CSE 473: Artificial Intelligence

Informed Search

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[These slides were adapted from Dan Klein and Pieter Abbeel for CS188 Intro to AI at UC Berkeley. All CS188 materials are available at http://ai.berkeley.edu]
Today

- Informed Search
  - Heuristics
  - Greedy Search
  - A* Search

- Graph Search
Recap: Search

- **Search problem:**
  - States (configurations of the world)
  - Actions and costs
  - Successor function (world dynamics)
  - Start state and goal test

- **Search tree:**
  - Nodes: represent plans for reaching states
  - Plans have costs (sum of action costs)

- **Search algorithm:**
  - Systematically builds a search tree
  - Chooses an ordering of the fringe (unexplored nodes)
  - Optimal: finds least-cost plans
Example: Pancake Problem

Cost: Number of pancakes flipped
BOUNDS FOR SORTING BY PREFIX REVERSAL

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For a permutation $\sigma$ of the integers from 1 to $n$, let $f(\sigma)$ be the smallest number of prefix reversals that will transform $\sigma$ to the identity permutation, and let $f(n)$ be the largest such $f(\sigma)$ for all $\sigma$ in (the symmetric group) $S_n$. We show that $f(n) \leq (5n + 5)/3$, and that $f(n) \geq 17n/16$ for $n$ a multiple of 16. If, furthermore, each integer is required to participate in an even number of reversed prefixes, the corresponding function $g(n)$ is shown to obey $3n/2 - 1 \leq g(n) \leq 2n + 3$. 
State space graph with costs as weights
function Tree-Search(problem, strategy) returns a solution, or failure
initialize the search tree using the initial state of problem
loop do
  if there are no candidates for expansion then return failure
  choose a leaf node for expansion according to strategy
  if the node contains a goal state then return the corresponding solution
  else expand the node and add the resulting nodes to the search tree
end

Action: flip top two
Cost: 2

Path to reach goal:
Flip four, flip three
Total cost: 7
All these search algorithms are the same except for fringe strategies

- Conceptually, all fringes are priority queues (i.e. collections of nodes with attached priorities)
- Practically, for DFS and BFS, you can avoid the $\log(n)$ overhead from an actual priority queue, by using stacks and queues
- Can even code one implementation that takes a variable queuing object
Uninformed Search
Uniform Cost Search

- Strategy: expand lowest path cost

- The good: UCS is complete and optimal!

- The bad:
  - Explores options in every “direction”
  - No information about goal location

[Demo: contours UCS empty (L3D1)]
[Demo: contours UCS pacman small maze (L3D3)]
Video of Demo Contours UCS Empty
Video of Demo Contours UCS Pacman Small Maze

SCORE: 0
Informed Search
A heuristic is:
- A function that estimates how close a state is to a goal
- Designed for a particular search problem
- Examples: Manhattan, Euclidean distance for pathing
Example: Heuristic Function

h(x)

Straight-line distance to Bucharest
- Arad: 366
- Bucharest: 0
- Craiova: 160
- Dobrogea: 242
- Eforie: 161
- Fagaras: 178
- Giurgiu: 77
- Hirsova: 151
- Iasi: 226
- Lugoj: 244
- Mehadi: 241
- Neamt: 234
- Oradea: 380
- Pitești: 98
- Rimnicu Vâlce: 193
- Sibiu: 253
- Timișoara: 329
- Urziceni: 80
- Vaslui: 199
- Zerind: 374
Example: Heuristic Function

Heuristic: the number of the largest pancake that is still out of place
Greedy Search
Example: Heuristic Function

h(x)
Greedy Search

- Expand the node that seems closest...

- What can go wrong?
Greedy Search

- **Strategy:** expand a node that you think is closest to a goal state
  - Heuristic: estimate of distance to nearest goal for each state

- **A common case:**
  - Best-first takes you straight to the (wrong) goal

- **Worst-case:** like a badly-guided DFS

[Demo: contours greedy empty (L3D1)]
[Demo: contours greedy pacman small maze (L3D4)]
Video of Demo Contours Greedy (Pacman Small Maze)
A* Search
A* Search
Combining UCS and Greedy

- **Uniform-cost** orders by path cost, or *backward cost* $g(n)$
- **Greedy** orders by goal proximity, or *forward cost* $h(n)$

- **A* Search** orders by the sum: $f(n) = g(n) + h(n)$

Example: Teg Grenager
When should A* terminate?

