Threat modeling warm-up

Previously, a link to:

twitter.com / definitely_a_real_user / status / 34273525929

Would appear as exactly that on twitter

Now they have started automatically string search-and-replacing "twitter.com" with 'x.com' *only in the rendering of the link*.

x.com / definitely_a_real_user / status / 34273525929

This particular example would then still go to twitter.com.

How could a malicious user abuse this search-and-replace behavior on link rendering?

CSE 484: Computer Security and Privacy

Cryptography basics

Spring 2024

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Logistics

- Lab 1a due tonight!
  - Remember, up to 3 late days (out of 5 for the quarter) per-assignment
  - Stuck on something? Try debugging and tracing _normal_ execution, then your corrupted execution!
  - Sploit3 confusing? Breakpoint at the last instruction you are confident in, and step instruction at a time. Examine the stack + registers at each one!
- Lab 1b is due in a week (and is significantly trickier. Start ASAP)
- Make sure you did your writeups independently

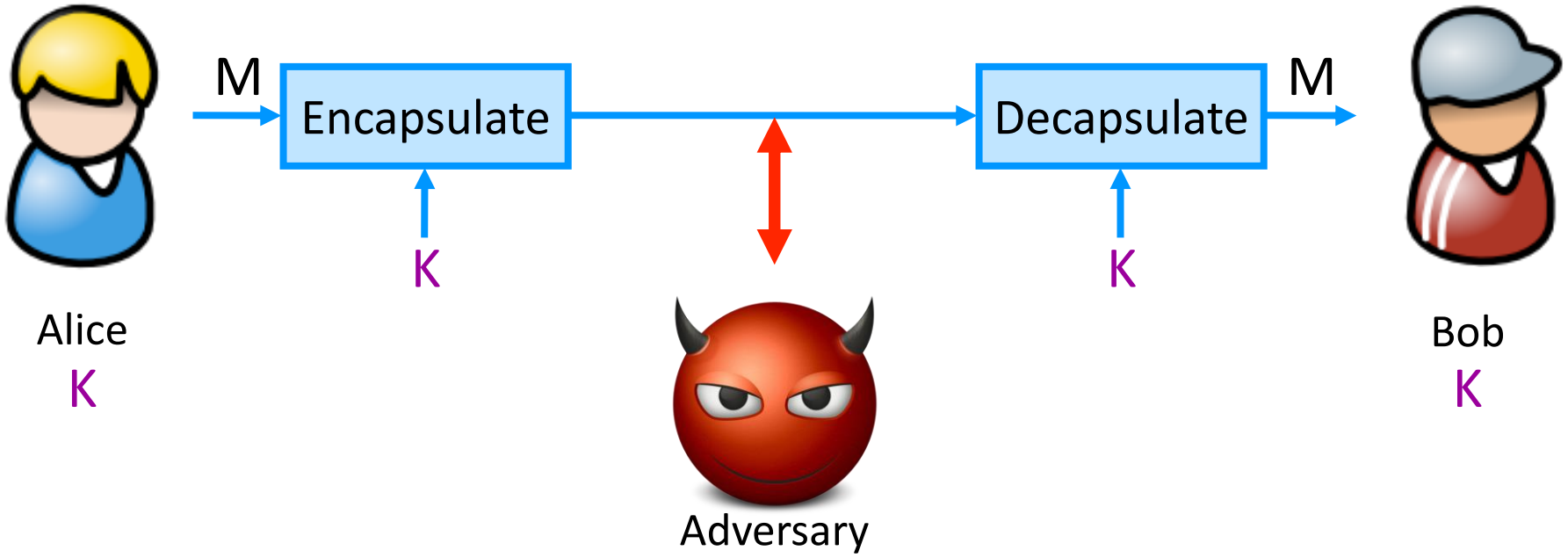
Flavors of Cryptography

- Symmetric cryptography
  - Both communicating parties have access to a shared random string K, called the key.

