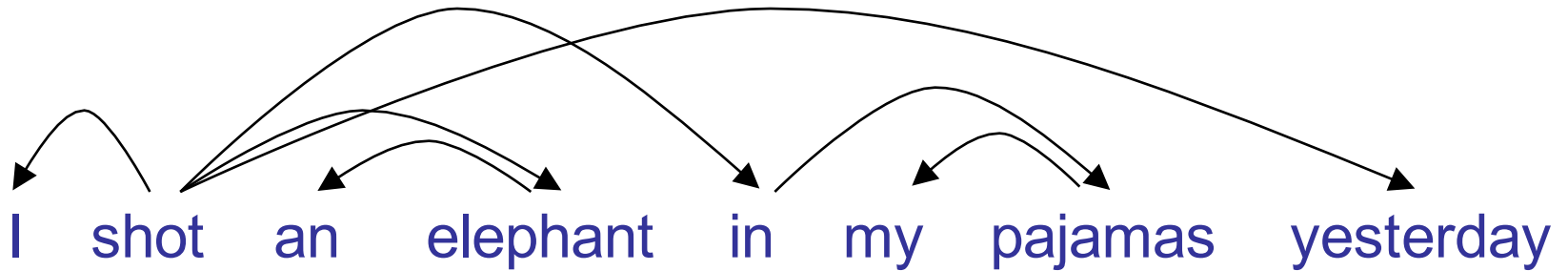
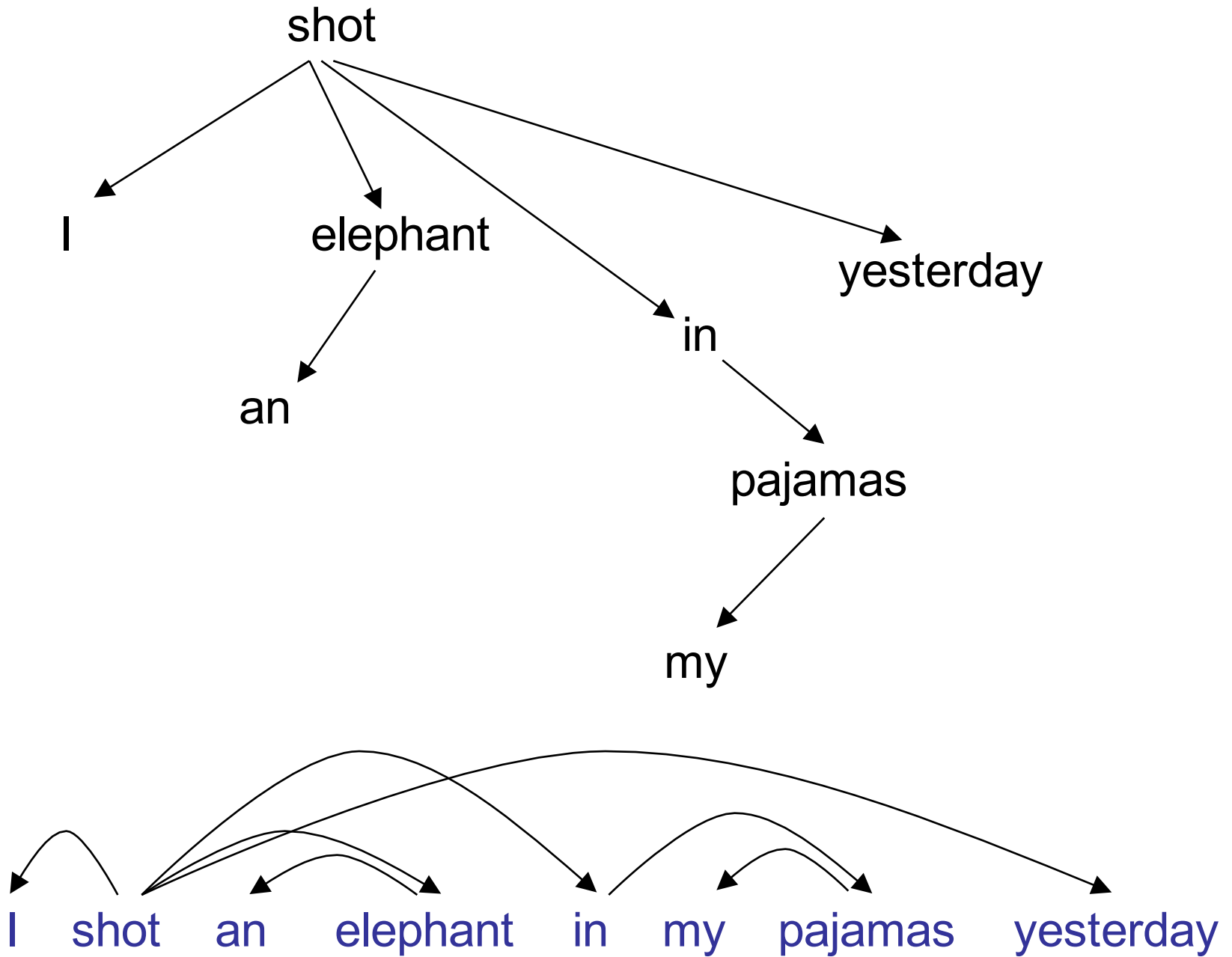
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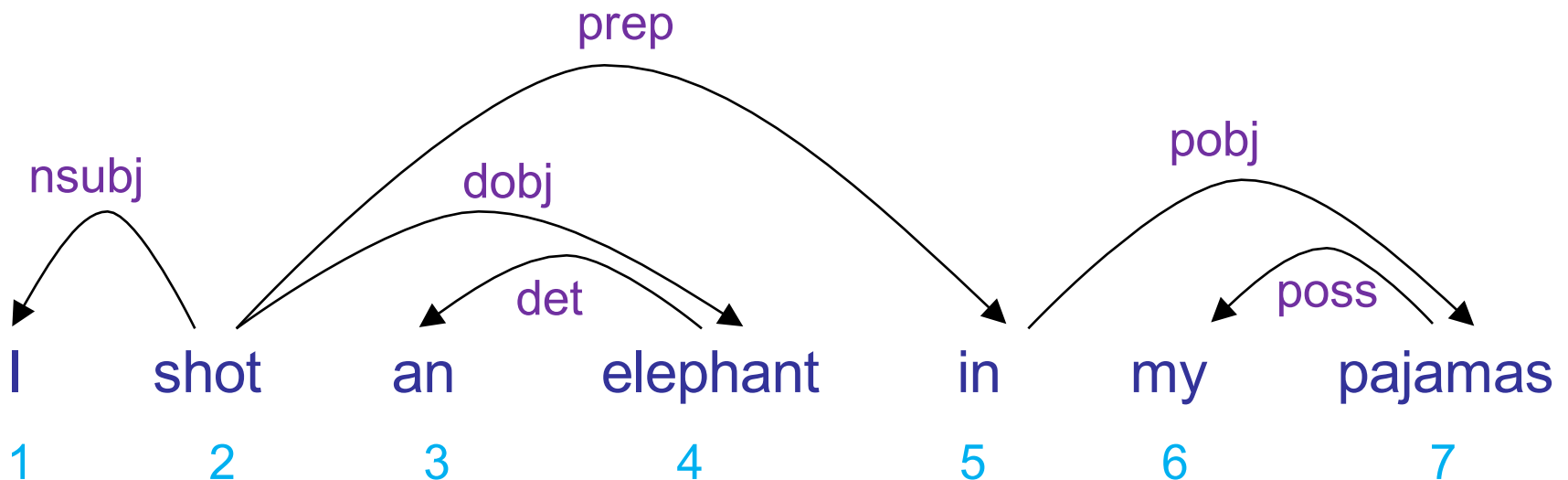




# Typed Dependencies

nsubj(shot-2, i-1)  
root(ROOT-0, shot-2)  
det(elephant-4, an-3)  
dobj(shot-2, elephant-4)

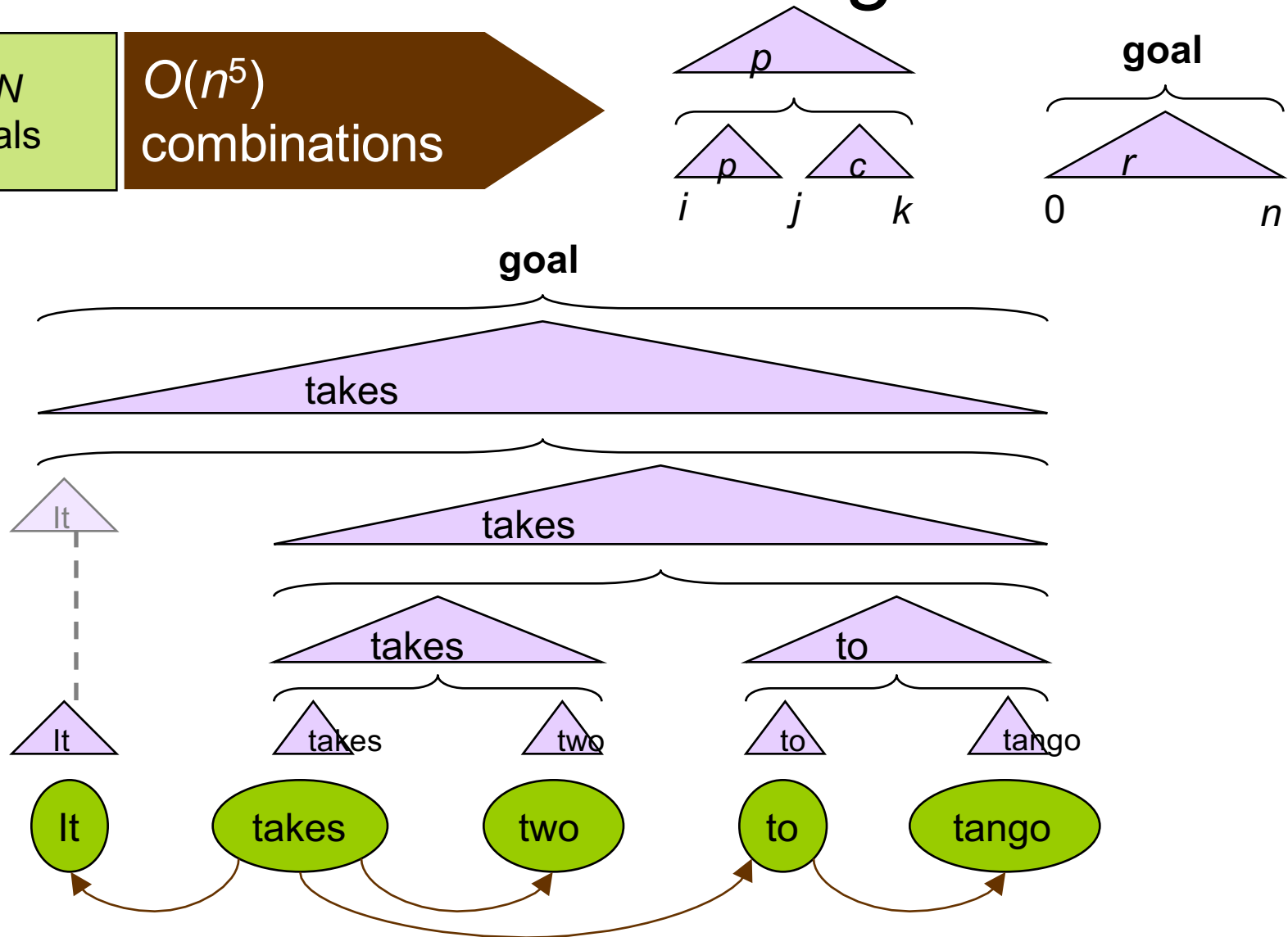
prep(shot-2, in-5)  
poss(pajamas-7, my-6)  
pobj(in-5, pajamas-7)



