# CSE 421: Introduction to Algorithms

Bipartiteness - DFS

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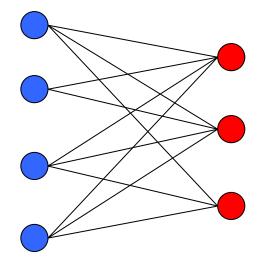
## **Bipartite Graphs**

Definition: An undirected graph G=(V,E) is bipartite if you can partition the node set into 2 parts (say, blue/red or left/right) so that

all edges join nodes in different parts i.e., no edge has both ends in the same part.

#### Application:

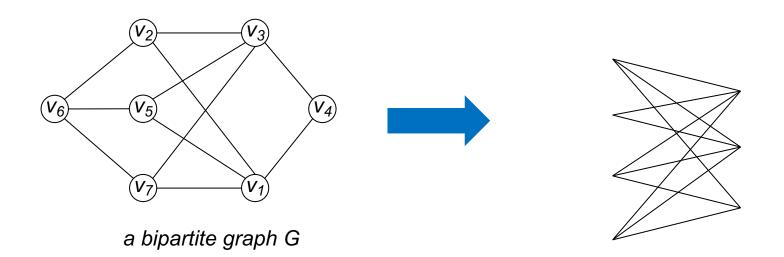
- Scheduling: machine=red, jobs=blue
- Stable Matching: men=blue, wom=red



a bipartite graph

## **Testing Bipartiteness**

Problem: Given a graph G, is it bipartite?



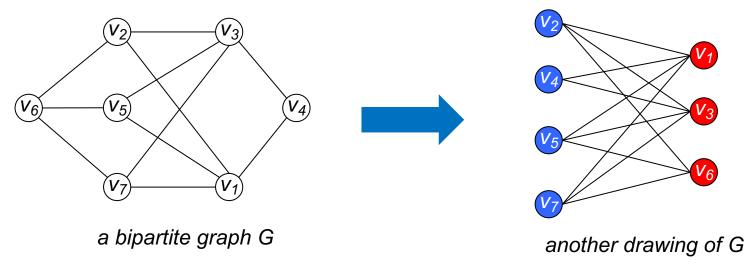
## **Testing Bipartiteness**

Problem: Given a graph G, is it bipartite?

Many graph problems become:

- Easier if the underlying graph is bipartite (matching)
- Tractable if the underlying graph is bipartite (independent set)

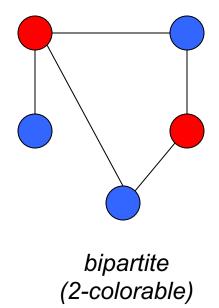
Before attempting to design an algorithm, we need to understand structure of bipartite graphs.

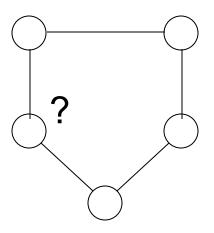


## An Obstruction to Bipartiteness

Lemma: If G is bipartite, then it does not contain an odd length cycle.

Pf: We cannot 2-color an odd cycle, let alone G.



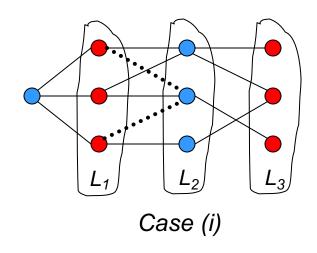


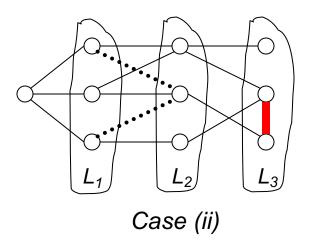
not bipartite (not 2-colorable)

## A Characterization of Bipartite Graphs

Lemma: Let G be a connected graph, and let  $L_0, ..., L_k$  be the layers produced by BFS(s). Exactly one of the following holds.

- (i) No edge of G joins two nodes of the same layer, and G is bipartite.
- (ii) An edge of G joins two nodes of the same layer, and G contains an odd-length cycle (and hence is not bipartite).





## A Characterization of Bipartite Graphs

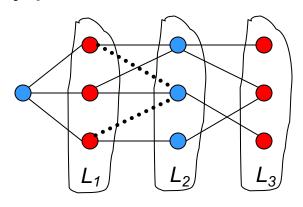
Lemma: Let G be a connected graph, and let  $L_0, ..., L_k$  be the layers produced by BFS(s). Exactly one of the following holds.

- (i) No edge of G joins two nodes of the same layer, and G is bipartite.
- (ii) An edge of G joins two nodes of the same layer, and G contains an odd-length cycle (and hence is not bipartite).

