What's Interesting about RPC?

• RPC = Remote Procedure Call
  – the most common means for remote communication
  – used both by operating systems and applications
  • NFS is implemented as a set of RPCs
  • HTTP is essentially RPC
  • DCOM, CORBA, Java RMI, etc., are just RPC systems
• Allows you to communicate over a network with syntax and semantics very similar to local procedure call

Client/Server communication

• The prevalent model for structuring distributed computation is the client/server paradigm
  – a server is a program (or collection of programs) that provides a service to other programs
    • e.g., file server, name server, web server, mail server …
    • servers/services may span multiple nodes (clusters)
    – often, nodes are called servers too
    e.g., the web server runs on a Dell server computer
  – a client is a program that uses the service
    • the client first binds to the server
      – locates it, establishes a network connection to it
    • the client then sends requests (with data) to perform actions, and the server sends responses (with data)
      – e.g., web browser sends a "GET" request, server responds with a web page
    TCP/IP is the transport, but what is the higher-level programming model?

Procedure calls

• Procedure calls are a natural way to structure multiple modules inside a single program
  – every language supports procedure calls
  – semantics are well-defined and well-understood
  – programmers are used to them
• "Server" (called procedure) exports an API
  – think about a file system / file server API: open, close, read, write, sync, etc.
• "Client" (calling procedure) calls the server procedure's API
• Linker binds the two together

Procedure call example

• If the server were just a library, then "Add" would just be a local procedure call
Remote Procedure Call

- Use procedure calls as the model for distributed (remote) communication
  - traditional procedure call syntax and semantics
  - have servers export a set of procedures that can be called by client programs
    - similar to library API, class definitions, etc.
    - clients do a local procedure call, as though they were directly linked with the server
      - under the covers, the procedure call is converted into a message exchange with the server
    - largely invisible to the programmer!

Remote Procedure Call

- There are a bunch of hard issues:
  - how do we make the “remote” part of RPC invisible to the programmer?
    - and is that a good idea?
  - what are the semantics of parameter passing?
    - what if we try to pass by reference?
  - how do we bind (locate/connect-to) servers?
  - how do we handle heterogeneity?
    - OS, language, architecture, …
  - how do we make it go fast?

Procedure Call

RPC model

- A server defines the service interface using an interface definition language (IDL)
  - the IDL specifies the names, parameters, and types for all client-callable server procedures
    - example: ASN.1 in the OSI reference model
    - example: Sun’s XDR (external data representation)
- A “stub compiler” reads the IDL declarations and produces two stub procedures for each server procedure
  - the server programmer implements the service’s procedures and links them with the server-side stubs
  - the client programmer implements the client program and links it with the client-side stubs
  - the stubs manage all of the details of remote communication between client and server using the RPC runtime system

RPC stubs

- A client-side stub is a procedure that looks to the client as if it were a callable server procedure
  - it has the same API as the server’s implementation of the procedure
  - a client-side stub is just called a “stub” in Java RMI
- A server-side stub looks like a caller to the server
  - it looks like a hunk of code that invokes the server procedure
  - a server-side stub is called a “skeleton” or “skel” in Java RMI
  - The client program thinks it’s invoking the server
    - but it’s calling into the client-side stub
  - The server program thinks it’s called by the client
    - but it’s really called by the server-side stub
  - The stubs send messages to each other, via the runtime, to make the RPC happen transparently

Remote Procedure Call

- Client-side stub:
  - int Add(int x, int y) {
    - alloc message buffer;
    - mark as “add” call;
    - store x,y in buffer;
    - send message;
    - receive response;
    - unpack response;
    - return response;
  }
- Server-side stub:
  - Message
    - Add_Stub(Message m) {
      - remove x,y from m;
      - r = Add(x,y);
      - allocate response buffer;
      - store r in response;
      - return response;
    }
- RPC runtime system:
  - receive message m;
  - response = Add_Stub(m);
  - send response to client;

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  - send response to client;
Remote Procedure Call

Client Program:
```
sum = server->Add(3, 4);
```

Server Program:
```
int Add(int x, int y) {
    return x + y;
}
```

Client-side stub:
```
int Add(int x, int y) {
    alloc message buffer;
    mark as "add" call;
    store x, y in buffer;
    send message;
    receive response;
    unpack response;
    return response;
}
```

RPC runtime system:
```
send message to server;
receive response;
```

Server-side stub:
```
Message Add_Stub(Message m) {
    remove x, y from m;
    r = Add(x, y);
    allocate response buffer;
    store r in response;
    return response;
}
```

RPC runtime system:
```
receive message m;
response = Add_Stub(m);
send response to client;
```

Topics:
- interface description
- stubs
- stub generation
- parameter marshalling
- binding
- runtime system
- error handling
- performance
- thread pools

RPC marshalling

- Marshalling is the packing of procedure parameters into a message packet
  - the RPC stubs call type-specific procedures to marshal or unmarshal the parameters of an RPC
    - the client stub marshals the parameters into a message
    - the server stub unmarshals the parameters and uses them to invoke the service's procedure
  - on return:
    - the server stub marshals the return value
    - the client stub unmarshals the return value, and returns them to the client program

RPC binding

- Binding is the process of connecting the client to the server
  - the server, when it starts up, exports its interface
    - identifies itself to a network name server
    - tells RPC runtime that it is alive and ready to accept calls
  - the client, before issuing any calls, imports (binds to) the server
    - RPC runtime uses the name server to find the location of the server and establish a connection
  - The import and export operations are explicit in the server and client programs
    - a slight breakdown in transparency
      - more to come...

RPC transparency

- One goal of RPC is to be as transparent as possible
  - make remote procedure calls look like local procedure calls
    - we've seen that binding breaks this transparency
  - What else breaks transparency?
    - failures: remote nodes/networks can fail in more ways than with local procedure calls
      - network partition, server crash
    - need extra support to handle failures
    - server can fail independently from client
      - "partial failure": a big issue in distributed systems
        - if an RPC fails, was it invoked on the server?
        - performance: remote communication is inherently slower than local communication

RPC and thread pools

- What happens if two client threads (or client programs) simultaneously invoke the same server using RPC?
  - ideally, two separate threads will run on the server
  - so, the RPC runtime system on the server needs to spawn or dispatch threads into server-side stubs when messages arrive
    - is there a limit on the number of threads?
      - if so, does this change semantics?
      - if not, what if 1,000,000 clients simultaneously RPC into the same server?