# Naïve CKY Parsing

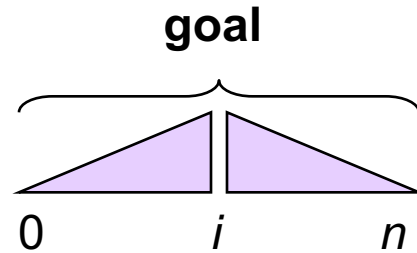
$O(n^5 N^3)$  if  $N$   
nonterminals

$O(n^5)$   
combinations

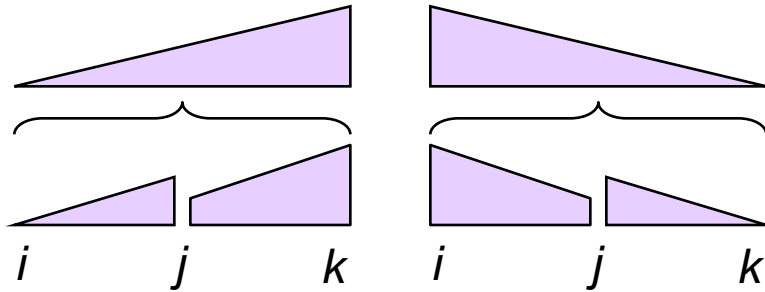


# Eisner Algorithm (Eisner & Satta, 1999)

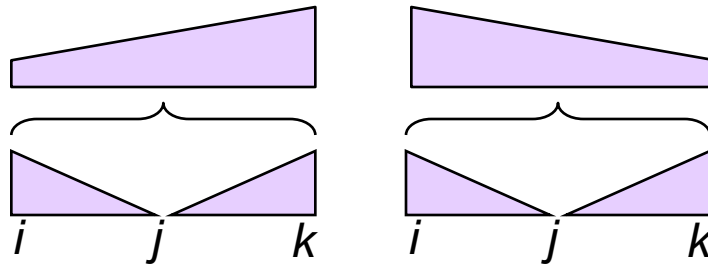
*This happens only once as the very final step*



*Without adding a dependency arc*



*When adding a dependency arc (head is higher)*

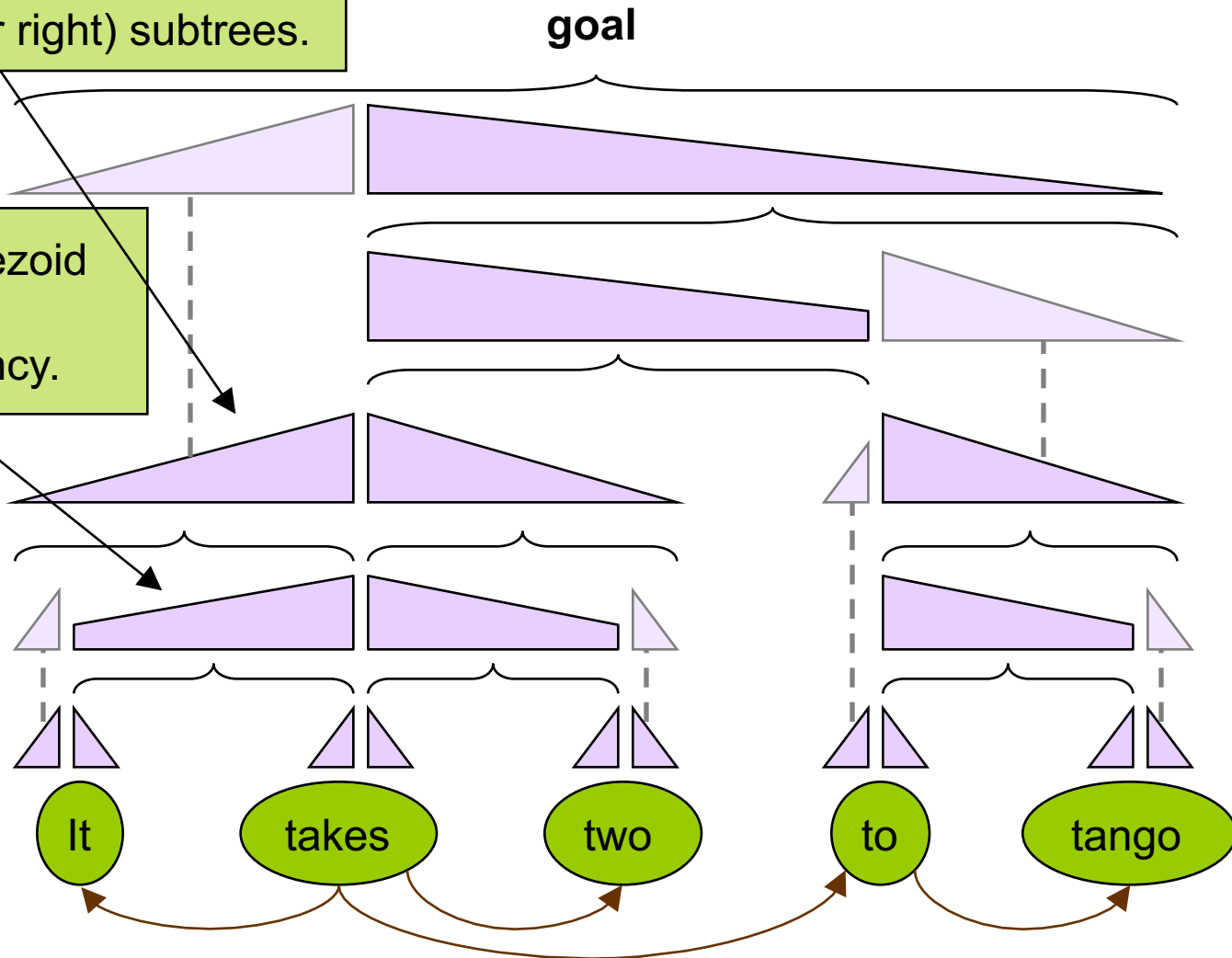




# Eisner Algorithm (Eisner & Satta, 1999)

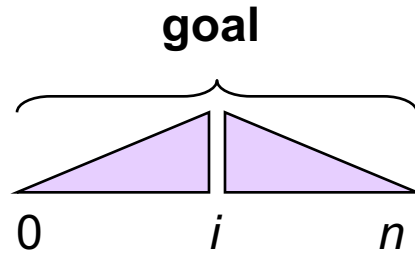
A triangle is a head with some left (or right) subtrees.

One trapezoid per dependency.

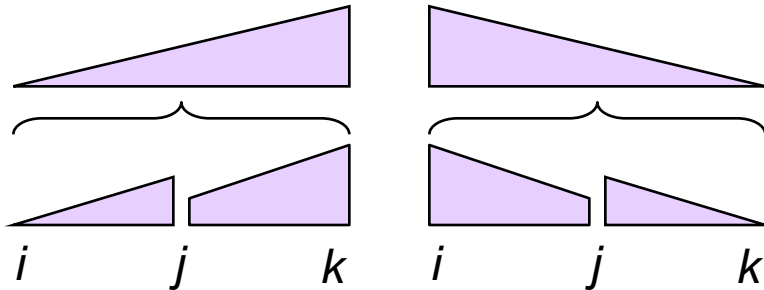


# Eisner Algorithm (Eisner & Satta, 1999)

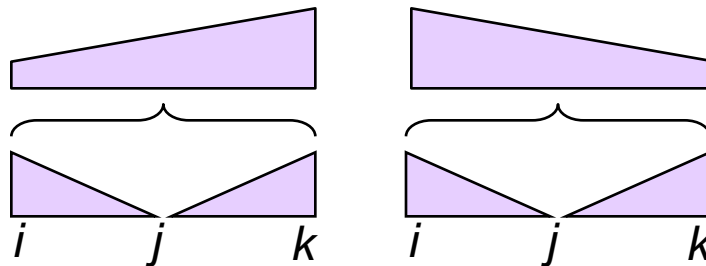
$O(n)$   
combinations



$O(n^3)$   
combinations



$O(n^3)$   
combinations



Gives  $O(n^3)$  dependency grammar  
parsing

# Eisner Algorithm

- Base case:

$$\forall t \in \{\underline{\triangleleft}, \underline{\triangleright}, \triangleleft, \triangleright\}, \pi(i, i, t) = 0$$

- Recursion:

$$\pi(i, j, \underline{\triangleleft}) = \max_{i \leq k < j} \left( \pi(i, k, \triangleright) + \pi(k + 1, j, \triangleleft) + \phi(w_j, w_i) \right)$$

$$\pi(i, j, \underline{\triangleright}) = \max_{i \leq k < j} \left( \pi(i, k, \triangleright) + \pi(k + 1, j, \triangleleft) + \phi(w_i, w_j) \right)$$

$$\pi(i, j, \triangleleft) = \max_{i \leq k < j} \left( \pi(i, k, \triangleleft) + \pi(k + 1, j, \underline{\triangleleft}) \right)$$

$$\pi(i, j, \triangleright) = \max_{i \leq k < j} \left( \pi(i, k, \underline{\triangleright}) + \pi(k + 1, j, \triangleright) \right)$$

- Final case:

$$\pi(1, n, \triangleleft \triangleright) = \max_{1 \leq k < n} \left( \pi(1, k, \triangleleft) + \pi(k + 1, n, \triangleright) \right)$$

# CFG vs Dependency Parse I

- CFG focuses on “**constituency**” (i.e., phrasal/clausal structure)
- Dependency focuses on “**head**” relations.
  
- CFG includes non-terminals. CFG edges are not typed.
- No non-terminals for dependency trees. Instead, dependency trees provide “dependency types” on edges.
  
