Dynamic Languages

CSE 501

Spring 15

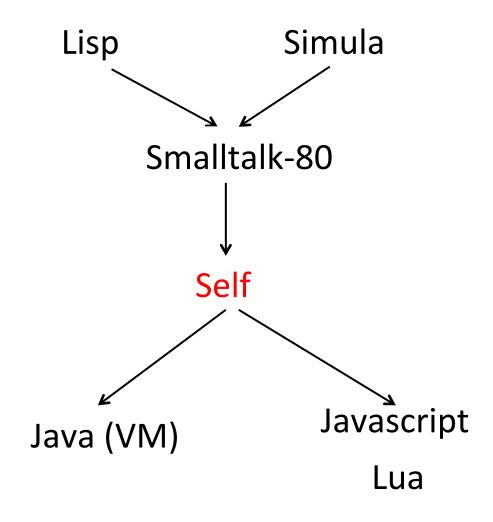
Dynamic Programming Languages

- Languages where program behavior, broadly construed, cannot be determined during compilation
 - Types
 - Code to be executed (eval in Javascript)
 - Loading external libraries
- Language examples
 - Javascript
 - Python
 - PHP
 - Smalltalk
 - Matlab

History of Self

- Prototype-based pure object-oriented language.
- Designed by Randall Smith (Xerox PARC) and David Ungar (Stanford) in 1987.
 - Successor to Smalltalk-80
 - Vehicle for implementation research
 - Later implementation by Craig Chambers and others at Stanford ← This is the one we are studying

History of SELF





fun through simplicity



Self Mallard Released!

The latest version of Self is Self "Mallard" 4.5.0 released January 2014. Download now!

Here is where to get Self:

Download for OS X



Includes the Self Control.app, Self VM and a prebuilt snapshot.

Download for Linux x86



Includes a Self VM and a prebuilt snapshot.

Use the Source, Luke



All of the Self sources for the VM and for the default Self World are on Github.

Design Goals

- Conceptual economy
 - Everything is an object
 - Everything done using messages
 - No classes
 - No variables
- Concreteness
 - Objects should seem "real"
 - GUI to manipulate objects directly

Language Overview

- Dynamically typed
 - Users do not declare types
- All computation via message passing
- Objects are organized into slots
- Operations on objects:
 - send messages
 - add new slots
 - replace old slots
 - remove slots

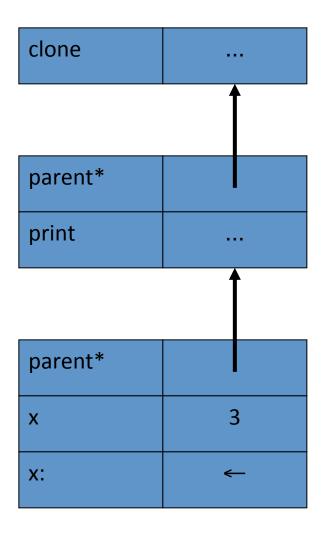
Objects and Slots

Object consists of named slots.

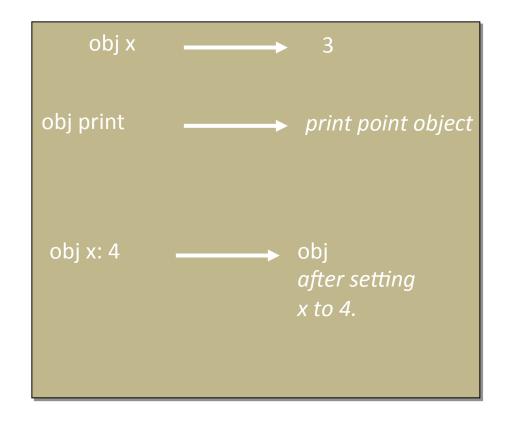
- Data
 - Such slots return contents upon evaluation; so act like instance variables
- Assignment
 - Set the value of associated slot
- Method
 - Slot contains Self code
- Parent
 - Point to existing object to inherit slots

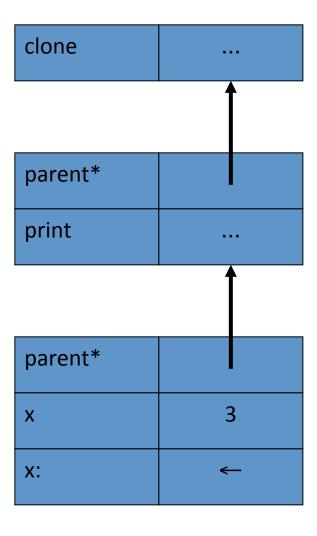
Messages and Methods

- When message is sent, object searched for slot with name.
- If none found, all parents are searched.
 - Runtime error if more than one parent has a slot with the same name.
- If slot is found, its contents evaluated and returned.
 - Runtime error otherwise

