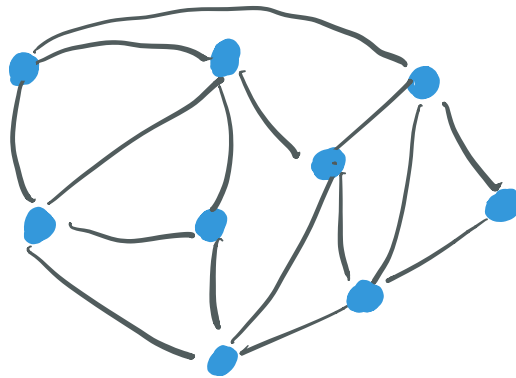


(1) A graph  $G = (V, E)$  is called planar if it can be drawn in the plane without edge crossings:



Define the language:

NONPLANAR

$$= \{ \langle G \rangle : G \text{ is an undirected graph that is } \underline{\text{not}} \text{ planar} \}$$

Use Kuratowski's theorem (see Wikipedia) to prove that  $\text{NONPLANAR} \in \text{NP}$

Extra credit: Can you show that

$$\text{PLANAR} = \{ \langle G \rangle : G \text{ is a planar graph} \}$$

is in NP? Use as simple a certificate as possible.

[Make sure the length of the certificate is polynomial in the input size!]

(2) Recall from class the language

$$3COL = \{ \langle G \rangle : G \text{ is a 3-colorable graph} \}$$

(a) Prove that  $3COL \leq_p 4COL$

(b) Consider the language:

$$TEN-3COL =$$

$$\{ \langle G \rangle : G \text{ is an undirected graph and has at least ten distinct proper 3-colorings} \}$$

Prove that  $3COL \leq_p TEN-3COL$

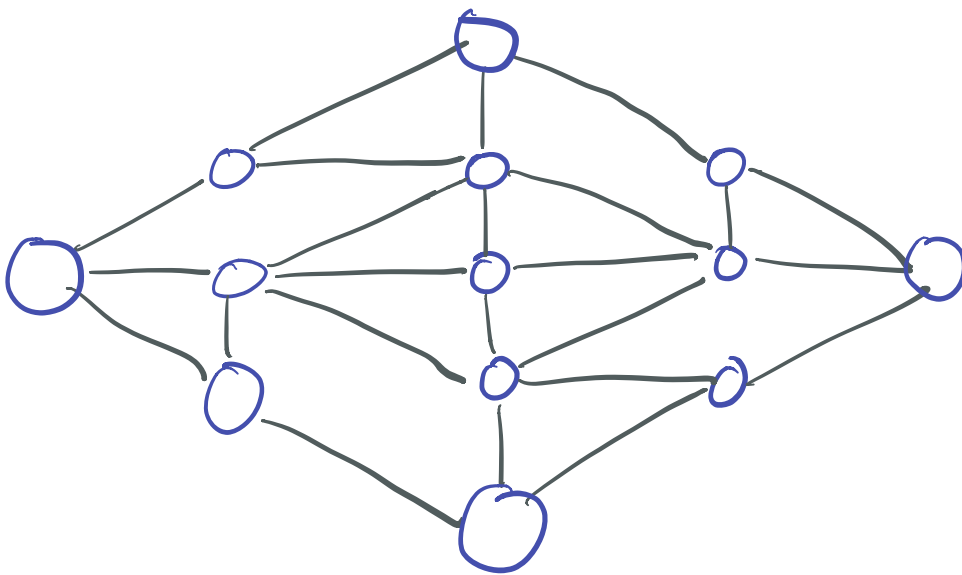
Extra credit: Consider the language:

PLANAR-3COL

=  $\{ \langle G \rangle : G \text{ is a planar, } 3\text{-colorable graph} \}$

Show that  $3COL \leq_p \text{PLANAR-3COL}$ .

Hint: Use the following gadget to replace edge crossings:



(3) Show that if 2COL is NP-complete, then  $P = NP$ .

Extra credit: Two undirected graphs

$H$  and  $G$  are isomorphic if the vertices of  $H$  can be reordered to obtain  $G$ .

Show that if  $P=NP$ , there is a polynomial-time algorithm that, given isomorphic graphs  $H$  and  $G$ , outputs the ordering.