- Asymmetric cryptography
  - Each party creates a public key pk and a secret key sk.
  - *Hard concept to understand, and revolutionary! Inventors won Turing Award* ☺

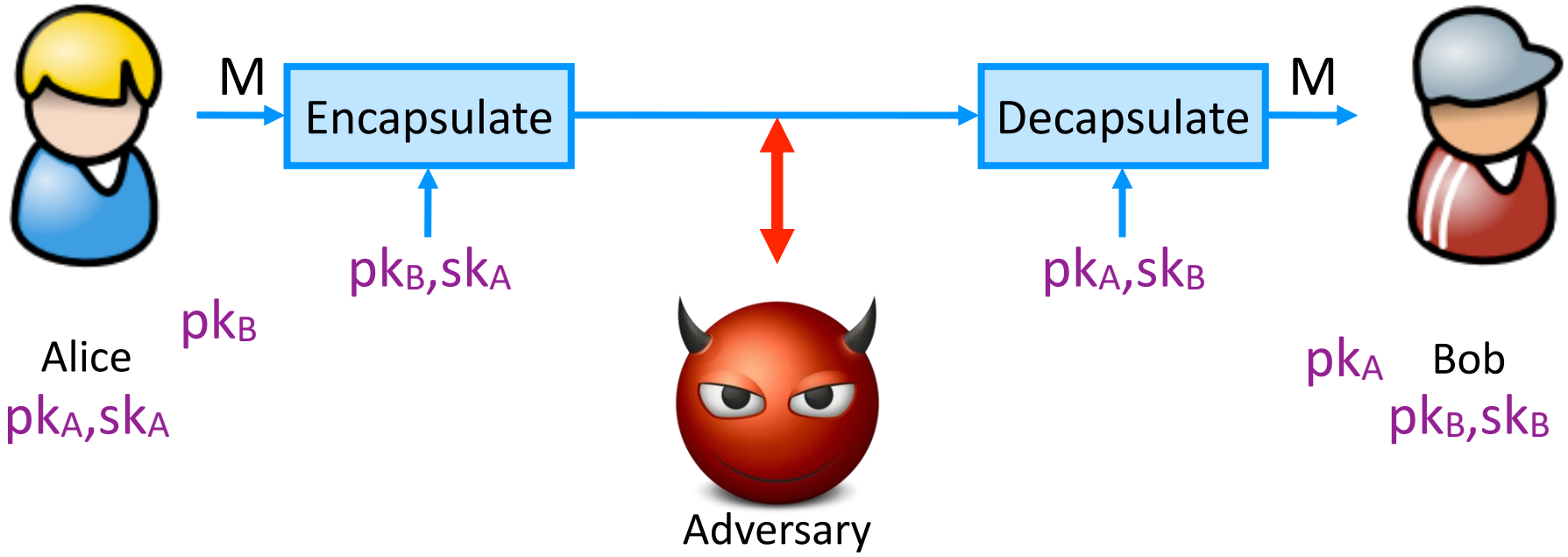
Symmetric Setting

Both communicating parties have access to a shared random string K, called the key.



Asymmetric Setting

Each party creates a public key pk and a secret key sk.



Public keys, Private keys, Secret keys…

- Secret key
  - The single key used in symmetric encryption
  - The non-public key in asymmetric
- Private keys
  - The non-public key in asymmetric
- Public key
  - The… public key in asymmetric
- Key
  - Generally means private/secret

Properties of asymmetric cryptography

- We have a funny situation here:
  - Public keys are shared with everyone
  - Secret keys are not

- What is are some security properties we would want of:
  - Knowing a public key?
  - Encrypting a message with a secret key?

Received April 4, 1977

A Method for Obtaining Digital Signatures and Public-Key Cryptosystems

R.L. Rivest, A. Shamir, and L. Adleman*

Abstract

An encryption method is presented with the novel property that publicly revealing an encryption key does not thereby reveal the corresponding decryption key. This has two important consequences:

1. Couriers or other secure means are not needed to transmit keys, since a message can be enciphered using an encryption key publicly revealed by the intended recipient. Only he can decipher the message, since only he knows the corresponding decryption key.
2. A message can be "signed" using a privately held decryption key. Anyone can verify this signature using the corresponding publicly revealed encryption key. Signatures cannot be forged, and a signer cannot later deny the validity of his signature. This has obvious applications in "electronic mail" and "electronic funds transfer" systems.

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Flavors of Cryptography

- Symmetric cryptography
  - Both communicating parties have access to a shared random string K, called the key.
  - Challenge: How do you privately share a key?
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  - Each party creates a public key pk and a secret key sk.
  - Challenge: How do you validate a public key?