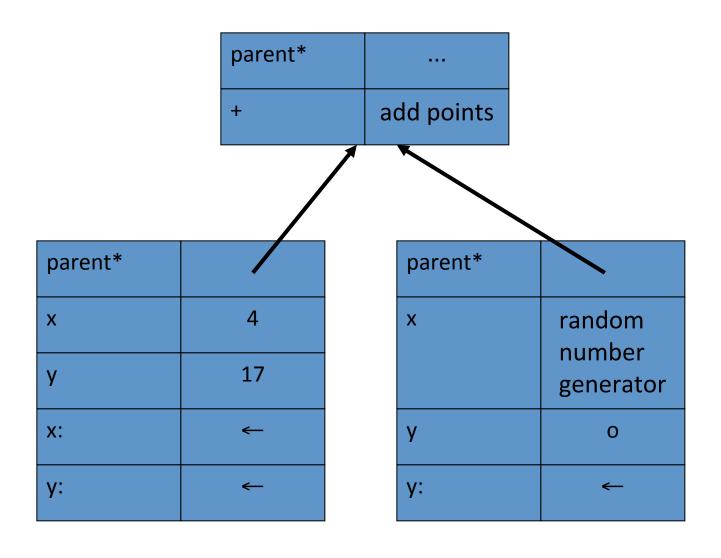


Messages and Methods





Mixing State and Behavior

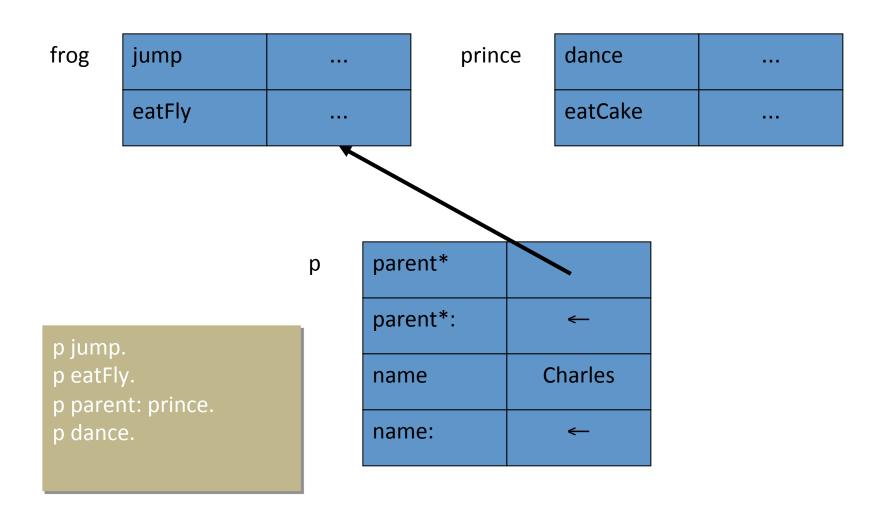


Creating Objects

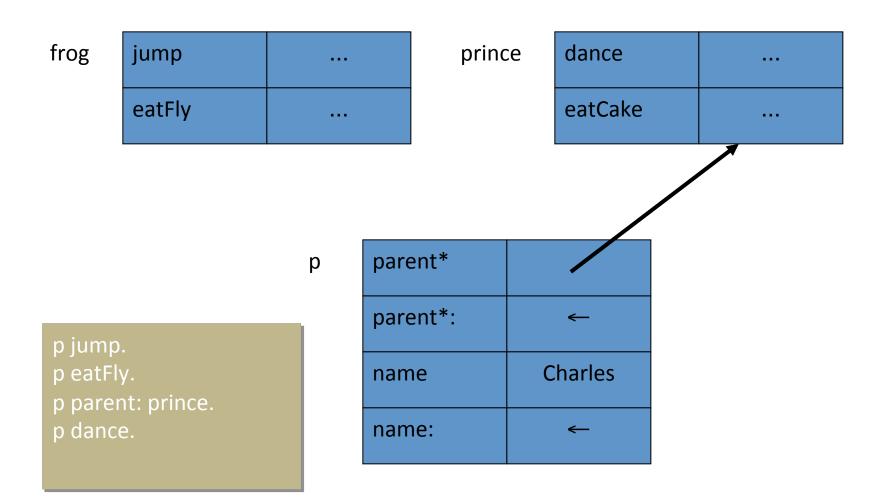
- To create an object, we copy an old one
- We can add new methods, override existing ones, or even remove methods as the program executes

 These operations also apply to parent slots as well

Changing Parent Pointers



Changing Parent Pointers



Why no classes?

