Scheduling

- In discussion process management, we talked about context switching between threads/process on the ready queue
  - but, we glossed over the details of which process or thread is chosen next
  - making this decision is called scheduling
    - scheduling is policy
    - context switching is mechanism

- Today, we’ll look at:
  - the goals of scheduling
    - starvation
  - well-known scheduling algorithms
    - standard UNIX scheduling
Multiprogramming and Scheduling

• Multiprogramming increases resource utilization and job throughput by overlapping I/O and CPU
  – today: look at scheduling policies
    • which process/thread to run, and for how long
  – schedulable entities are usually called jobs
    • processes, threads, people, disk arm movements, …

• There are two time scales of scheduling the CPU:
  – long term: determining the multiprogramming level
    • how many jobs are loaded into primary memory
    • act of loading in a new job (or loading one out) is swapping
  – short-term: which job to run next to result in “good service”
    • happens frequently, want to minimize context-switch overhead
    • good service could mean many things
Scheduling

• The **scheduler** is the module that moves jobs from queue to queue
  – the **scheduling algorithm** determines which job(s) are chosen to run next, and which queues they should wait on
  – the scheduler is typically run when:
    • a job switches from running to waiting
    • when an interrupt occurs
      – especially a timer interrupt
    • when a job is created or terminated

• There are two major classes of scheduling systems
  – in **preemptive** systems, the scheduler can interrupt a job and force a context switch
  – in non-**preemptive** systems, the scheduler waits for the running job to explicitly (voluntarily) block
Scheduling Goals

• Scheduling algorithms can have many different goals (which sometimes conflict)
  – maximize CPU utilization
  – maximize job throughput (#jobs/s)
  – minimize job turnaround time ($T_{\text{finish}} - T_{\text{start}}$)
  – minimize job waiting time ($\text{Avg}(T_{\text{wait}})$: average time spent on wait queue)
  – minimize response time ($\text{Avg}(T_{\text{resp}})$: average time spent on ready queue)

• Goals may depend on type of system
  – batch system: strive to maximize job throughput and minimize turnaround time
  – interactive systems: minimize response time of interactive jobs (such as editors or web browsers)
Scheduler Non-goals

- Schedulers typically try to prevent starvation
  - starvation occurs when a process is prevented from making progress, because another process has a resource it needs
- A poor scheduling policy can cause starvation
  - e.g., if a high-priority process always prevents a low-priority process from running on the CPU
- Synchronization can also cause starvation
  - we’ll see this in a future class
  - roughly, if somebody else always gets a lock I need, I can’t make progress
Algorithm #1: FCFS/FIFO

• First-come first-served (FCFS)
  – jobs are scheduled in the order that they arrive
  – “real-world” scheduling of people in lines
    • e.g. supermarket, bank tellers, MacDonalds, …
  – typically non-preemptive
    • no context switching at supermarket!
  – jobs treated equally, no starvation
    • except possibly for infinitely long jobs

• Sounds perfect!
  – what’s the problem?
FCFS picture

- Problems:
  - average response time and turnaround time can be large
    - e.g., small jobs waiting behind long ones
    - results in high turnaround time
  - may lead to poor overlap of I/O and CPU
Algorithm #2: SJF

- Shortest job first (SJF)
  - choose the job with the smallest expected CPU burst
  - can prove that this has optimal min. average waiting time
- Can be preemptive or non-preemptive
  - preemptive is called shortest remaining time first (SRTF)
- Sounds perfect!
  - what’s the problem here?
SJF Problem

- Problem: impossible to know size of future CPU burst
  - from your theory class, equivalent to the halting problem
  - can you make a reasonable guess?
    - yes, for instance looking at past as predictor of future
    - but, might lead to starvation in some cases!
Priority Scheduling

• Assign priorities to jobs
  – choose job with highest priority to run next
    • if tie, use another scheduling algorithm to break (e.g. FCFS)
    • to implement SJF, priority = expected length of CPU burst

• Abstractly modeled as multiple “priority queues”
  – put ready job on queue associated with its priority

• Sound perfect!
  – what’s wrong with this?
Priority Scheduling: problem

• The problem: starvation
  – if there is an endless supply of high priority jobs, no low-priority job will ever run

• Solution: “age” processes over time
  – increase priority as a function of wait time
  – decrease priority as a function of CPU time
  – many ugly heuristics have been explored in this space
Round Robin

• Round Robin scheduling (RR)
  – ready queue is treated as a circular FIFO queue
  – each job is given a time slice, called a quantum
    • job executes for duration of quantum, or until it blocks
    • time-division multiplexing (time-slicing)
  – great for timesharing
    • no starvation
    • can be preemptive or non-preemptive

• Sounds perfect!
  – what’s wrong with this?
RR problems

• Problems:
  – what do you set the quantum to be?
    • no setting is “correct”
      – if small, then context switch often, incurring high overhead
      – if large, then response time drops
  – treats all jobs equally
    • if I run 100 copies of SETI@home, it degrades your service
    • how can I fix this?
Combining algorithms

• Scheduling algorithms can be combined in practice
  – have multiple queues
  – pick a different algorithm for each queue
  – and maybe, move processes between queues
• Example: multi-level feedback queues (MLFQ)
  – multiple queues representing different job types
    • batch, interactive, system, CPU-bound, etc.
  – queues have priorities
    • schedule jobs within a queue using RR
  – jobs move between queues based on execution history
    • “feedback”: switch from CPU-bound to interactive behavior
• Pop-quiz:
  – is MLFQ starvation-free?
UNIX Scheduling

• Canonical scheduler uses a MLFQ
  – 3-4 classes spanning ~170 priority levels
    • timesharing: first 60 priorities
    • system: next 40 priorities
    • real-time: next 60 priorities
  – priority scheduling across queues, RR within
    • process with highest priority always run first
    • processes with same priority scheduled RR
  – processes dynamically change priority
    • increases over time if process blocks before end of quantum
    • decreases if process uses entire quantum

• Goals:
  – reward interactive behavior over CPU hogs
    • interactive jobs typically have short bursts of CPU