- Dependency types encode “**grammatical roles**” like
  - nsubj -- nominal subject
  - dobj – direct object
  - pobj – prepositional object
  - nsubjpass – nominal subject in a passive voice

# CFG vs Dependency Parse II

- Can we get “heads” from CFG trees?
  - Yes. In fact, modern statistical parsers based on CFGs use hand-written “head rules” to assign “heads” to all nodes.
- Can we get constituents from dependency trees?
  - Yes, with some efforts.
- Can we transform CFG trees to dependency parse trees?
  - Yes, and transformation software exists. (stanford toolkit based on [de Marneffe et al. LREC 2006])
- Can we transform dependency trees to CFG trees?
  - Mostly yes, but (1) dependency parse can capture **non-projective** dependencies, while CFG cannot, and (2) people rarely do this in practice

# CFG vs Dependency Parse III

- Both are context-free.
- Both are used frequently today, but dependency parsers are more recently popular.
- CKY Parsing algorithm:
  - $O(N^3)$  using CKY & unlexicalized grammar
  - $O(N^5)$  using CKY & lexicalized grammar ( $O(N^4)$  also possible)
- Dependency parsing algorithm:
  - $O(N^5)$  using naïve CKY
  - $O(N^3)$  using Eisner algorithm
  - $O(N^2)$  based on minimum directed spanning tree algorithm (arborescence algorithm, aka, Edmond-Chu-Liu algorithm – see edmond.pdf)
- Linear-time  $O(N)$  Incremental parsing (shift-reduce parsing) possible for both grammar formalisms

# Non Projective Dependencies

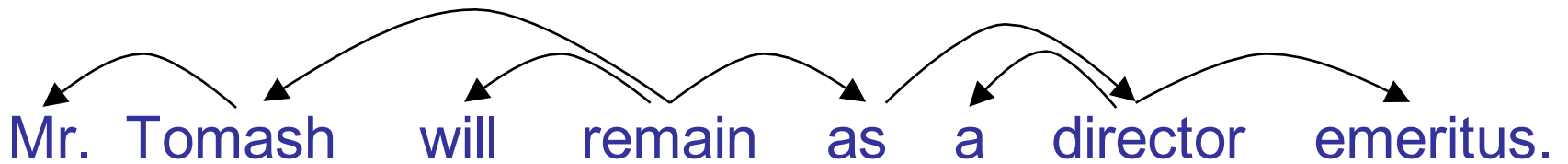
- Mr. Tomash will remain as a director emeritus.
- A hearing is scheduled on the issue today.

# Non Projective Dependencies

- Projective dependencies: when the tree edges are drawn directly on a sentence, it forms a tree (without a cycle), and there is no crossing edge.

- Projective Dependency:

- Eg:

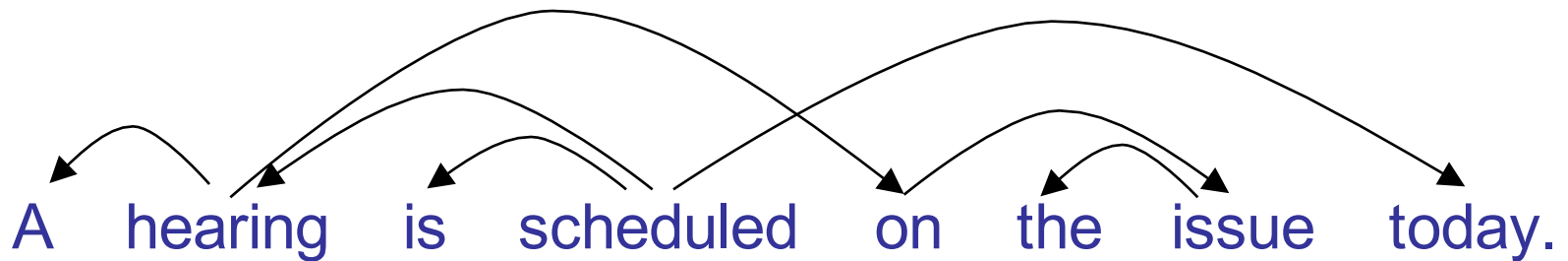




# Non Projective Dependencies

- Projective dependencies: when the tree edges are drawn directly on a sentence, it forms a tree (without a cycle), and there is no crossing edge.
- Non-projective dependency:

■ Eg:



# Non Projective Dependencies

- which word does “on the issue” modify?
  - We scheduled a meeting on the issue today.
  - A meeting is scheduled on the issue today.
- CFGs capture only projective dependencies (why?)

# Coordination across Constituents

- Right-node raising:

- [[She **bought**] and [he ate]] bananas.

- Argument-cluster coordination:

- I **give** [[you an apple] and [him a pear]].

- Gapping:

- She **likes** sushi, and he sashimi

→ CFGs don't capture coordination across constituents:

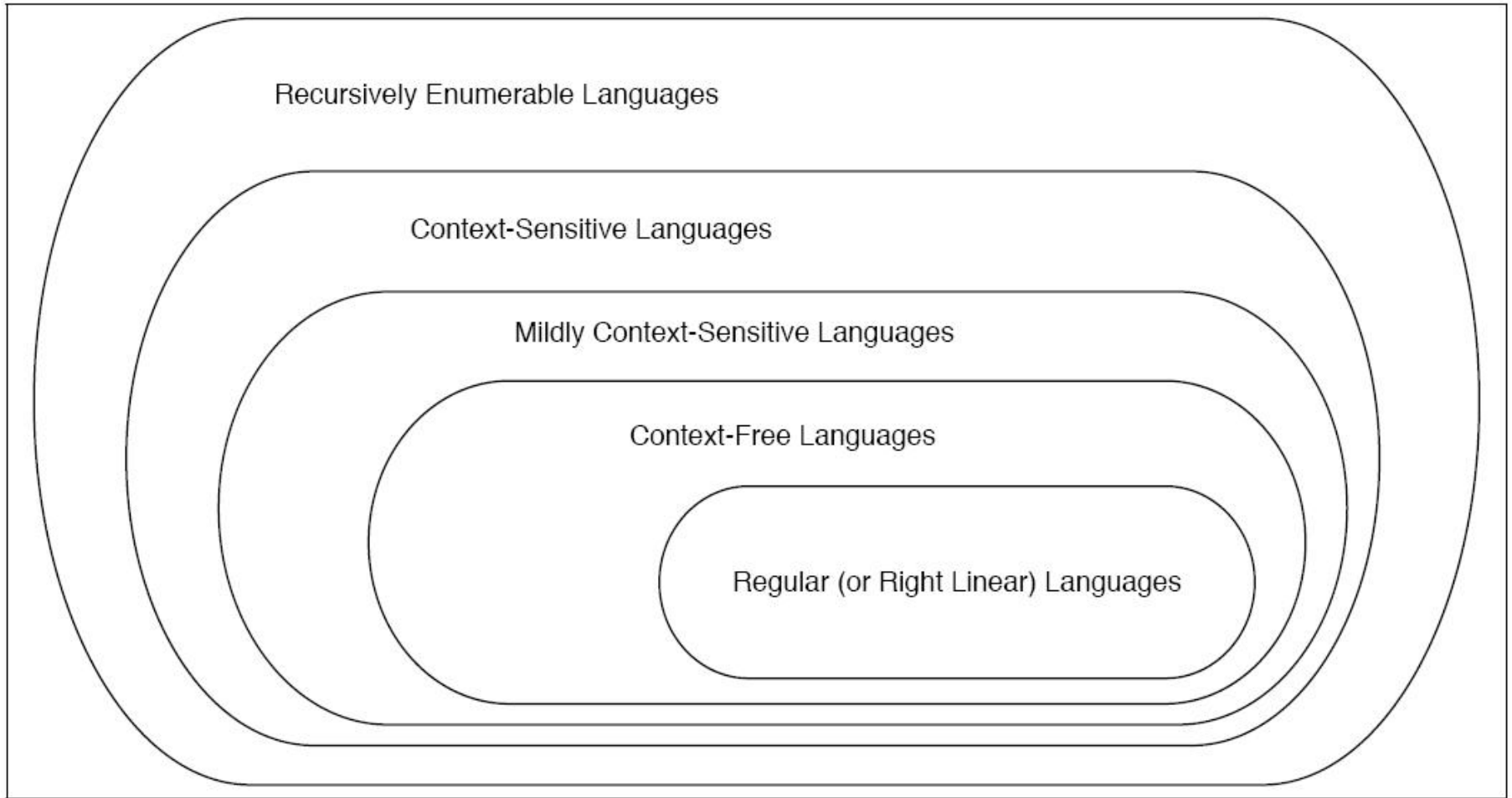
# Coordination across Constituents

- She bought and he ate bananas.
- I give you an apple and him a pear.

Compare above to:

- She bought and ate bananas.
- She bought bananas and apples.
- She bought bananas and he ate apples.

# The Chomsky Hierarchy



# The Chomsky Hierarchy

| Type | Common Name              | Rule Skeleton  | Linguistic Example        |
|------|--------------------------|--|---------------------------|
| 0    | Turing Equivalent        | $\alpha \rightarrow \beta$ , s.t. $\alpha \neq \epsilon$                       | HPSG, LFG, Minimalism     |
| 1    | Context Sensitive        | $\alpha A \beta \rightarrow \alpha \gamma \beta$ , s.t. $\gamma \neq \epsilon$ |                           |
| –    | Mildly Context Sensitive |  | TAG, CCG                  |
| 2    | Context Free             | $A \rightarrow \gamma$   | Phrase-Structure Grammars |
| 3    | Regular                  | $A \rightarrow xB$ or $A \rightarrow x$  | Finite-State Automata     |

- Head-Driven Phrase Structure Grammar (HPSG) (Pollard and Sag, 1987, 1994)
- Lexical Functional Grammar (LFG) (Bresnan, 1982)
- Minimalist Grammar (Stabler, 1997)
  
- Tree-Adjoining Grammars (TAG) (Joshi, 1969)
- Combinatory Categorical Grammars (CCG) (Steedman, 1986)

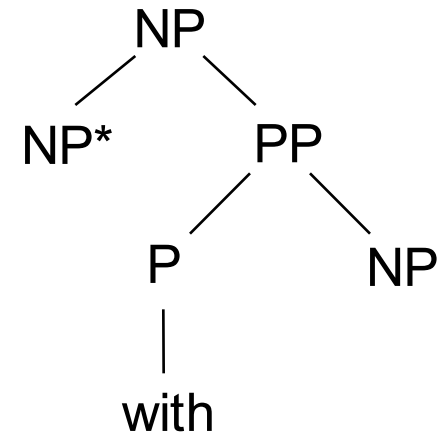
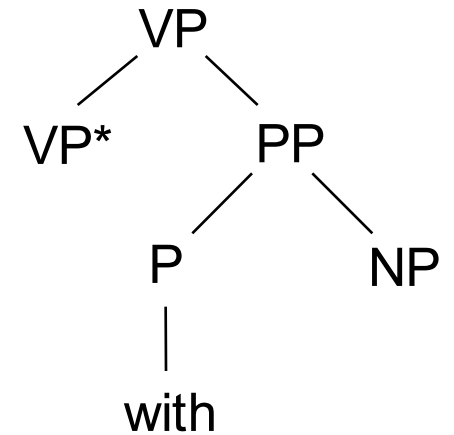
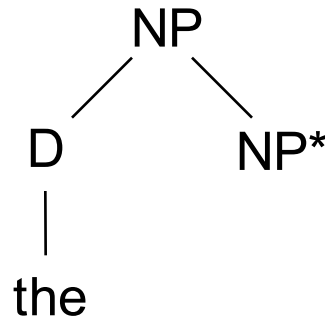
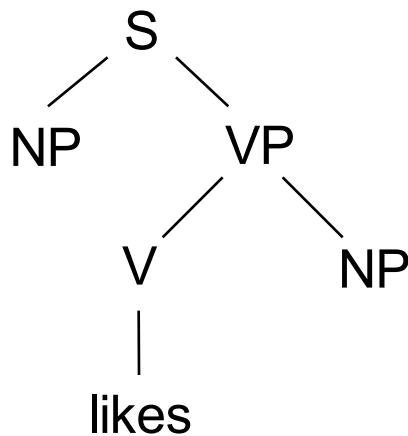
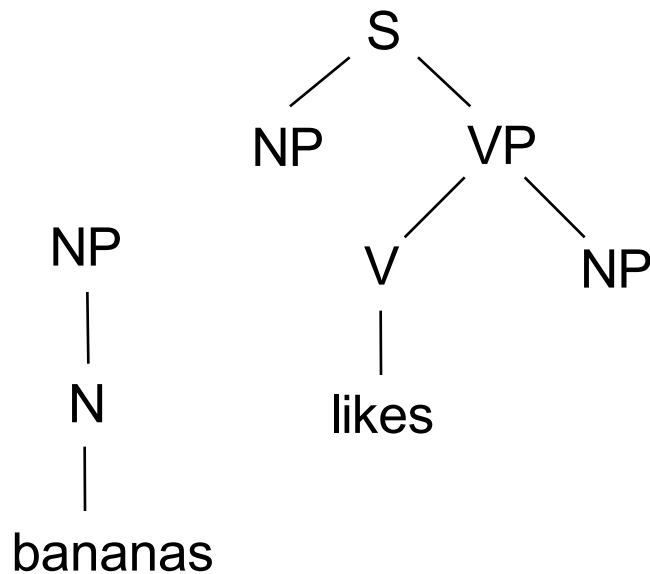
# Mildly Context-Sensitive Grammar Formalisms