Flavors of Cryptography

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- Asymmetric cryptography
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  - Challenge: How do you validate a public key?
- Key building block: Randomness – something that the adversaries won't know and can't predict and can't figure out

Detour: Randomness

Ingredient: Randomness

- Many applications (especially security ones) require randomness
- Explicit uses:
  - Generate secret cryptographic keys
  - Generate random initialization vectors for encryption
- Other “non-obvious” uses:
  - Generate passwords for new users
  - Shuffle the order of votes (in an electronic voting machine)
  - Shuffle cards (for an online gambling site)

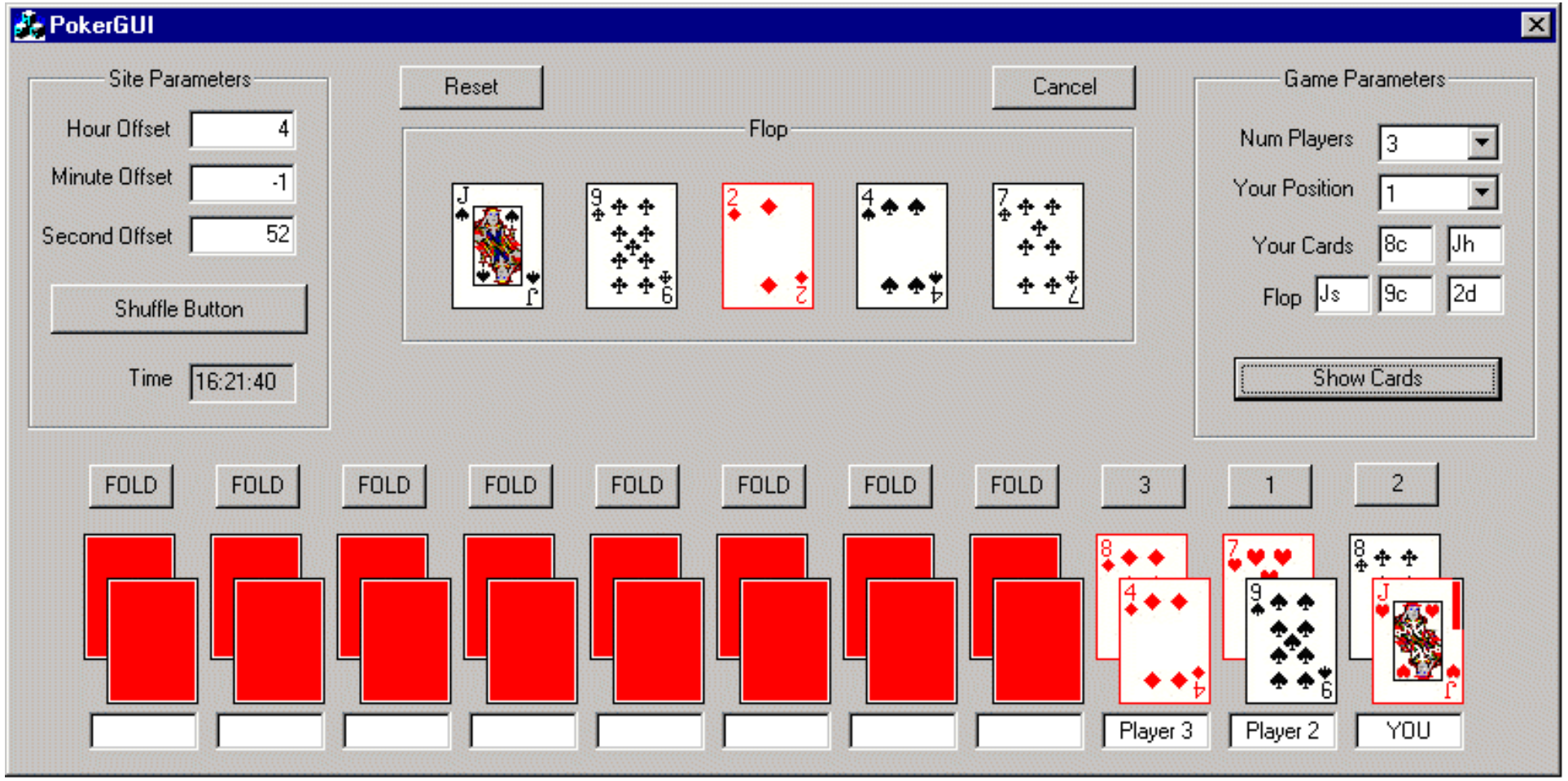
C's rand() Function

- C has a built-in random function: rand()

```
unsigned long int next = 1;
/* rand:  return pseudo-random integer on 0..32767 */
int rand(void) {
    next = next * 1103515245 + 12345;
    return (unsigned int)(next/65536) % 32768;
}
/* srand:  set seed for rand() */
void srand(unsigned int seed) {
    next = seed;
}
```