- Should we stop when we enqueue a goal?
  - No: only stop when we dequeue a goal
What went wrong?
- Actual bad goal cost < estimated good goal cost
- We need estimates to be less than actual costs!
Admissible Heuristics
Idea: Admissibility

Inadmissible (pessimistic) heuristics break optimality by trapping good plans on the fringe

Admissible (optimistic) heuristics slow down bad plans but never outweigh true costs
A heuristic $h$ is **admissible** (optimistic) if:

$$0 \leq h(n) \leq h^*(n)$$

where $h^*(n)$ is the true cost to a nearest goal.

Examples:

Coming up with admissible heuristics is most of what’s involved in using A* in practice.
Optimality of A* Tree Search
Assume:
- A is an optimal goal node
- B is a suboptimal goal node
- h is admissible

Claim:
- A will exit the fringe before B
Proof:

- Imagine B is on the fringe
- Some ancestor \( n \) of A is on the fringe, too (maybe A!)
- Claim: \( n \) will be expanded before B
  1. \( f(n) \) is less or equal to \( f(A) \)

\[
\begin{align*}
  f(n) &= g(n) + h(n) & \text{Definition of f-cost} \\
  f(n) &\leq g(A) & \text{Admissibility of } h \\
  g(A) &= f(A) & \text{h = 0 at a goal}
\end{align*}
\]
Optimality of A* Tree Search: Blocking

Proof:
- Imagine B is on the fringe
- Some ancestor $n$ of A is on the fringe, too (maybe A!)
- Claim: $n$ will be expanded before B
  1. $f(n)$ is less or equal to $f(A)$
  2. $f(A)$ is less than $f(B)$

\[ g(A) < g(B) \quad \text{B is suboptimal} \]
\[ f(A) < f(B) \quad h = 0 \text{ at a goal} \]
Optimality of A* Tree Search: Blocking

Proof:

- Imagine B is on the fringe
- Some ancestor $n$ of A is on the fringe, too (maybe A!)
- Claim: $n$ will be expanded before B
  1. $f(n)$ is less or equal to $f(A)$
  2. $f(A)$ is less than $f(B)$
  3. $n$ expands before B
- All ancestors of A expand before B
- A expands before B
- A* search is optimal
Properties of A*
Properties of A*

Uniform-Cost

A*
UCS vs A* Contours

- Uniform-cost expands equally in all "directions"

- A* expands mainly toward the goal, but does hedge its bets to ensure optimality

[Demo: contours UCS / greedy / A* empty (L3D1)]
[Demo: contours A* pacman small maze (L3D5)]
Video of Demo Contours (Empty) -- UCS
Video of Demo Contours (Empty) -- Greedy
Video of Demo Contours (Empty) – A*
Video of Demo Contours (Pacman Small Maze) – A*
Comparison

Greedy	Uniform Cost	A*
A* Applications

- Video games
- Pathing / routing problems
- Resource planning problems
- Robot motion planning
- Language analysis
- Machine translation
- Speech recognition
- ...

[Demo: UCS / A* pacman tiny maze (L3D6,L3D7)]
[Demo: guess algorithm Empty Shallow/Deep (L3D8)]
Video of Demo Pacman (Tiny Maze) – UCS / A*
Video of Demo Empty Water Shallow/Deep – Guess Algorithm
Creating Heuristics
Creating Admissible Heuristics

- Most of the work in solving hard search problems optimally is in coming up with admissible heuristics.
- Often, admissible heuristics are solutions to relaxed problems, where new actions are available.
- Inadmissible heuristics are often useful too.
Example: 8 Puzzle

- What are the states?
- How many states?
- What are the actions?
- How many successors from the start state?
- What should the costs be?
8 Puzzle I

- Heuristic: Number of tiles misplaced
- Why is it admissible?
- $h(\text{start}) = 8$
- This is a *relaxed-problem* heuristic

**Start State**

**Goal State**

<table>
<thead>
<tr>
<th>Average nodes expanded when the optimal path has...</th>
<th>...4 steps</th>
<th>...8 steps</th>
<th>...12 steps</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>UCS</strong></td>
<td>112</td>
<td>6,300</td>
<td>$3.6 \times 10^6$</td>
</tr>
<tr>
<td><strong>TILES</strong></td>
<td>13</td>
<td>39</td>
<td>227</td>
</tr>
</tbody>
</table>

Statistics from Andrew Moore
What if we had an easier 8-puzzle where any tile could slide any direction at any time, ignoring other tiles?

Total Manhattan distance

Why is it admissible?

\[ h(\text{start}) = 3 + 1 + 2 + ... = 18 \]

<table>
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</thead>
<tbody>
<tr>
<td>TILES</td>
<td>13</td>
<td>39</td>
<td>227</td>
</tr>
<tr>
<td>MANHATTAN</td>
<td>12</td>
<td>25</td>
<td>73</td>
</tr>
</tbody>
</table>
How about using the *actual cost* as a heuristic?
- Would it be admissible?
- Would we save on nodes expanded?
- What’s wrong with it?

