CSE 484 (Winter 2008)

Software Security: Attacks, Defenses, and Design Principles

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Thanks to Dan Boneh, Dieter Gollmann, John Manferdelli, John Mitchell, Vitaly Shmatikov, Bennet Yee, and many others for sample slides and materials ...

### Goals for Today

- ♦ TOCTOU
- Integer Overflow, Casting
- Randomness
- Timing Attacks
- Defensive Mechanisms
- ◆ Software Development Design Principles

#### ΤΟCΤΟυ

- ◆ TOCTOU == Time of Check to Time of Use
- int openfile(char \*path) {
   struct stat s;
- if (stat(path, &s) < 0)
   return -1;</pre>
- return -1;
  if (!S\_ISRREG(s.st\_mode)) {
   error("only allowed to regular files!");
   return -1;

return open(path, O\_RDONLY);

- Goal: Open only regular files (not symlink, etc)
- Attacker can change meaning of path between stat and open (and access files he or she shouldn't)

### Integer Overflow and Implicit Cast

char buf[80]; void vulnerable() { int len = read int from network(); char \*p = read string from network(); if (len > sizeof buf) { error("length too large, nice try!"); }; return; memcpy(buf, p, len); }

void \*memcpy(void \*dst, const void \* src, size\_t n); typedef unsigned int size\_t;

◆ If len is negative, may copy huge amounts of input into buf (from www-inst.eecs.berkeley.edu-implflaws.pdf)

#### Integer Overflow and Implicit Cast

size\_t len = read\_int\_from\_network();
char \*buf;
buf = malloc(len+5);
read(fd, buf, len);

- What if len is large (e.g., len = 0xFFFFFFF)?
- Then len + 5 = 4 (on many platforms)
- Result: Allocate a 4-byte buffer, then read a lot of data into that buffer.

(from www-inst.eecs.berkeley.edu-implflaws.pdf)

#### Randomness issues

- Many applications (especially security ones) require randomness
- "Obvious" uses:
- Generate secret cryptographic keysGenerate random initialization vectors for encryption
- Other "non-obvious" uses:
- Generate passwords for new users
- Shuffle the order of votes (in an electronic voting machine)
- Shuffle cards (for an online gambling site)

## C's rand() Function

#### ◆ C has a built-in random function: rand()

- unsigned long int next = 1; /\* rand: return pseudo-random integer on 0..32767 \*/
- int rand(void) {
   next \* 1103515245 + 12345;
   return (unsigned int)(next/65536) % 32768;
- }
  /\* srand: set seed for rand() \*/
  void srand(unsigned int seed) {
   next = seed;

next = s

# Problem: don't use rand() for security-critical applications!

 Given a few sample outputs, you can predict subsequent ones

#### Dr.Dobb's Portal

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NEWSLETTERS   RESOL	IRCES   BLOGS   PODCAST	S   CAREERS

RINT ISSUE

Email Discuss

#### Windows/.NET

Randomness	and	the	Netscape	
Browser			-	

How secure is the World Wide Web?

Ian Goldberg and David Wagner

No one was more surprised than Netscape Communications when a pair of computer-science students broke the Netscape encryption scheme. Ian and David describe how they attacked the popular Web browser and what they found out.

# Problems in Practice

- One institution used (something like) rand() to generate passwords for new users
- Given your password, you could predict the passwords of other users
- ◆ Kerberos (1988 1996)
- Random number generator improperly seeded
  Possible to trivially break into machines that rely upon Kerberos for authentication
- Online gambling websites
- Random numbers to shuffle cards
- Real money at stake
- But what if poor choice of random numbers?









Big news... CNN, etc..

#### **Obtaining Pseudorandom Numbers**

- For security applications, want "cryptographically secure pseudorandom numbers"
- Libraries include:
  - OpenSSL
- CryptoAPI (Microsoft)
- Linux:
  - /dev/random
  - /dev/urandom
- ◆ Internally:
- Pool from multiple sources (interrupt timers, keyboard, ...)
- Physical sources (radioactive decay, ...)

#### Timing Attacks

- Assume there are no "typical" bugs in the software
- No buffer overflow bugs
- No format string vulnerabilities
- Good choice of randomness
- Good design
- The software may still be vulnerable to timing attacks
- Software exhibits input-dependent timings
- Complex and hard to fully protect against

# Password Checker

- Functional requirements
- PwdCheck(RealPwd, CandidatePwd) should:
   Return TRUE if RealPwd matches CandidatePwd
- Return FALSE otherwise
- RealPwd and CandidatePwd are both 8 characters long
- Implementation (like TENEX system)
  - PwdCheck(RealPwd, CandidatePwd) // both 8 chars for i = 1 to 8 do
    - if (RealPwd[i] != CandidatePwd[i]) then
    - return FALSE
    - return TRUE
- Clearly meets functional description

#### Attacker Model

- PwdCheck(RealPwd, CandidatePwd) // both 8 chars for i = 1 to 8 do if (RealPwd[i] != CandidatePwd[i]) then return FALSE
- return TRUE
- Attacker can guess CandidatePwds through some standard interface
- ◆ Naive: Try all 256<sup>8</sup> = 18,446,744,073,709,551,616 possibilities
- Better: Time how long it takes to reject a CandidatePasswd. Then try all possibilities for first character, then second, then third, .... • Total tries: 256\*8 = 2048

#### Other Examples

- Plenty of other examples of timings attacks
  - AES cache misses
  - AES is the "Advanced Encryption Standard" - It is used in SSH, SSL, IPsec, PGP, ...
  - RSA exponentiation time
  - RSA is a famous public-key encryption scheme
     It's also used in many cryptographic protocols and products