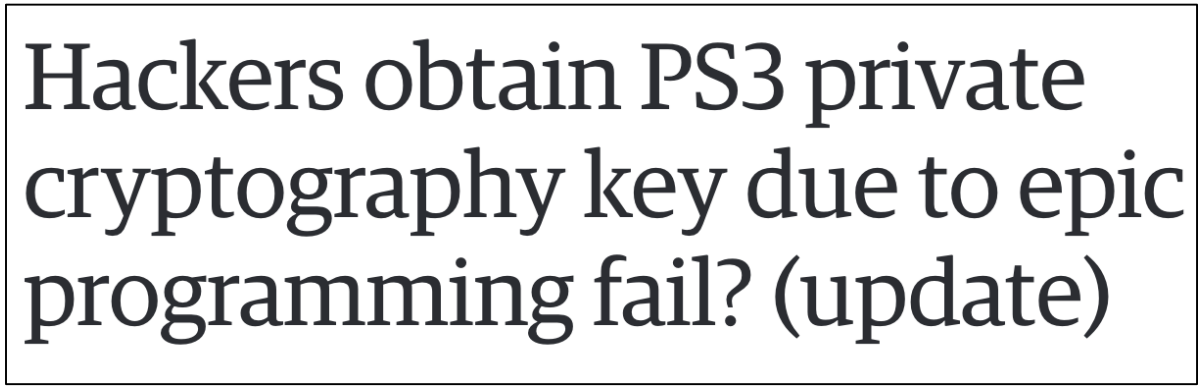
- Problem: don't use rand() for security-critical applications!
  - Given a few sample outputs, you can predict subsequent ones


A World of Action!
Help
Suggestions
Fold in Turn
$0
jonnyboy
brute
Fold
Check
Bet ($5)
mamajoe
RickG
Options
Sit Out
Leave
584330
mamajoe: Hey guys, Big B is in!
Dealer Text



More details: “How We Learned to Cheat at Online Poker: A Study in Software Security”
https://web.archive.org/web/20120301022831/http://www.cigital.com/papers/download/developer_gambling.php

PS3 and Randomness



http://www.engadget.com/2010/12/29/hackers-obtain-ps3-private-cryptography-key-due-to-epic-programm/

- 2010/2011: Hackers found/released private root key for Sony's PS3
- Key used to sign software – now can load any software on PS3 and it will execute as "trusted"
- Due to bad random number: same "random" value used to sign all system updates

A recent example: keypair

https://securitylab.github.com/advisories/GHSL-2021-1012-keypair/

- keypair is a JS library for generating (asymmetric) keypairs

The output from the Lehmer LCG is encoded incorrectly. The specific line with the flaw is:

```
b.putByte(String.fromCharCode(next & 0xFF))
```

The definition of putByte is

```
[…]putByte = function(b) { this.data += String.fromCharCode(b); };
```

Since we are masking with 0xFF, we can determine that 97% of the output from the LCG are converted to zeros. The only outputs that result in meaningful values are outputs 48 through 57, inclusive.

The impact is that each byte in the RNG seed has a 97% chance of being 0 due to incorrect conversion. When it is not, the bytes are 0 through 9.

How might we get “good” random numbers?