# I. Tree Adjoining Grammar (TAG)

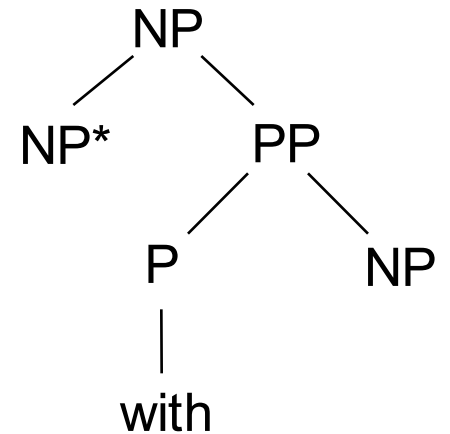
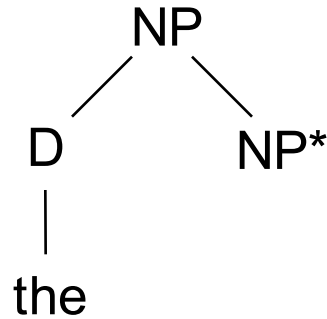
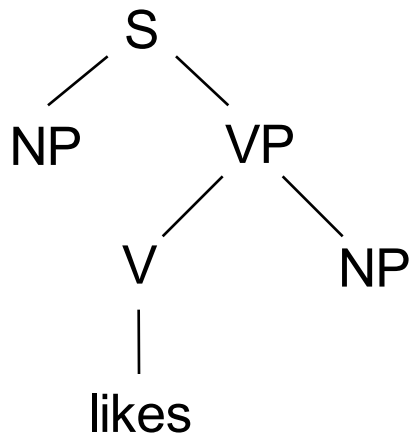
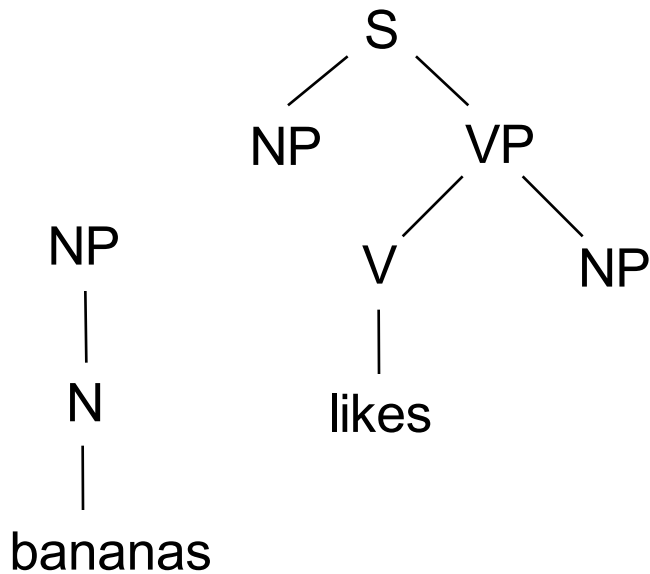
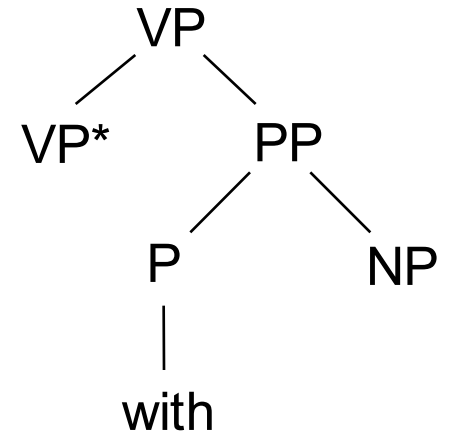
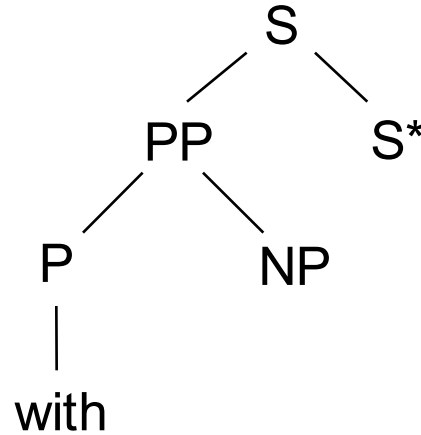
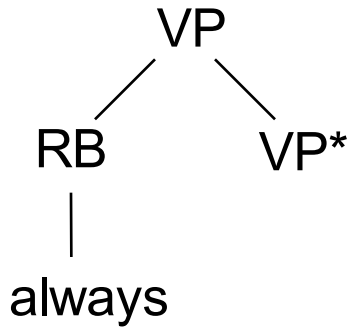


# TAG Lexicon (Supertags)

- Tree-Adjoining Grammars (TAG) (Joshi, 1969)
- “... *super parts of speech (supertags): almost parsing*” (Joshi and Srinivas 1994)
- POS tags enriched with syntactic structure
- also used in other grammar formalisms (e.g., CCG)

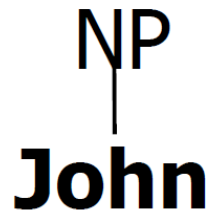


# TAG Lexicon (Supertags)

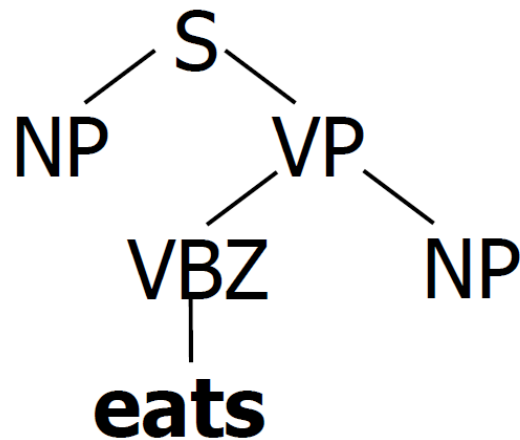


# Example: TAG Lexicon

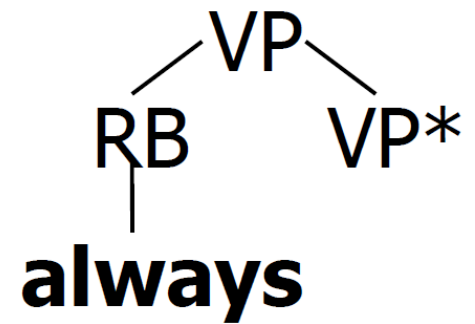
$\alpha_2$ :



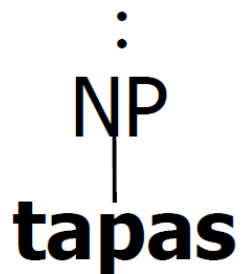
$\alpha_1$ :



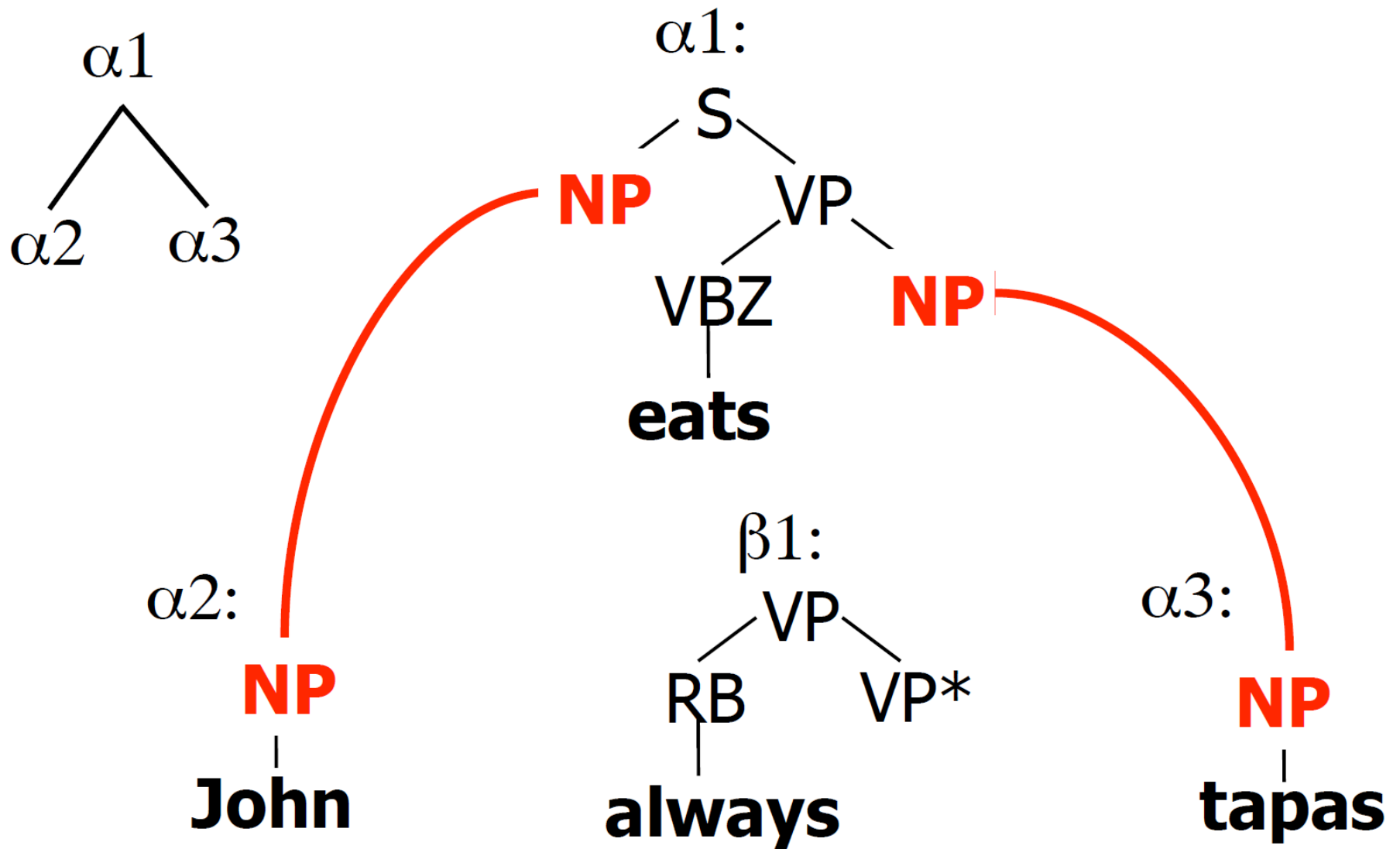
$\beta_1$ :