With A*: a trade-off between quality of estimate and work per node
- As heuristics get closer to the true cost, you will expand fewer nodes but usually do more work per node to compute the heuristic itself
Semi-Lattice of Heuristics
Trivial Heuristics, Dominance

- **Dominance:**\( h_a \geq h_c \) if
  \[
  \forall n : h_a(n) \geq h_c(n)
  \]

- **Heuristics form a semi-lattice:**
  - Max of admissible heuristics is admissible
    \[
    h(n) = \max(h_a(n), h_b(n))
    \]

- **Trivial heuristics**
  - Bottom of lattice is the zero heuristic (what does this give us?)
  - Top of lattice is the exact heuristic
Graph Search
Failure to detect repeated states can cause exponentially more work.
In BFS, for example, we shouldn’t bother expanding the circled nodes (why?)
Graph Search

- **Idea:** never expand a state twice
- **How to implement:**
  - Tree search + set of expanded states ("closed set")
  - Expand the search tree node-by-node, but...
  - Before expanding a node, check to make sure its state has never been expanded before
    - If not new, skip it, if new add to closed set
- **Important:** store the closed set as a set, not a list
- **Can graph search wreck completeness?** Why/why not?
- **How about optimality?**
A* Graph Search Gone Wrong?

State space graph

Search tree

S (0+2)

A (1+4)  B (1+1)

C (2+1)  C (3+1)

G (5+0)  G (6+0)
Consistency of Heuristics

- **Main idea:** heuristic costs ≤ actual costs
  - **Admissibility:** heuristic cost ≤ actual cost to goal
    \[ h(A) \leq \text{actual cost from A to G} \]
  - **Consistency:** heuristic “arc” cost ≤ actual cost for each arc
    \[ h(A) - h(C) \leq \text{cost}(A \text{ to } C) \]

- **Consequences of consistency:**
  - The f value along a path never decreases
    \[ h(A) \leq \text{cost}(A \text{ to } C) + h(C) \]
  - A* graph search is optimal
Optimality of A* Graph Search
Optimality of A* Graph Search

Sketch: consider what A* does with a consistent heuristic:

- Fact 1: In tree search, A* expands nodes in increasing total f value (f-contours)
- Fact 2: For every state s, paths that reach s optimally are expanded before paths that reach s suboptimally
- Result: A* graph search is optimal
Optimality

- **Tree search:**
  - A* is optimal if heuristic is admissible
  - UCS is a special case ($h = 0$)

- **Graph search:**
  - A* optimal if heuristic is consistent
  - UCS optimal ($h = 0$ is consistent)

- Consistency implies admissibility

- In general, most natural admissible heuristics tend to be consistent, especially if from relaxed problems
A*: Summary
A*: Summary

- A* uses both backward costs and (estimates of) forward costs
- A* is optimal with admissible / consistent heuristics
- Heuristic design is key: often use relaxed problems
function Tree-Search(problem, fringe) return a solution, or failure
    fringe ← Insert(make-node(initial-state[problem]), fringe)
    loop do
        if fringe is empty then return failure
        node ← REMOVE-FRONT(fringe)
        if GOAL-TEST(problem, STATE[node]) then return node
        for child-node in EXPAND(STATE[node], problem) do
            fringe ← Insert(child-node, fringe)
        end
    end
Graph Search Pseudo-Code

function \textsc{Graph-Search}(\textit{problem}, fringe) return a solution, or failure

\begin{itemize}
  \item \texttt{closed} $\leftarrow$ an empty set
  \item \texttt{fringe} $\leftarrow$ \textsc{Insert}(\textsc{Make-Node}(\textsc{Initial-State}[\textit{problem}]), fringe)
\end{itemize}

\textbf{loop} do

  \begin{itemize}
  \item \textbf{if} \texttt{fringe} is empty \textbf{then} \textbf{return} failure
  \item \texttt{node} $\leftarrow$ \textsc{Remove-Front}(fringe)
  \item \textbf{if \textsc{Goal-Test}(\textit{problem}, \textsc{State}[\texttt{node}]) then \textbf{return} node}
  \item \textbf{if \textsc{State}[\texttt{node}] is not in \texttt{closed} then}
  \begin{itemize}
    \item add \textsc{State}[\texttt{node}] to \texttt{closed}
    \item \textbf{for \texttt{child-node} in \textsc{Expand}(\textsc{State}[\texttt{node}], \textit{problem}) do}
    \begin{itemize}
      \item \texttt{fringe} $\leftarrow$ \textsc{Insert}(\texttt{child-node}, fringe)
    \end{itemize}
  \end{itemize}
  \end{itemize}

\textbf{end}