Obtaining Pseudorandom Numbers

- For security applications, want "cryptographically secure pseudorandom numbers"
- Libraries include cryptographically secure pseudorandom number generators (CSPRNG)

Obtaining Pseudorandom Numbers

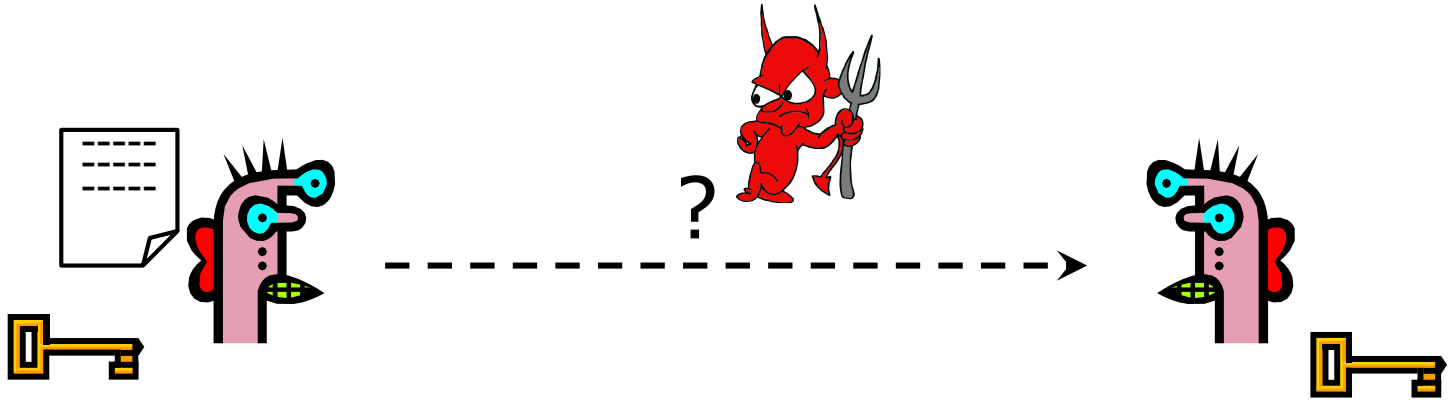
- Linux:
  - /dev/random – blocking (waits for enough entropy)
  - /dev/urandom – nonblocking, possibly less entropy
  - getrandom() – syscall! – by default, blocking
- Internally:
  - Entropy pool gathered from multiple sources
    - e.g., mouse/keyboard/network timings
- Challenges with embedded systems, saved VMs

Obtaining *Random* Numbers

- Better idea:
  - AMD/Intel's on-chip random number generator
    - RDRAND
- Hopefully no hardware bugs!

Back to encryption

Confidentiality: Basic Problem



Given (Symmetric Crypto): both parties know the same secret.

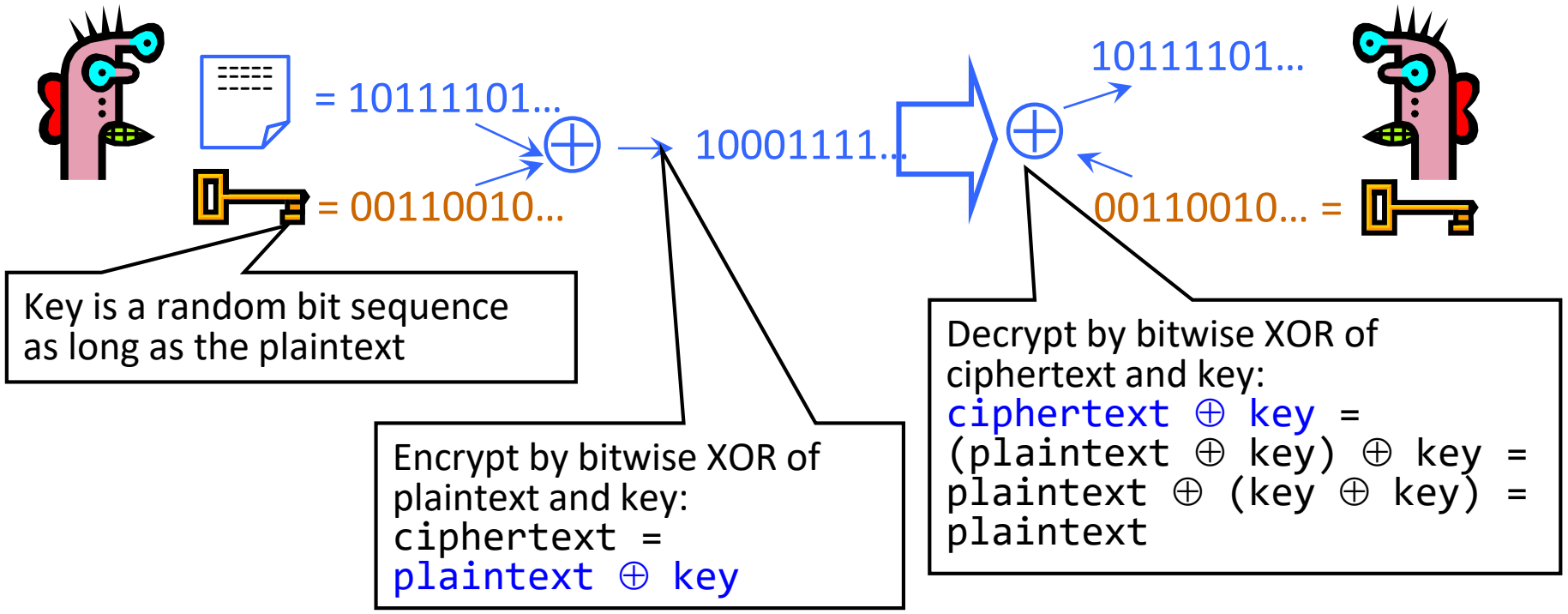
Goal: send a message confidentially.