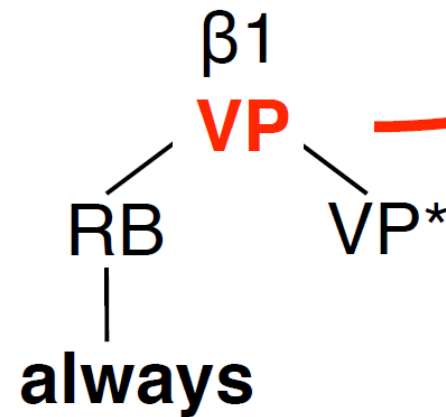
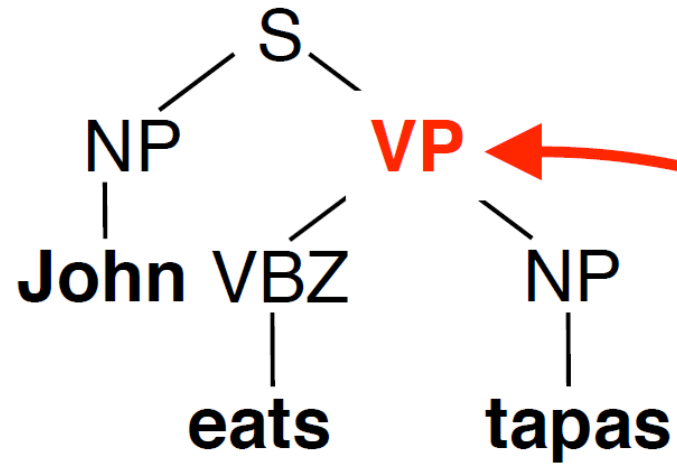
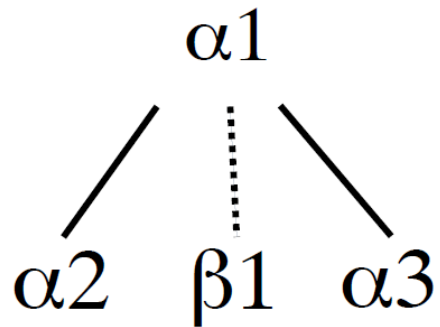
$\alpha_3$



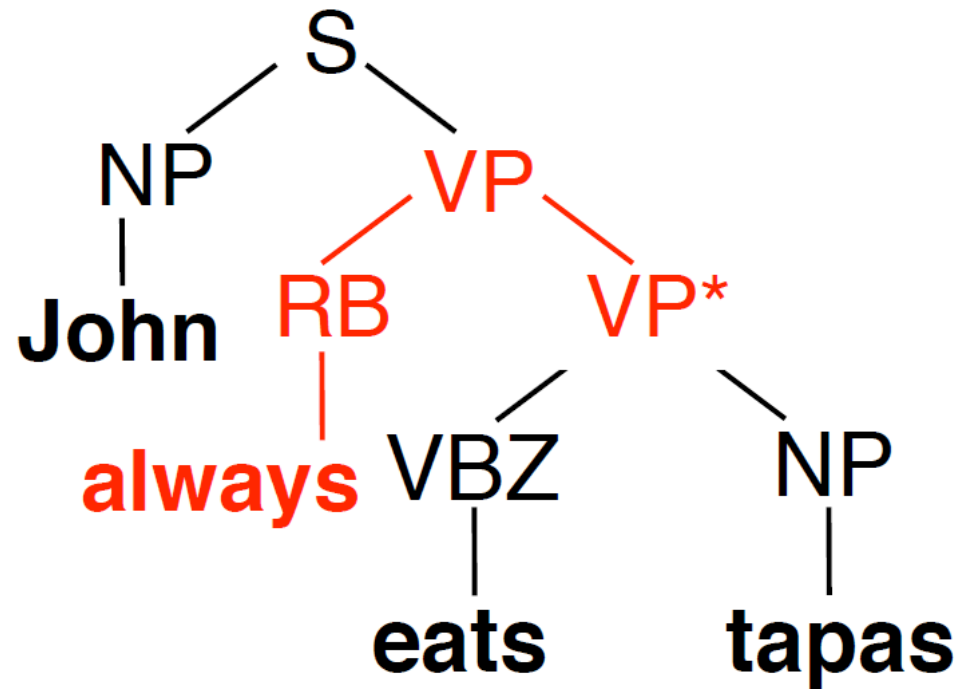
# Example: TAG Derivation



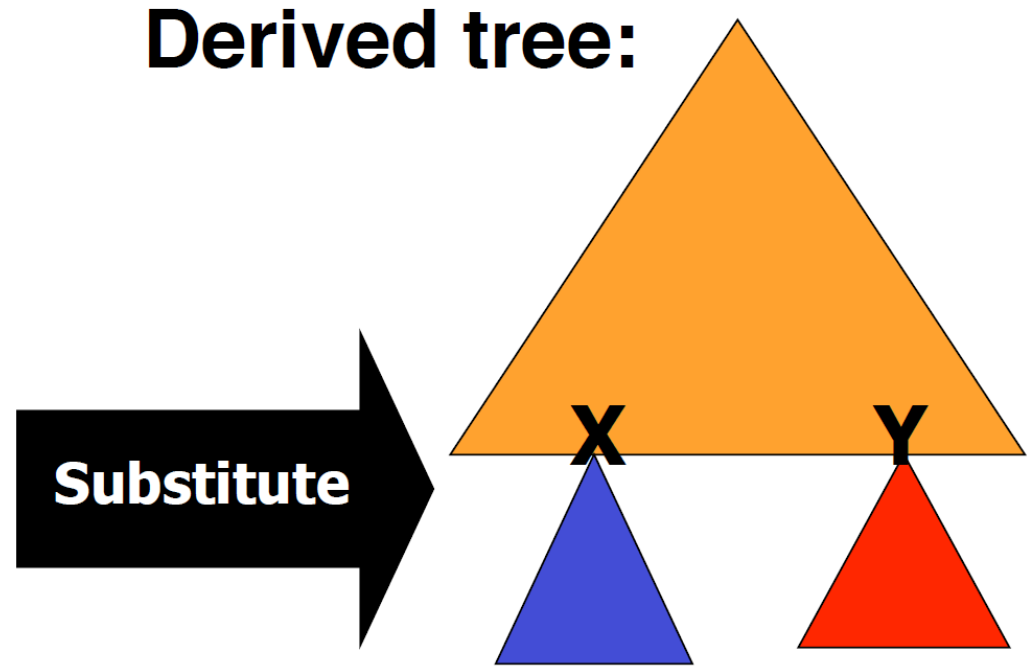
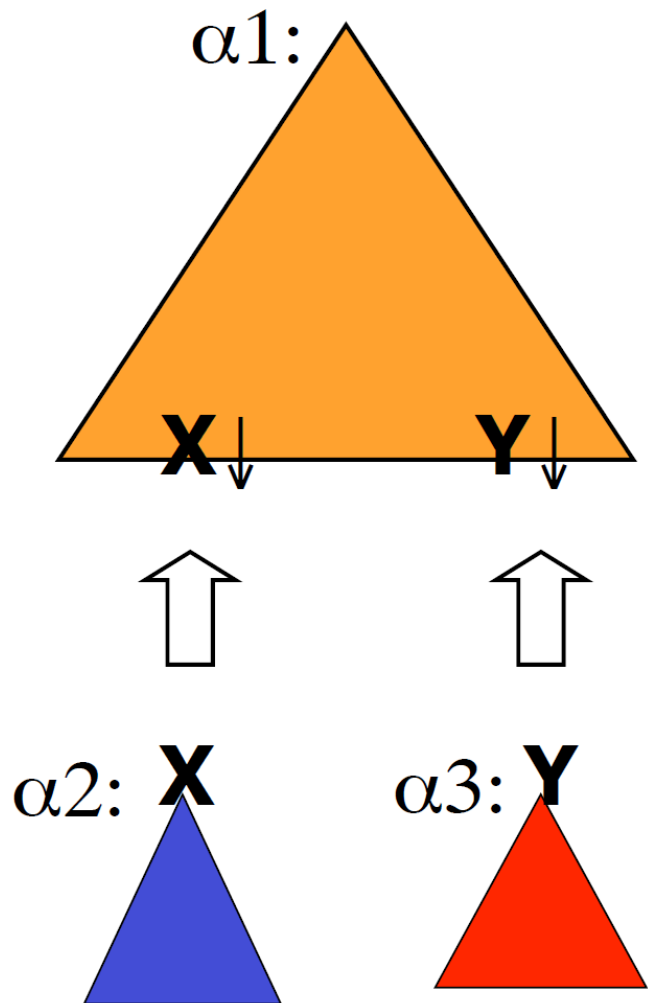
# Example: TAG Derivation