#### **Pf.** (i)

Suppose no edge joins two nodes in the same layer.

By previous lemma, all edges join nodes on adjacent levels.



Case (i)

#### Bipartition:

blue = nodes on odd levels,
red = nodes on even levels.

## A Characterization of Bipartite Graphs

Lemma: Let G be a connected graph, and let  $L_0, ..., L_k$  be the layers produced by BFS(s). Exactly one of the following holds.

- (i) No edge of G joins two nodes of the same layer, and G is bipartite.
- (ii) An edge of G joins two nodes of the same layer, and G contains an odd-length cycle (and hence is not bipartite).

#### Pf. (ii)

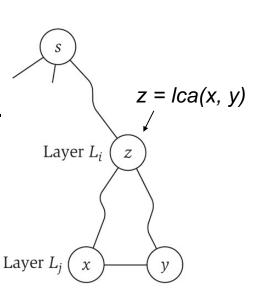
Suppose (x, y) is an edge & x, y in same level  $L_j$ .

Let z =their lowest common ancestor in BFS tree.

Let  $L_i$  be level containing z.

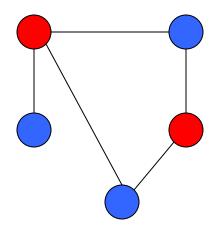
Consider cycle that takes edge from x to y, then tree from y to z, then tree from z to x.

Its length is 1 + (j-i) + (j-i), which is odd.

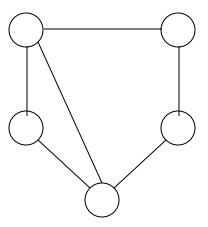


## Obstruction to Bipartiteness

Cor: A graph G is bipartite iff it contains no odd length cycles.



bipartite (2-colorable)



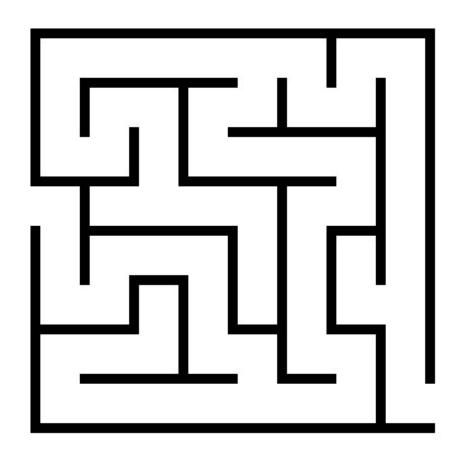
not bipartite (not 2-colorable)

### In class Exercise

Let G be a graph with n vertices and at least n edges. Show that G has a cycle.

## Depth First Search

Follow the first path you find as far as you can go; back up to last unexplored edge when you reach a dead end, then go as far you can



Naturally implemented using recursive calls or a stack

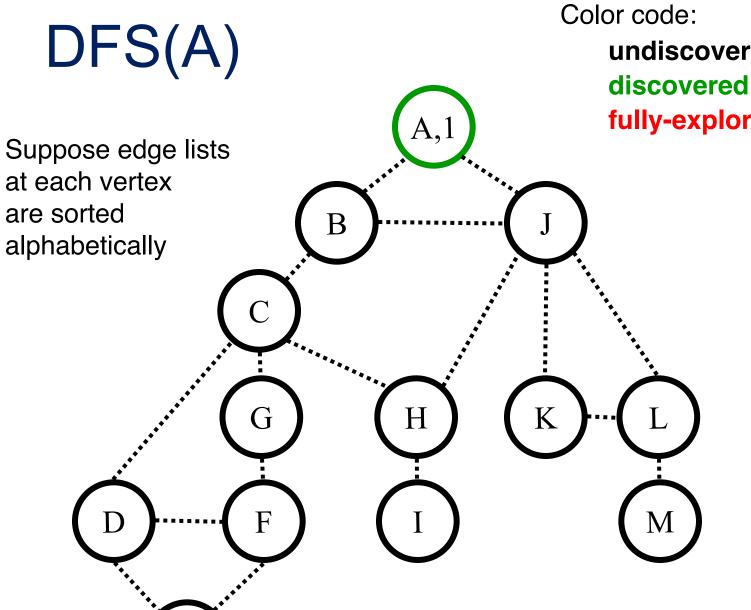
## DFS(s) – Recursive version

Global Initialization: mark all vertices undiscovered

```
DFS(v)
Mark v discovered

for each edge {v,x}
    if (x is undiscovered)
        Mark x discovered
        DFS(x)

Mark v full-discovered
```