- Classes require programmers to understand a more complex model.
 - To make a new kind of object, we have to create a new class first.
 - To change an object, we have to change the class.
 - Infinite meta-class regression.
- But: Does Self require programmer to reinvent structure?
 - Common to structure Self programs with traits:
 objects that simply collect behavior for sharing.

Contrast with C++

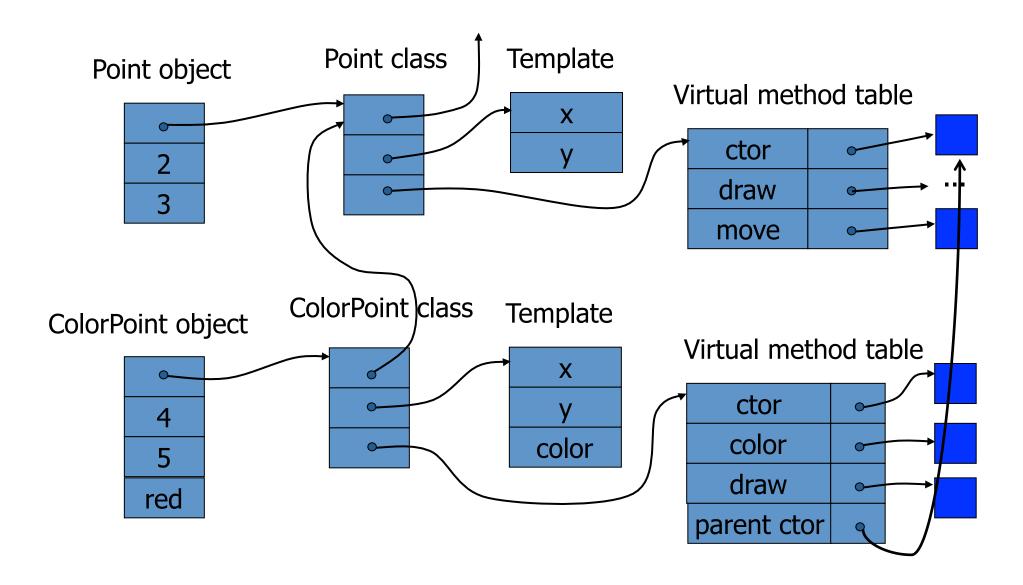
- C++
 - Restricts expressiveness to ensure efficient implementation
 - Class hierarchy is fixed during development
- Self
 - Provides high-level abstraction of underlying machine
 - Compiler does fancy optimizations to obtain acceptable performance

Implementation Challenges I

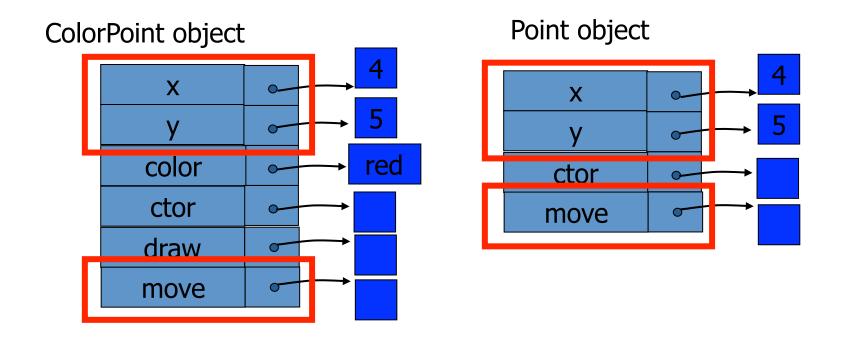
- Many, many slow function calls:
 - Function calls generally expensive.
 - Dynamic dispatch makes message invocation even slower than typical procedure calls.
 - OO programs tend to have lots of small methods.
 - Everything is a message: even variable access!

"The resulting call density of pure objectoriented programs is staggering, and brings naïve implementations to their knees" [Chambers & Ungar, PLDI 89]

C++ Object Layout



Naive Self Object Layout



Implementation Challenges II

- No static type system
 - Each reference could point to any object, making it hard to find methods statically.
- No class structure to enforce sharing
 - Each object having a copy of its methods leads to space overheads.