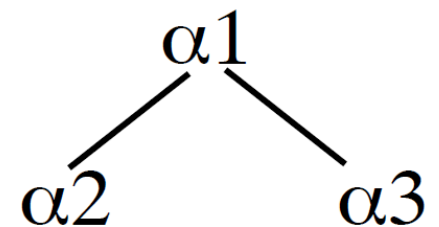
# Example: TAG Derivation



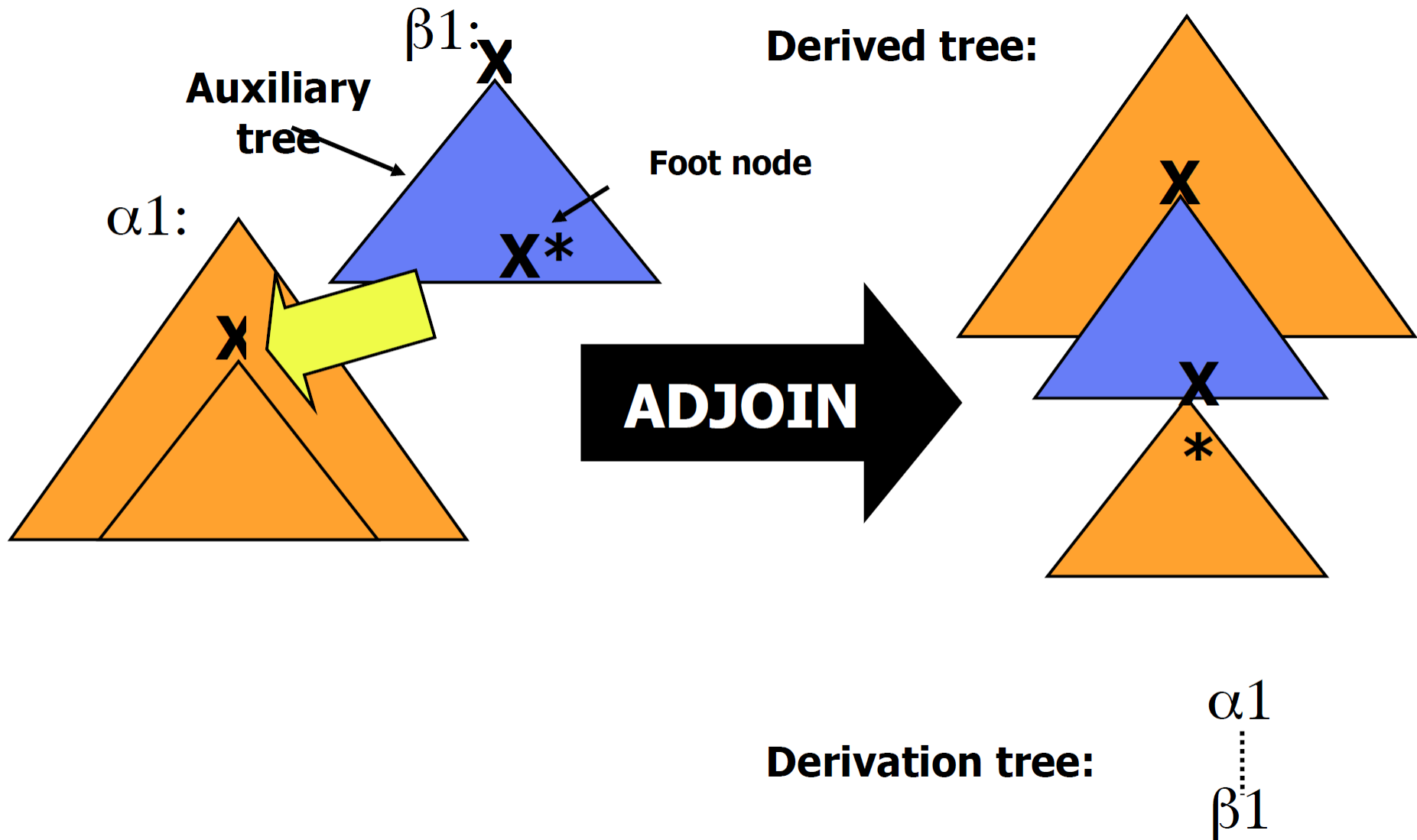
# TAG rule 1: Substitution



Derivation tree:

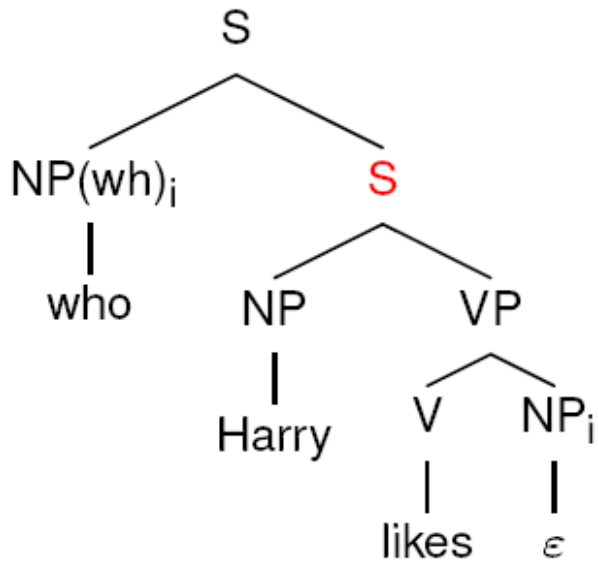
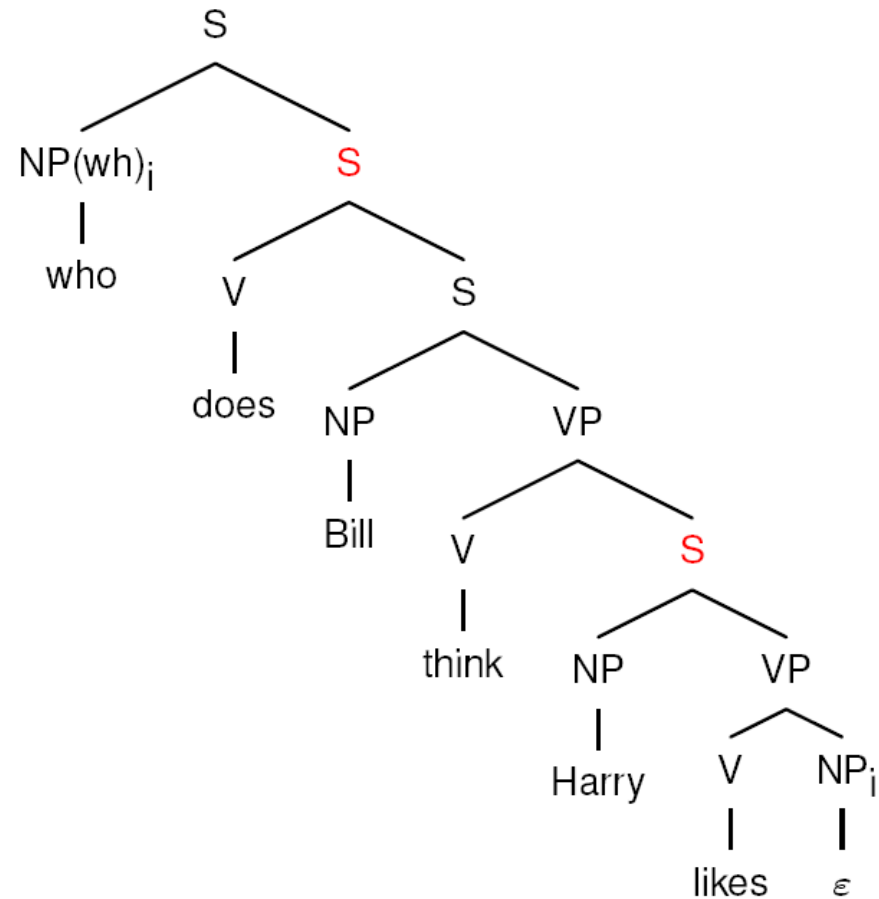
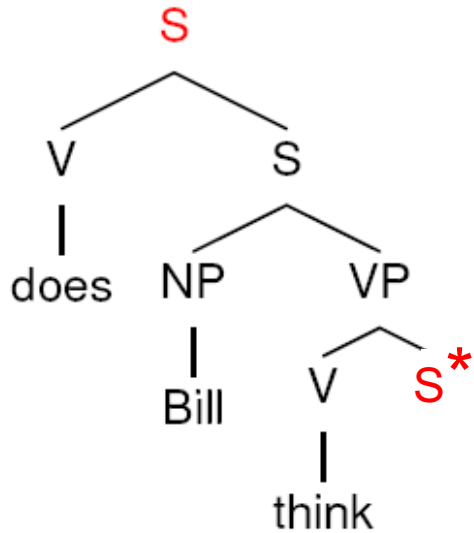


# TAG rule 2: Adjunction



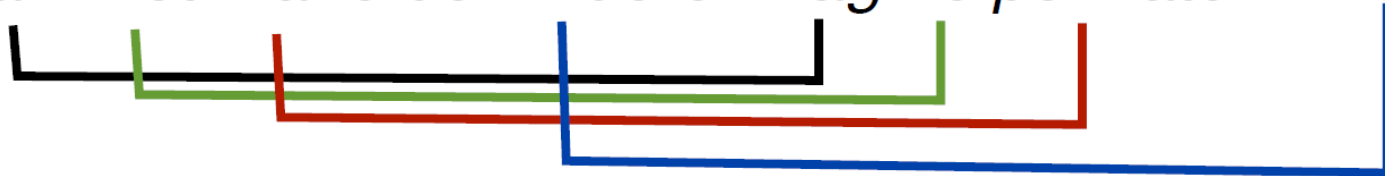


# (1) Can handle long distance dependencies



## (2) Cross-serial Dependencies

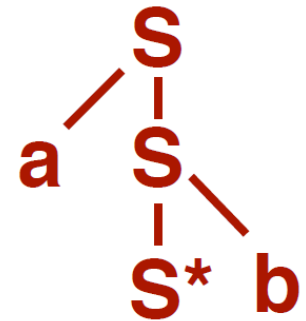
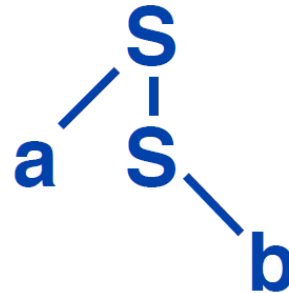
*dat Jan Piet Marie de kinderen zag helpen laten zwemmen*



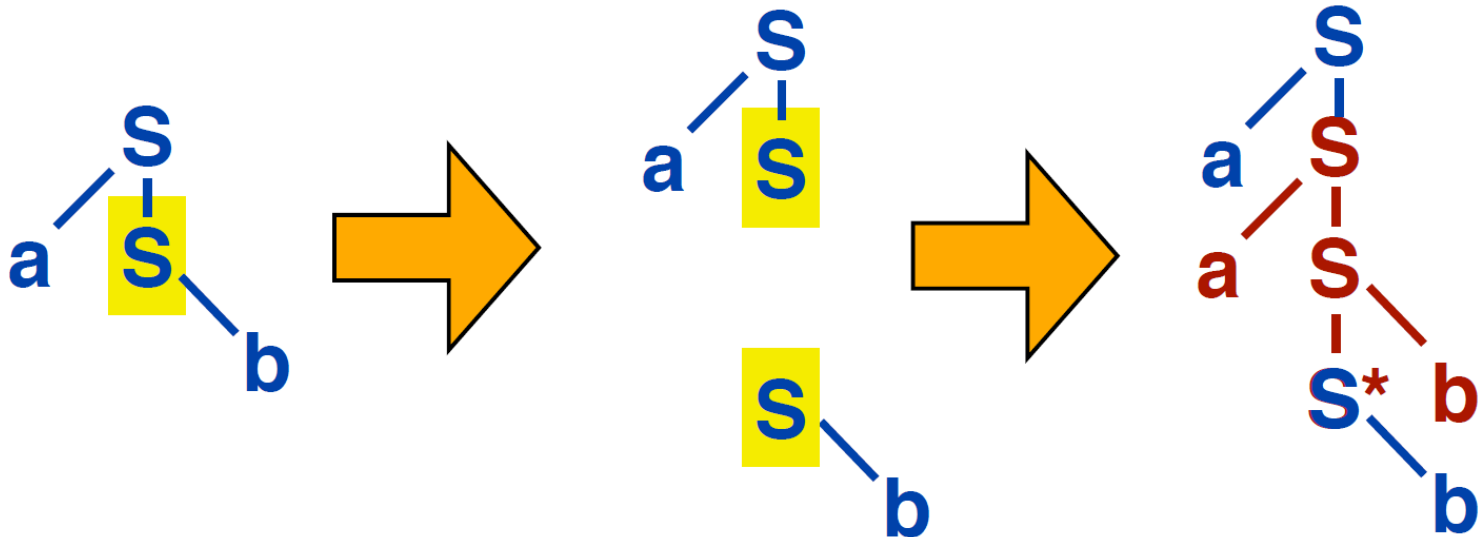
- Dutch and Swiss-German
- Can this be generated from context-free grammar?