Ignore for now: How is this achieved in practice??

One weird bit-level trick

- XOR!
  - Just XOR with a random bit!
- Why?
  - Uniform output
  - Independent of ‘message’ bit

One-Time Pad



Cipher achieves perfect secrecy if and only if
there are as many possible keys as possible plaintexts,
and every key is equally likely (Claude Shannon, 1949)

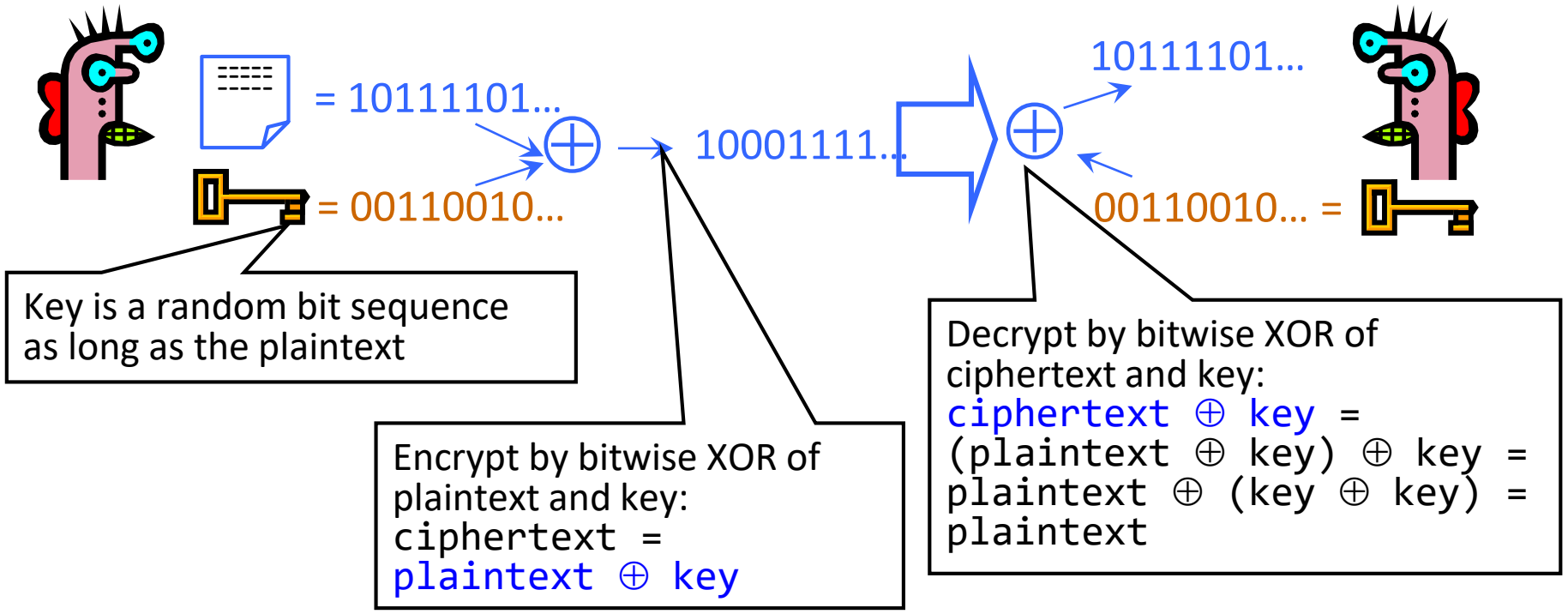
Advantages of One-Time Pad

- Easy to compute
  - Encryption and decryption are the same operation
  - Bitwise XOR is very cheap to compute
- As secure as theoretically possible
  - Given a ciphertext, all plaintexts are equally likely, regardless of attacker's computational resources
