CSE 484 (Winter 2008)

Applied Cryptography

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Goals for Today

- Symmetric
- Reminder: Midterm on Friday. (Closed book.)
  - Contents up through the material for today
  - Not as hard as last year's midterm.
  - Make sure you understand the core concepts so far in this course:
    - Threat modeling
    - Software security
      - Problems
      - Defensive approaches
    - Symmetric cryptography
      - Components, definitions, security properties, classic problems

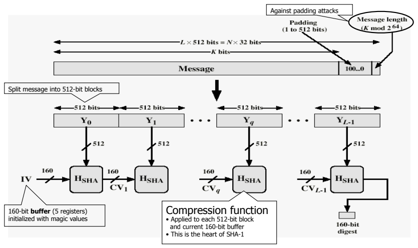
Which Property Do We Need?

- UNIX passwords stored as hash(password)
  - One-wayness: hard to recover password
- Integrity of software distribution
  - Weak collision resistance
  - But software images are not really random… maybe need full collision resistance
- Auction bidding
  - Alice wants to bid B, sends H(B), later reveals B
  - One-wayness: rival bidders should not recover B
  - Collision resistance: Alice should not be able to change her mind to bid B' such that H(B)=H(B')

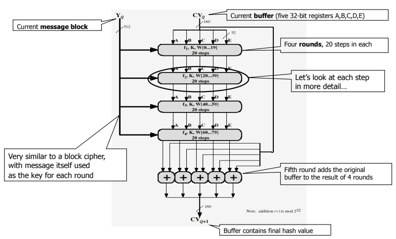
Common Hash Functions

- MD5
  - 128-bit output
  - Designed by Ron Rivest, used very widely
  - Collision-resistance broken (summer of 2004)
- RIPEMD-160
  - 160-bit variant of MD5
- SHA-1 (Secure Hash Algorithm)
  - 160-bit output
  - US government (NIST) standard as of 1993-95
    - Also the hash algorithm for Digital Signature Standard (DSS)

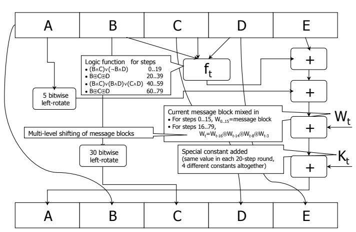
Basic Structure of SHA-1 (Skip)



SHA-1 Compression Function (Skip)



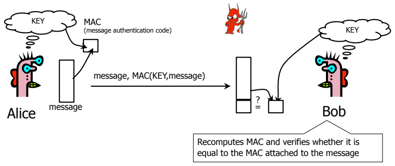
One Step of SHA-1 (80 steps total) (Skip)



How Strong Is SHA-1?

- Every bit of output depends on every bit of input
  - Very important property for collision-resistance
- Brute-force inversion requires 2^{160} ops, birthday attack on collision resistance requires 2^{80} ops
- Some very recent weaknesses (2005)
  - Collisions can be found in 2^{63} ops

Authentication Without Encryption

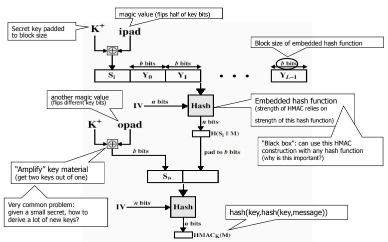


Integrity and authentication: only someone who knows KEY can compute MAC for a given message

HMAC

- Construct MAC by applying a cryptographic hash function to message and key
  - Could also use encryption instead of hashing, but…
  - Hashing is faster than encryption in software
  - Library code for hash functions widely available
  - Can easily replace one hash function with another
  - There used to be US export restrictions on encryption
- Invented by Bellare, Canetti, and Krawczyk (1996)
  - HMAC strength established by cryptographic analysis
- Mandatory for IP security, also used in SSL/TLS

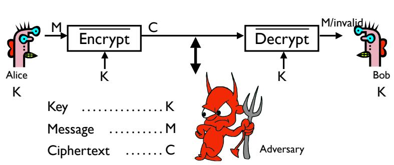
Structure of HMAC



Achieving Both Privacy and Integrity

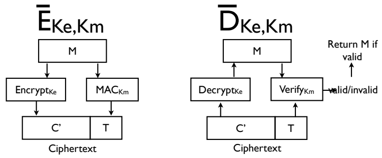
Authenticated encryption scheme

Recall: Often desire both privacy and integrity. (For SSH, SSL, IPsec, etc.)



