

# Polymorphic Type Inference Review Problems

CSE505

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Determine the (principal) types of the following expressions (using the Hindley-Milner algorithm), or say how the algorithm fails:

a.  $\lambda x. \lambda y. \lambda z. ((x\ z)\ (y\ z))$   
 $(\alpha \rightarrow \beta \rightarrow \gamma) \rightarrow (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \gamma$

b. let  $f = \lambda x. (\text{if } x \text{ then } x \text{ else } f(1))$  in  $f(0)$

The type inferred for  $\lambda x. (\text{if } x \text{ then } x \text{ else } f(1))$  is  $\text{Bool} \rightarrow \text{Bool}$ . In the process, the type of  $f$  is inferred to be  $\text{Int} \rightarrow \text{Bool}$  because of the call to  $f(1)$  in the body of the function (since  $f$  is treated as monomorphic within its own definition). These two types must be unified to get the final type of  $f$ , and this unification fails because  $\text{Int}$  and  $\text{Bool}$  cannot be unified.

c.  $\lambda x. (x\ x)$

When inferring the type of  $(x\ x)$ , we attempt to unify  $\alpha$  (the assumed type of  $x$ ) with  $\alpha \rightarrow \beta$ , for some fresh  $\beta$ . This unification fails because of the occurs check.

d.  $\lambda f. \lambda x. \text{if } x \text{ then } f(x) \text{ else } f(0)+1$

After inferring the type of  $f(x)$ , we have learned that  $f$  has a type of the form  $\text{Bool} \rightarrow \alpha$ . When inferring the type of  $f(0)$ , we attempt to unify  $\text{Bool} \rightarrow \alpha$  with  $\text{Int} \rightarrow \beta$ , for some fresh type variable  $\beta$ , which fails because  $\text{Bool}$  and  $\text{Int}$  cannot be unified.

e. let  $f = \lambda z. 0$  in  $\lambda x. \text{if } x \text{ then } f(x) \text{ else } f(0)+1$   
 $\text{Bool} \rightarrow \text{Int}$