  - …as long as the key sequence is truly random
    - True randomness is expensive to obtain in large quantities
  - …as long as each key is same length as plaintext
    - But how does sender communicate the key to receiver?

Problems with the One-Time Pad?

- Gradescope!
- What potential security problems do you see with the one-time pad?
- (Try not to look ahead and next slides)
- Recall two key goals of cryptography: confidentiality and integrity
  - Assume we are sending the message over an untrusted medium
  - Remember our different adversaries!

One-Time Pad - Reminder



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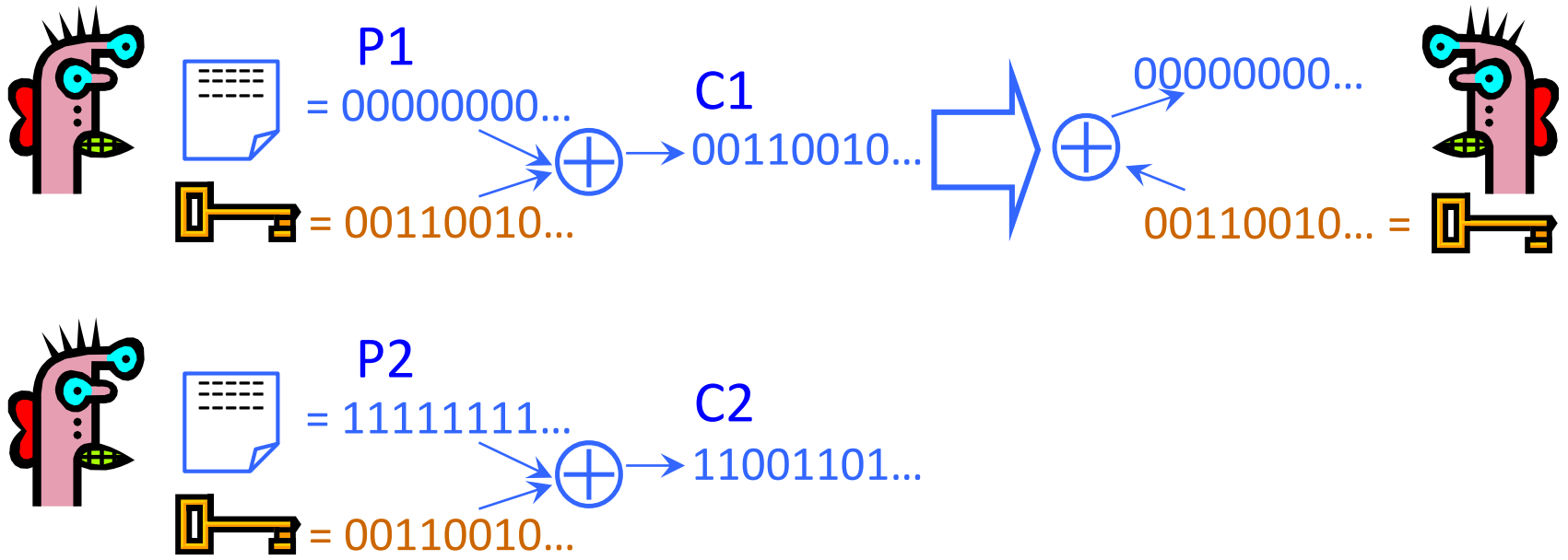
Dangers of Reuse



