All solutions should be neatly written or type set. All major steps in proofs and algorithms must be justified.

1. (10 points) In this problem you will get practice in doing a diagonal argument to show undecidability. Consider the language

\[ H_{TM} = \{ \langle M, w \rangle : M \text{ halts on input } w \}. \]

Use a diagonal argument to show that \( H_{TM} \) is undecidable.

2. (10 points) In this problem you will get practice in doing a reduction argument to show undecidability. Consider the language

\[ I_{TM} = \{ \langle M \rangle : \text{both } L(M) \text{ and its complement are infinite} \}. \]

Show that \( I_{TM} \) is undecidable by a reduction from \( A_{TM} \).

3. (15 points) Thue systems were invented in 1914. A Thue system consists of a finite set of rules \( \mathcal{T} \) of the form

\[ u \rightarrow v \]

where \( u, v \) are strings from a finite alphabet \( \Sigma \). One string can be derived in one step from another via \( \mathcal{T} \) using the following definition

\[ ux \Rightarrow xv \text{ if for some } i, u = u_i \text{ and } v = v_i. \]

For example, suppose the rules in \( \mathcal{T} \) are

\[ (0,0), (1,1), (\#,\#), (c1,0c), (e0,d1), (c\#,d0\#), (0d,d0), (1d,d1), (\#d,\#c) \]

then we could have the multiple step derivation

\[ \#c\# \Rightarrow c\#\# \Rightarrow \#d0\# \Rightarrow 0\#\#c \Rightarrow \#\#c0 \Rightarrow \#c0\# \Rightarrow \]

\[ e0\#\# \Rightarrow \#\#d1 \Rightarrow \#d1\# \Rightarrow 1\#\#c \Rightarrow \#\#c1 \Rightarrow \#c1\# \]

(a) Continue the derivation above for 10 more steps. Describe in words what the Thue system \( \mathcal{T} \) is doing.

(b) Show that the problem of determining if given Thue system \( \mathcal{T} \) and start string \( x \), \( x \) derives the empty string in \( \mathcal{T} \). Hint: Do a reduction from the acceptance problem for Turing machines. That is, given a Turing \( M \) and input \( w \) show how to construct a Thue system \( \mathcal{T} \) and start string \( x \) with the property that \( M \) accepts \( w \) if and only if \( x \) derives the empty string in \( \mathcal{T} \). There will be some resemblance to the construction of a general grammar equivalent to a given Turing machine.