CSE 484 / CSE M 584
Computer Security: Buffer Overflows II

TA: Adrian Sham
adrsham@cs
Original slides by Franzi
Lab 1 Deadline Reminders

• Lab 1 Checkpoint (Sploits 1-3) due next Monday (4/20) at 5pm!
  – Turn in text file of md5sums for sploits 1-3
• Lab 1 Final due in two weeks (5/1, 5pm).
• If you don’t have a group or ssh access yet, talk to me today!
• Upcoming office hours:
  – Tomorrow (Friday) 9:30 am – Michael & Adrian
  – Monday 9:30 pm – Franzi
  – Wednesday 3:30 pm – Adrian & Peter
  – Thursday 12:30 pm – Peter & Michael
Lab 1 Notes/Hints

• If you get stuck, move on!
• Don’t procrastinate on Sploits 4-7. Some of them are much harder.
• Sploit 3: No frame pointer, so you can only change last byte of saved EIP. Think about an existing instruction you could point to that would have desirable side effects.
• You have more than one copy of your buffer: (1) as argument to function, (2) where it gets copied.
• Sploit 4 is not necessarily harder than Sploit 3.
Sploit 5 Tips

• Buffer copied to the heap.
• Target 5 uses the implementation that’s found in ~/sources/tmalloc.c.
• Read “Once upon a free()”:
Dynamic Memory Management in C

• Memory allocation: `malloc(size_t n)`
  – Allocates n bytes and returns a pointer to the allocated memory; memory not cleared.

• Memory deallocation: `free(void * p)`
  – Frees the memory space pointed to by p, which must have been returned by a previous call to `malloc()` (or similar).
    – If `free(p)` has been called before ("double free"), undefined behavior occurs.
  – If p is null, no operation is performed.

(Some memory management slides adapted from Vitaly Shmatikov)
Target5: What’s the problem?

```c
char *p; char *q;

if ( (p = tmalloc(128)) == NULL)
{ exit(EXIT_FAILURE); }

if ( (q = tmalloc(128)) == NULL)
{exit(EXIT_FAILURE); }

tfree(p);
tfree(q);

if ( (p = tmalloc(256)) == NULL)
{exit(EXIT_FAILURE); }

obsd_strlcpy(p, arg, 256);
{tfree(q); /* Undefined” behavior on second free()*/
```
Free Chunks (as used in tmalloc.c)

- Chunks organized into doubly-linked list.
- Each chunk on list contains forward/back pointers to next/previous chunks in the list.
  - LSB of right pointer contains free bit.
  - Adjacent free chunks are consolidated.
Chunk Maintenance

One big free chunk:

Split to malloc:

Split to malloc (twice):

Free (twice):

Consolidate free chunks:
Chunks in tmalloc.c

- Lines 20-28 give chunk structure:

```c
q = p->s.l;
...
q->s.r = p->s.r;
p->s.r->s.l = q;
```

- Look at chunk consolidation in `tfree(p)`:

Hey look, if we control chunks `p` (and `q`), this code will write the value of `q` (address of buffer?) to a location we specify (location of saved EIP?).

- Goal: populate (fake) chunks appropriately.
General Questions?