# $a^n b^n$ : Cross-serial dependencies

Elementary trees:

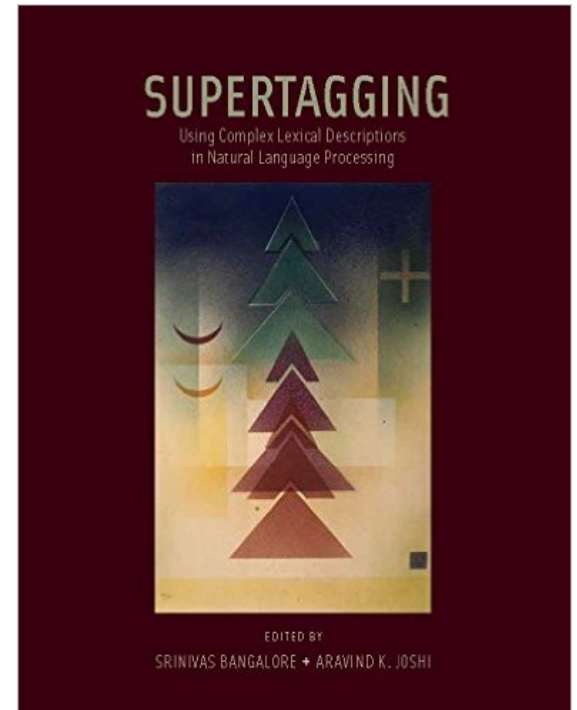


Deriving  $aabb$



# Tree Adjoining Grammar (TAG)

- TAG: Aravind Joshi in 1969
- Supertagging for TAG: Joshi and Srinivas 1994
- Pushing grammar down to lexicon.
- With just two rules: substitution & adjunction
- Parsing Complexity:
  - $O(N^7)$
- Xtag Project (TAG Penntree) (<http://www.cis.upenn.edu/~xtag/>)
- Local expert!
  - Fei Xia @ Linguistics (<https://faculty.washington.edu/fxia/>)



# II. Combinatory Categorical Grammar (CCG)

# Categories

- Categories = types
  - Primitive categories
    - N, NP, S, etc
  - Functions
    - a combination of primitive categories
    - S/NP, (S/NP) / (S/NP), etc
    - V, VP, Adverb, PP, etc

# Combinatory Rules

- ➔ **Application**
  - forward application:  $x/y \ y \rightarrow x$
  - backward application:  $y \ x \backslash y \rightarrow x$
- **Composition**
  - forward composition:  $x/y \ y/z \rightarrow x/z$
  - backward composition:  $y \backslash z \ x \backslash y \rightarrow x \backslash z$
  - (forward crossing composition:  $x/y \ y \backslash z \rightarrow x \backslash z$ )
  - (backward crossing composition:  $x \backslash y \ y/z \rightarrow x/z$ )
- **Type-raising**
  - forward type-raising:  $x \rightarrow y / (y \backslash x)$
  - backward type-raising:  $x \rightarrow y \backslash (y/x)$
- **Coordination <&>**
  - $x \ \text{conj} \ x \rightarrow x$

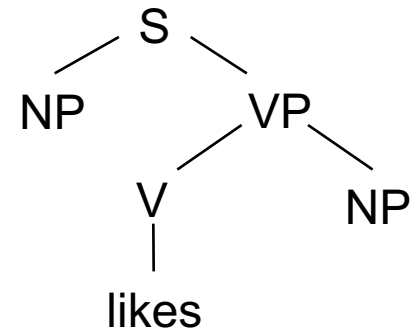
# Combinatory Rules 1 : Application

- Forward application “>”
  - $X/Y \ Y \rightarrow X$
  - $(S\NP)/NP \ NP \rightarrow S\NP$
- Backward application “<”
  - $Y \ X\Y \rightarrow X$
  - $NP \ S\NP \rightarrow S$

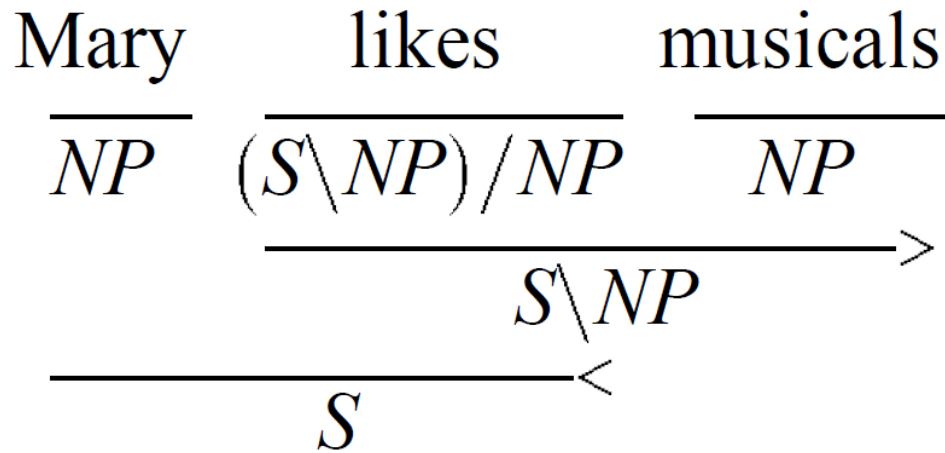


# Function

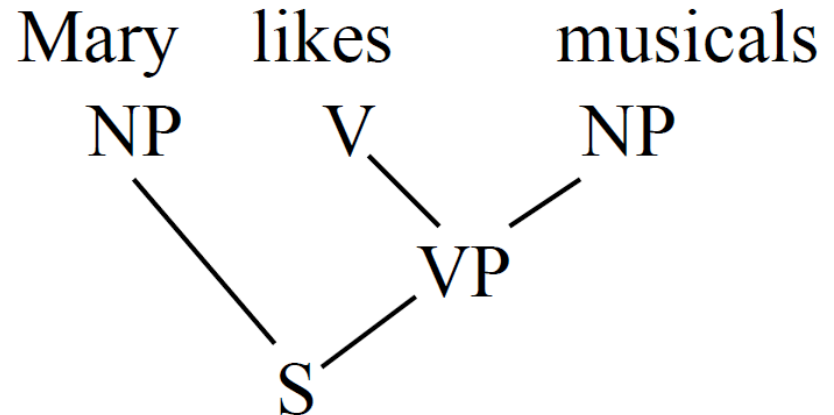
- likes := (S\NP) / NP
  - A transitive verb is a function from NPs into predicate S. That is, it accepts two NPs as arguments and results in S.
- Transitive verb: (S\NP) / NP
- Intransitive verb: S\NP
- Adverb: (S\NP) \ (S\NP)
- Preposition: (NP\NP) / NP
- Preposition: ((S\NP) \ (S\NP)) / NP



# CCG Derivation:



CFG Derivation:



# Combinatory Rules

- Application

- forward application:  $x/y \ y \rightarrow x$
- backward application:  $y \ x \backslash y \rightarrow x$

- Composition

- forward composition:  $x/y \ y/z \rightarrow x/z$
- backward composition:  $y \backslash z \ x \backslash y \rightarrow x \backslash z$
- forward crossing composition:  $x/y \ y \backslash z \rightarrow x \backslash z$
- backward crossing composition:  $x \backslash y \ y/z \rightarrow x/z$

- Type-raising

- forward type-raising:  $x \rightarrow y / (y \backslash x)$
- backward type-raising:  $x \rightarrow y \backslash (y/x)$



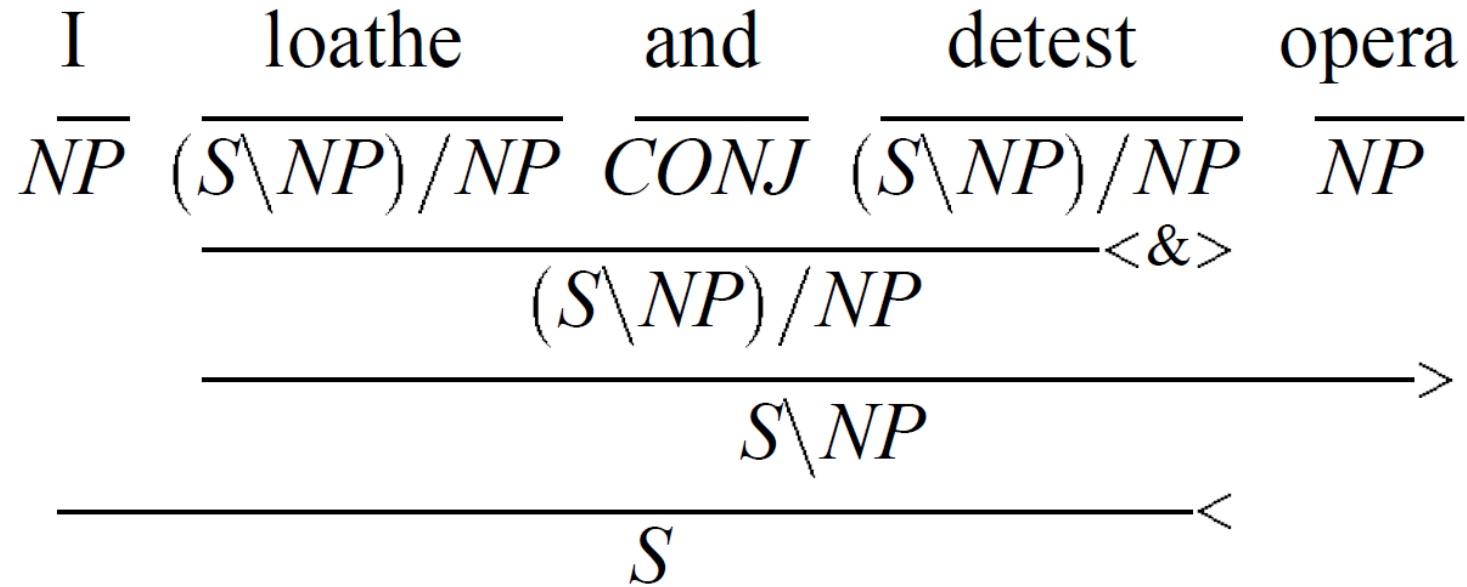
## Coordination <&>

- $x \ \text{conj} \ x \rightarrow x$

# Combinatory Rules 4 : Coordination

- $X \text{ conj } X \rightarrow X$
- Alternatively, we can express coordination by defining conjunctions as functions as follows:
  - $\text{and} := (X \setminus X) / X$

# Coordination with CCG



# Coordination with CCG

$\frac{\text{Marcel}}{\text{NP}}$

$\frac{\text{conjectured}}{(\text{S} \setminus \text{NP}) / \text{NP}}$

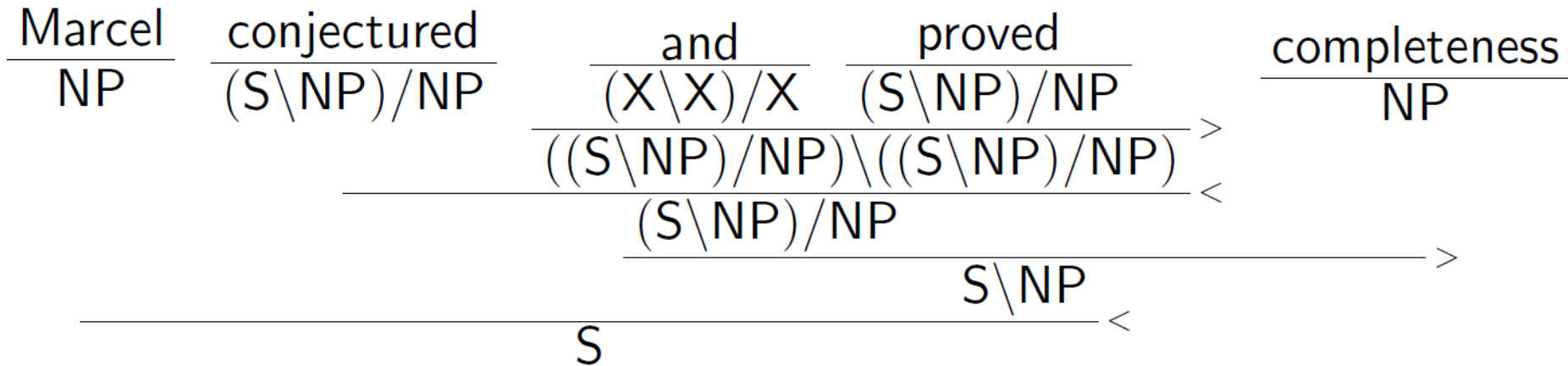
$\frac{\text{and}}{(\text{X} \setminus \text{X}) / \text{X}}$