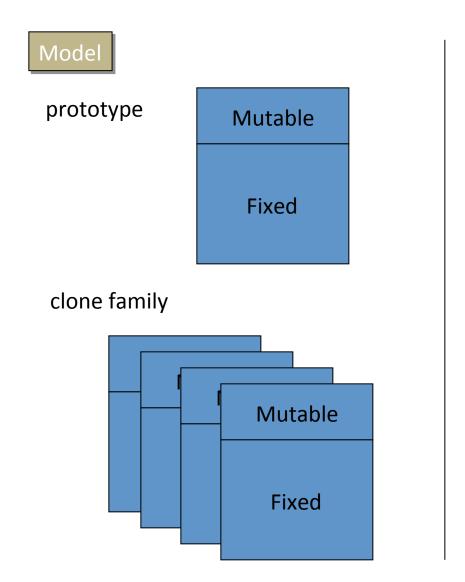
Optimized Smalltalk-80 roughly 10 times slower than optimized C

Optimization Strategies

- Avoid per object space requirements
- Avoid interpreting
 - Compile code instead
- Avoid method lookup
- Inline methods wherever possible
 - Saves method call overhead
 - Enables further optimizations

Clone Families

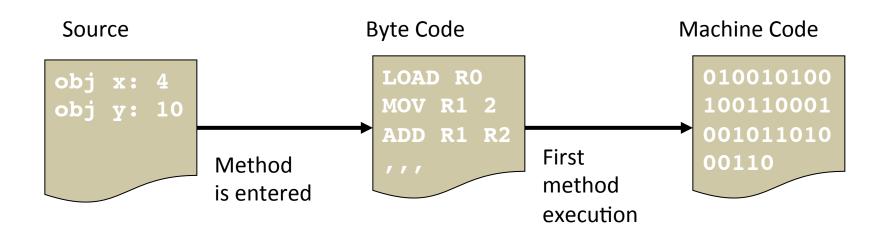
Avoid per object data



Implementation map **Fixed** Info Mutable Map

Dynamic Compilation

Avoid interpreting



- Method is converted to byte codes when entered
- Compiled to machine code when first executed
- Code stored in cache
 - if cache fills, previously compiled method flushed
- Requires entire source (byte) code to be available

Lookup Cache

Avoid method lookup

- Cache of recently used methods, indexed by (receiver type, message name) pairs.
- When a message is sent, compiler first consults cache
 - if found: invokes associated code
 - if absent: performs general lookup and potentially updates cache

- Compiler predicts types that are unknown but likely:
 - Arithmetic operations (+, -, <, etc.) have small integers as their receivers 95% of time in Smalltalk-80.
 - ifTrue had Boolean receiver 100% of the time.
- Compiler inlines code (and test to confirm guess):

```
if type = smallInt jump to method_smallInt
else call general_lookup
```

Inline Caches

Avoid method lookup

- First message send from a call site :
 - general lookup routine invoked
 - call site back-patched
 - is previous method still correct?
 - yes: invoke code directly
 - no: proceed with general lookup & backpatch
- Successful about 95% of the time
- All compiled implementations of Smalltalk and Self use inline caches

Polymorphic Inline Caches

- Typical call site has <10 distinct receiver types
 - So often can cache all receivers
- At each call site, for each new receiver, extend patch code:

```
if type = rectangle jump to method_rect
if type = circle     jump to method_circle
call general_lookup
```

- After some threshold, revert to simple inline cache (megamorphic site)
- Order clauses by frequency
- Inline short methods

Customized Compilation

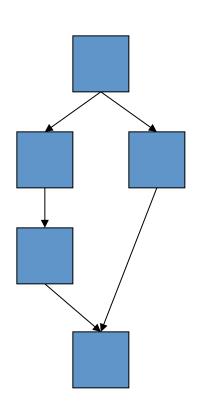
Inline methods

- Compile several copies of each method, one for each receiver type
- Within each copy:
 - Compiler knows the type of self
 - Calls through self can be statically selected and inlined
- Enables downstream optimizations
- Increases code size

Type Analysis

Inline methods

- Constructed by compiler by flow analysis.
- Type: set of possible maps for object
 - Singleton: know map statically
 - Union/Merge: know expression has one of a fixed collection of maps.
 - Unknown: know nothing about expression.
- If singleton, we can inline method.
- If type is small, we can insert type test and create branch for each possible receiver (type casing)



Performance Improvements

 Initial version of Self was 4-5 times slower than optimized C.

 After optimizations, implementation described in paper is within a factor of 2 of optimized C.

How successful?

- Few users: not a popular success
 - No compelling application, until JavaScript
 - Influenced development of object calculi w/o classes
- However, many research innovations
 - Very simple computational model
 - Enormous advances in compilation techniques
 - Influenced the design of Java compilers
 - Direct influence on design of Javascript

Lessons

Pochoir / Halide (DSL)

- Design specific constructs for domain
- Constructs need to easily map to underlying target language
 - Otherwise implementation might be a nightmare
- Expose high-level structure allows domain-specific optimizations

<u>Self / Javascript</u> (Dynamic Languages)

- "Power of simplicity"
 - Everything is an object
 - No classes, no variables
- Implementation specific to program constructs

 Uses various optimization tricks to recover performance