### **Issues in Multiprocessors**

### Which programming model for interprocessor communication

- · shared memory
  - · regular loads & stores
- · message passing
  - explicit sends & receives

#### Which execution model

- · control parallel
  - identify & synchronize different asynchronous threads
- data parallel
  - same operation on different parts of the shared data space

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## **Issues in Multiprocessors**

#### How to express parallelism

- automatic (compiler) detection
  - implicitly parallel C & Fortran programs, e.g., SUIF compiler
- · language support
  - HPF, ZPL
- · runtime library constructs
  - · coarse-grain, explicitly parallel C programs

### **Algorithm development**

- · embarrassingly parallel programs could be easily parallelized
- · development of different algorithms for same problem

### **Issues in Multiprocessors**

#### How to get good parallel performance

- recognize parallelism
- transform programs to increase parallelism without decreasing processor locality
- · decrease sharing costs

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## **Flynn Classification**

SISD: single instruction stream, single data stream

• single-context uniprocessors

SIMD: single instruction stream, multiple data streams

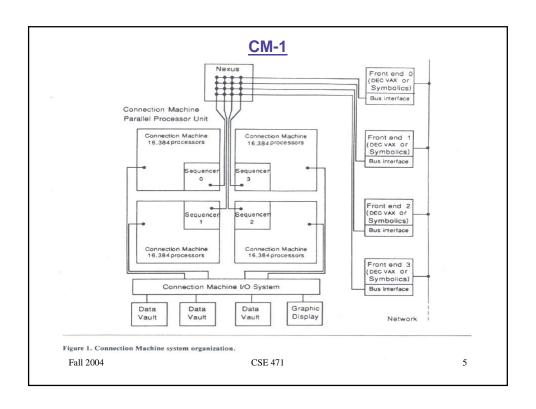
- · exploits data parallelism
- example: Thinking Machines CM-1...

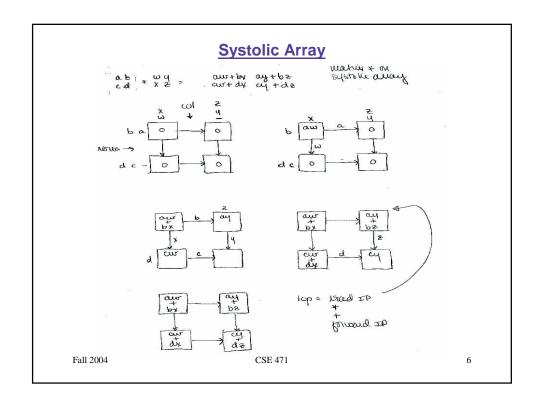
MISD: multiple instruction streams, single data stream

- · systolic arrays
- · example: Intel iWarp

MIMD: multiple instruction streams, multiple data streams

- · multiprocessors
- · multithreaded processors
- relies on control parallelism: execute & synchronize different asynchronous threads of control
- parallel programming & multiprogramming
- example: most processor companies have MP configurations

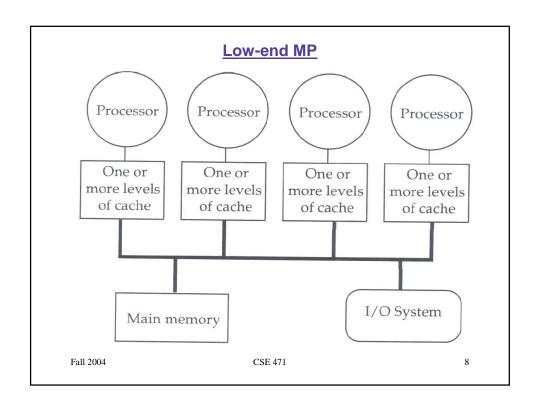




## **MIMD**

#### Low-end

- bus-based
  - simple, but a bottleneck
  - simple cache coherency protocol
- · physically centralized memory
- uniform memory access (UMA machine)
- Sequent Symmetry, SPARCCenter, Alpha-, PowerPC- or SPARC-based servers

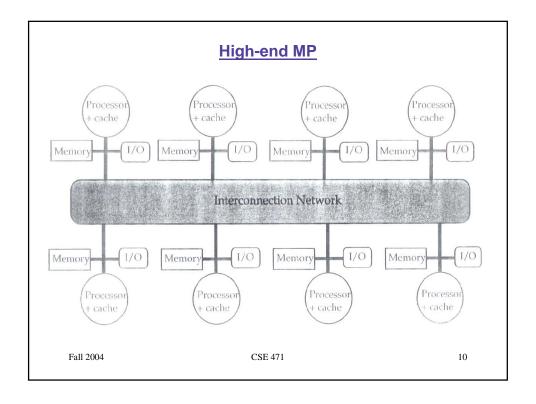


## **MIMD**

### **High-end**

- higher bandwidth, multiple-path interconnect
  - more scalable
  - more complex cache coherency protocol (if shared memory)
  - · longer latencies
- · physically distributed memory
- non-uniform memory access (NUMA machine)
- could have processor clusters
- SGI Challenge, Convex Examplar, Cray T3D, IBM SP-2, Intel Paragon

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## **MIMD Programming Models**

#### Address space organization for physically distributed memory

- · distributed shared memory
  - 1 (logical) global address space
- multicomputers
  - private address space/processor

#### Inter-processor communication

- · shared memory
  - · accessed via load/store instructions
  - SPARCCenter, SGI Challenge, Cray T3D, Convex Exemplar, KSR-1&2
- · message passing
  - · explicit communication by sending/receiving messages
  - TMC CM-5, Intel Paragon, IBM SP-2

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### **Shared Memory vs. Message Passing**

### **Shared memory**

- + simple parallel programming model
  - global shared address space
  - not worry about data locality but

get better performance when program for data placement lower latency when data is local

less communication when avoid false sharing

- but can do data placement if it is crucial, but don't have to
- · hardware maintains data coherence
- synchronize to order processor's accesses to shared data
- like uniprocessor code so parallelizing by programmer or compiler is easier
  - $\Rightarrow$  can focus on program semantics, not interprocessor communication

### **Shared Memory vs. Message Passing**

#### **Shared memory**

- + low latency (no message passing software) but overlap of communication & computation latency-hiding techniques can be applied to message passing machines
- + higher bandwidth for small transfers but usually the only choice

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### **Shared Memory vs. Message Passing**

### Message passing

- + abstraction in the programming model encapsulates the communication costs *but* 
  - more complex programming model additional language constructs
  - need to program for nearest neighbor communication
- + no coherency hardware
- + good throughput on large transfers **but** what about small transfers?
- more scalable (memory latency doesn't scale with the number of processors) but

large-scale SM has distributed memory also

 hah! so you're going to adopt the message-passing model?

# **Shared Memory vs. Message Passing**

Why there was a debate

- little experimental data
- not separate implementation from programming model
- can emulate one paradigm with the other
  - MP on SM machine message buffers in local (to each processor) memory copy messages by ld/st between buffers
  - SM on MP machine Id/st becomes a message copy slooooooooow

Who won?