$\frac{\text{proved}}{(\text{S} \setminus \text{NP}) / \text{NP}}$

$\frac{\text{completeness}}{\text{NP}}$

- Application
  - forward application:  $x/y \ y \rightarrow x$
  - backward application:  $y \ x \setminus y \rightarrow x$

# Coordination with CCG



- Application

- forward application:  $x/y \ y \rightarrow x$
- backward application:  $y \ x\backslash y \rightarrow x$

# Combinatory Rules

- Application

- forward application:  $x/y \ y \rightarrow x$
- backward application:  $y \ x \backslash y \rightarrow x$



## Composition

- forward composition:  $x/y \ y/z \rightarrow x/z$
- backward composition:  $y \backslash z \ x \backslash y \rightarrow x \backslash z$
- forward crossing composition:  $x/y \ y \backslash z \rightarrow x \backslash z$
- backward crossing composition:  $x \backslash y \ y/z \rightarrow x/z$

- Type-raising

- forward type-raising:  $x \rightarrow y / (y \backslash x)$
- backward type-raising:  $x \rightarrow y \backslash (y/x)$

- Coordination <&>

- $x \ \text{conj} \ x \rightarrow x$



# Coordination with CCG

$\frac{\text{Marcel}}{\text{NP}}$ 
 $\frac{\text{conjectured}}{(S \setminus \text{NP}) / \text{NP}}$ 
 $\frac{\text{and}}{(X \setminus X) / X}$ 
 $\frac{\text{might}}{(S \setminus \text{NP}) / ((S \setminus \text{NP}))}$ 
 $\frac{\text{prove}}{(S \setminus \text{NP}) / \text{NP}}$ 
 $\frac{\text{completeness}}{\text{NP}}$

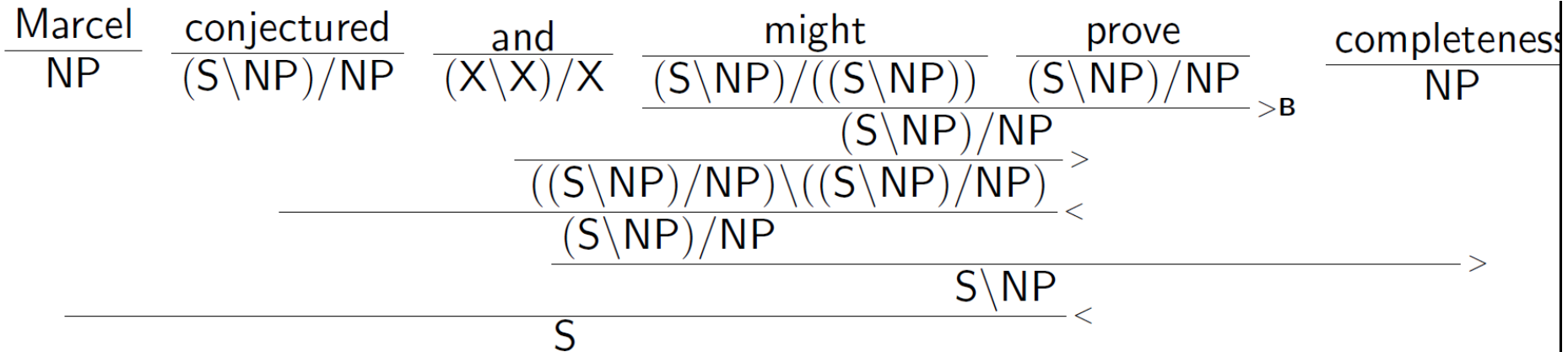
- Application

- forward application:  $x/y \ y \rightarrow x$
- backward application:  $y \ x \setminus y \rightarrow x$

- Composition

- forward composition:  $x/y \ y/z \rightarrow x/z$
- backward composition:  $y \setminus z \ x \setminus y \rightarrow x \setminus z$
- forward crossing composition:  $x/y \ y \setminus z \rightarrow x \setminus z$
- backward crossing composition:  $x \setminus y \ y/z \rightarrow x/z$

# Coordination with CCG



- Application

- forward application:  $x/y \ y \rightarrow x$
- backward application:  $y \ x \backslash y \rightarrow x$

- Composition

- forward composition:  $x/y \ y/z \rightarrow x/z$
- backward composition:  $y \backslash z \ x \backslash y \rightarrow x \backslash z$
- forward crossing composition:  $x/y \ y \backslash z \rightarrow x \backslash z$
- backward crossing composition:  $x \backslash y \ y/z \rightarrow x/z$

# Combinatory Rules

- Application

- forward application:  $x/y \ y \rightarrow x$
- backward application:  $y \ x \backslash y \rightarrow x$

- Composition

- forward composition:  $x/y \ y/z \rightarrow x/z$
- backward composition:  $y \backslash z \ x \backslash y \rightarrow x \backslash z$
- forward crossing composition:  $x/y \ y \backslash z \rightarrow x \backslash z$
- backward crossing composition:  $x \backslash y \ y/z \rightarrow x/z$

 Type-raising

- forward type-raising:  $x \rightarrow y / (y \backslash x)$
- backward type-raising:  $x \rightarrow y \backslash (y/x)$

- Coordination <&>

- $x \ \text{conj} \ x \rightarrow x$

# Combinatory Rules 3 : Type-Raising

- Turns an argument into a function
- Forward type-raising:  $X \rightarrow T / (T \backslash X)$
- Backward type-raising:  $X \rightarrow T \backslash (T / X)$

For instance...

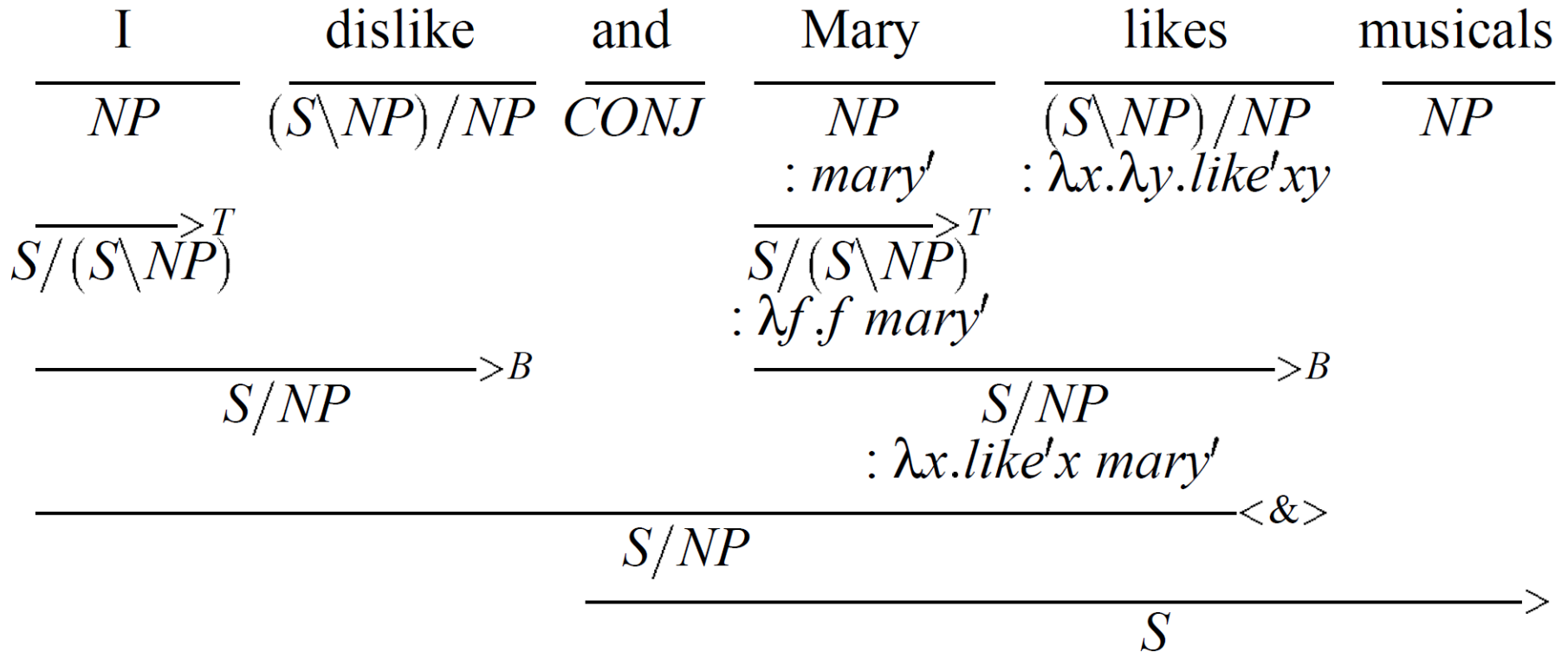
- Subject type-raising:  $NP \rightarrow S / (S \backslash NP)$
- Object type-raising:  $NP \rightarrow (S \backslash NP) \backslash ((S \backslash NP) / NP)$

# Combinatory Rules 3 : Type-Raising

|      |                         |        |      |                         |          |
|------|-------------------------|--------|------|-------------------------|----------|
| I    | dislike                 | and    | Mary | likes                   | musicals |
| $NP$ | $(S \setminus NP) / NP$ | $CONJ$ | $NP$ | $(S \setminus NP) / NP$ | $NP$     |

- Application
  - forward application:  $x/y \ y \rightarrow x$
  - backward application:  $y \ x \setminus y \rightarrow x$
- Type-raising
  - forward type-raising:  $x \rightarrow y / (y \setminus x)$
  - backward type-raising:  $x \rightarrow y \setminus (y/x)$
  - Subject type-raising:  $NP \rightarrow S / (S \setminus NP)$
  - Object type-raising:  $NP \rightarrow (S \setminus NP) \setminus ((S \setminus NP) / NP)$
- Coordination <&>
  - $x \ conj \ x \rightarrow x$

# Combinatory Rules 3 : Type-Raising



# Combinatory Categorical Grammar (CCG)

- CCG: Steedman in 1986
- Pushing grammar down to lexicon.
- With just a few rules: application, composition, type-raising
- We've looked at only syntactic part of CCG
- A lot more in the semantic part of CCG (using lambda calculus)
- Parsing Complexity:
  - $O(N^6)$
- Local expert!
  - Luke Zettlemoyer (<https://www.cs.washington.edu/people/faculty/lisz>)