Learn relationship between plaintexts

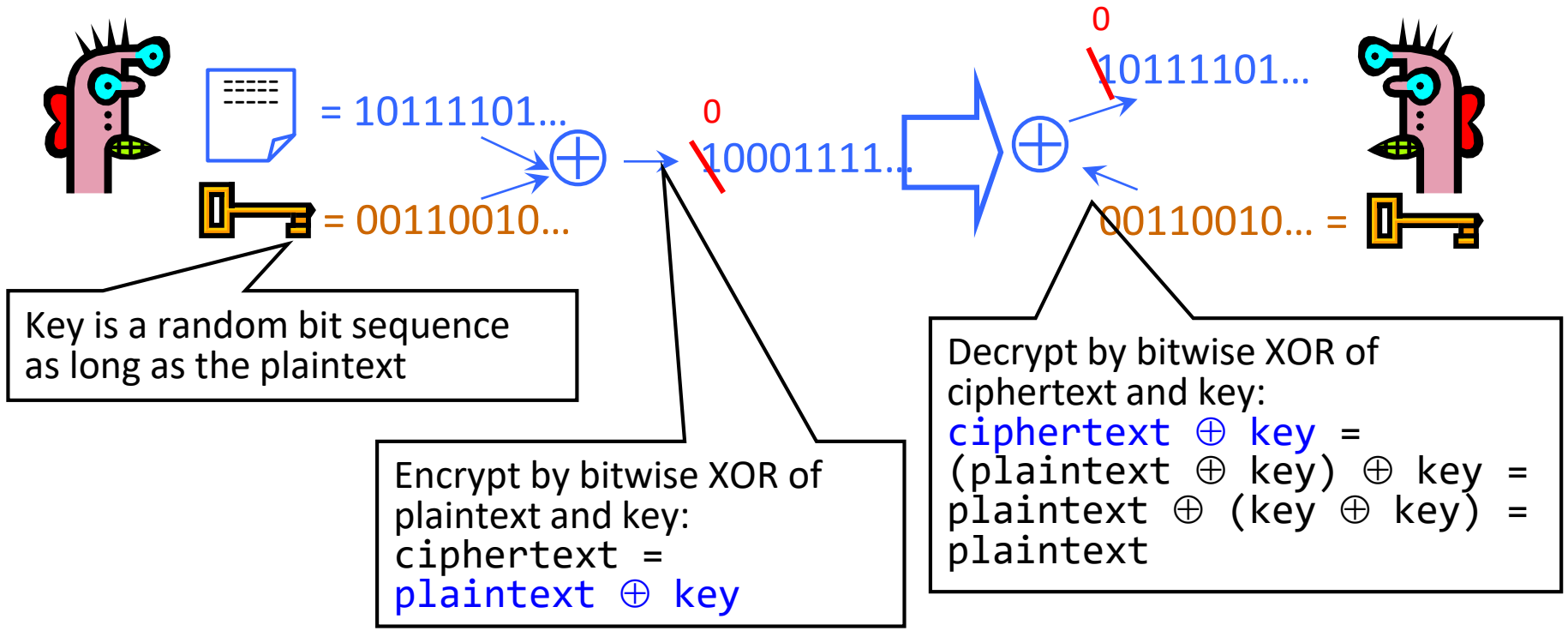
C1⊕C2 = (P1⊕K)⊕(P2⊕K) =

(P1⊕P2)⊕(K⊕K) = P1⊕P2

Problems with One-Time Pad

- (1) Key must be as long as the plaintext
  - Impractical in most realistic scenarios
  - Still used for diplomatic and intelligence traffic
- **(2) Insecure if keys are reused**
  - **Attacker can obtain XOR of plaintexts**

Integrity?


= 10111101…
= 00110010…
Key is a random bit sequence as long as the plaintext
0
10001111…
Encrypt by bitwise XOR of plaintext and key:
ciphertext = plaintext ⊕ key
0
10111101…
00110010… =
Decrypt by bitwise XOR of ciphertext and key:
ciphertext ⊕ key = (plaintext ⊕ key) ⊕ key = plaintext ⊕ (key ⊕ key) = plaintext

Problems with One-Time Pad

- (1) Key must be as long as the plaintext
  - Impractical in most realistic scenarios
  - Still used for diplomatic and intelligence traffic
- (2) Insecure if keys are reused
  - Attacker can obtain XOR of plaintexts
- **(3) Does not guarantee integrity**
  - **One-time pad only guarantees confidentiality**
  - **Attacker cannot recover plaintext, but can easily change it to something else**