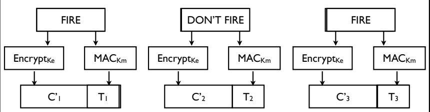
Some subtleties! Encrypt-and-MAC

Natural approach for authenticated encryption: Combine an encryption scheme and a MAC.



But insecure! [BN, Kra]

Assume Alice sends messages:



If $T_i = T_j$ then $M_i = M_j$

Adversary learns whether two plaintexts are equal.

Especially problematic when $M_1, M_2, \ldots$ take on only a small number of possible values.

The Secure Shell (SSH) protocol is designed to provide:

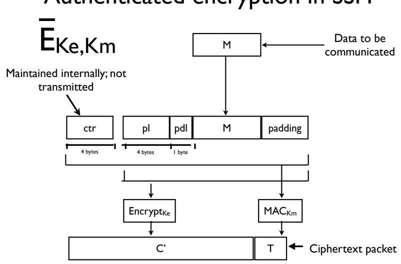
- Secure remote logins.
- Secure file transfers.

Where security includes:

- Protecting the privacy of users' data.
- Protecting the integrity of users' data.

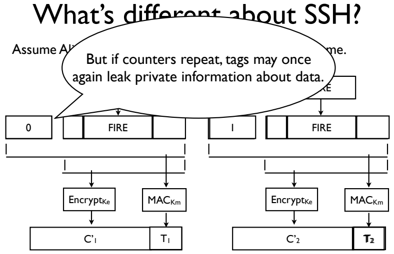
OpenSSH is included in the default installations of OS X and many Linux distributions.

Authenticated encryption in SSH



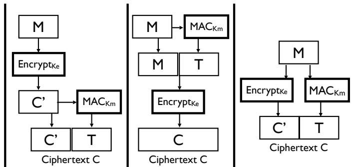
What's different about SSH?

Assume A… …me.



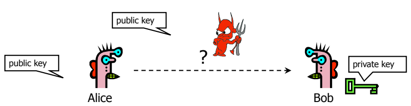
Then the tags T_1 and T_2 will be different with high probability.

Results of [BN00,Kra01]



	Encrypt-then-MAC	MAC-then-Encrypt	Encrypt-and-MAC
Privacy	Strong (CCA)	Weak (CPA)	Insecure
Integrity	Strong (CTXT)	Weak (PTXT)	Weak (PTXT)

Basic Problem



Given: Everybody knows Bob's public key
Only Bob knows the corresponding private key

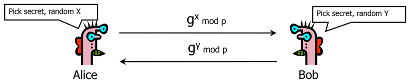
Goals: 1. Alice wants to send a secret message to Bob
2. Bob wants to authenticate himself

Applications of Public-Key Crypto

- Encryption for confidentiality
  - Anyone can encrypt a message
    - With symmetric crypto, must know secret key to encrypt
  - Only someone who knows private key can decrypt
  - Key management is simpler (maybe)
    - Secret is stored only at one site: good for open environments
- Digital signatures for authentication
  - Can "sign" a message with your private key
- Session key establishment
  - Exchange messages to create a secret session key
  - Then switch to symmetric cryptography (why?)

Diffie-Hellman Protocol (1976)

- Alice and Bob never met and share no secrets
- <u>Public</u> info: p and g
  - p is a large prime number, g is a generator of Z_p^*
    - $Z_p^* = \{1, 2 \ldots p-1\}$; $\forall a \in Z_p^*\ \exists i$ such that $a = g^i \bmod p$
    - <u>Modular arithmetic</u>: numbers "wrap around" after they reach p



Compute $k=(g^y)^x=g^{xy} \bmod p$ Compute $k=(g^x)^y=g^{xy} \bmod p$

Why Is Diffie-Hellman Secure?

- Discrete Logarithm (DL) problem:
  given $g^x \bmod p$, it's hard to extract x
  - There is no known <u>efficient</u> algorithm for doing this
  - This is <u>not</u> enough for Diffie-Hellman to be secure!
- Computational Diffie-Hellman (CDH) problem:
  given g^x and g^y, it's hard to compute $g^{xy} \bmod p$
  - … unless you know x or y, in which case it's easy
- Decisional Diffie-Hellman (DDH) problem:
  given g^x and g^y, it's hard to tell the difference between $g^{xy} \bmod p$ and $g^r \bmod p$ where r is random

Properties of Diffie-Hellman

- Assuming DDH problem is hard, Diffie-Hellman protocol is a secure key establishment protocol against <u>passive</u> attackers
  - Eavesdropper can't tell the difference between established key and a random value
  - Can use new key for symmetric cryptography
    - Approx. 1000 times faster than modular exponentiation
- Diffie-Hellman protocol (by itself) does not provide authentication