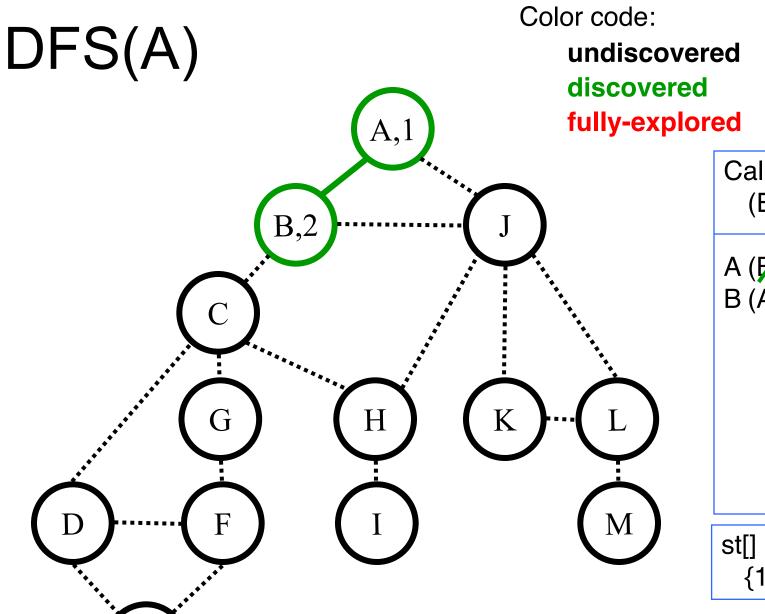


fully-explored

Call Stack (Edge list):

A (B,J)

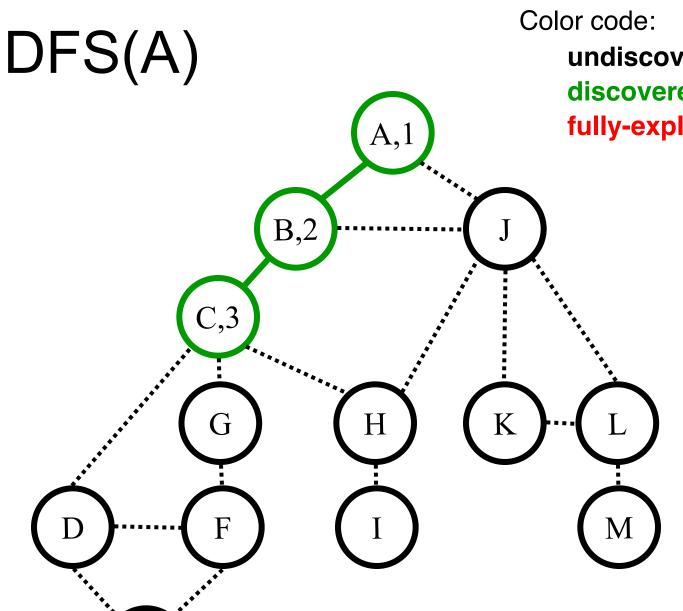
st[] = {1}



Call Stack: (Edge list)

A (**⅓**,J) B (A,C,J)

st[] = {1,2}



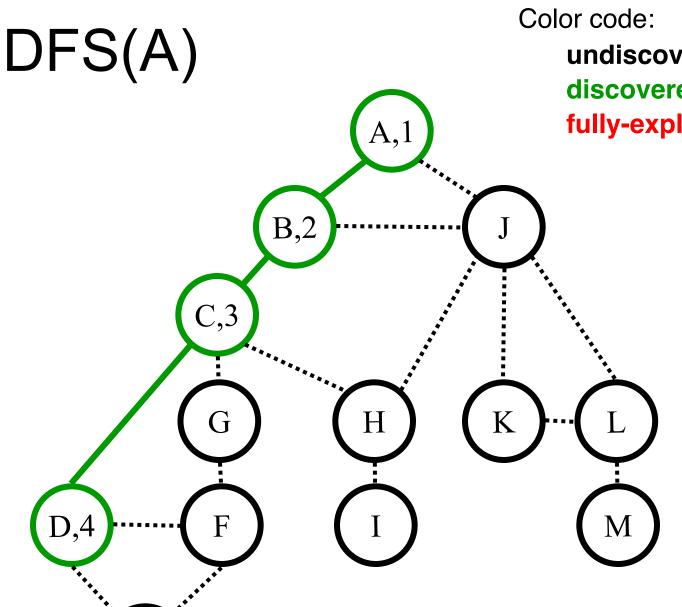
discovered

fully-explored

Call Stack: (Edge list)

A(P,J)B (**A**,**C**,J) C (B,D,G,H)

st[] = {1,2,3}



discovered

fully-explored

Call Stack: (Edge list)

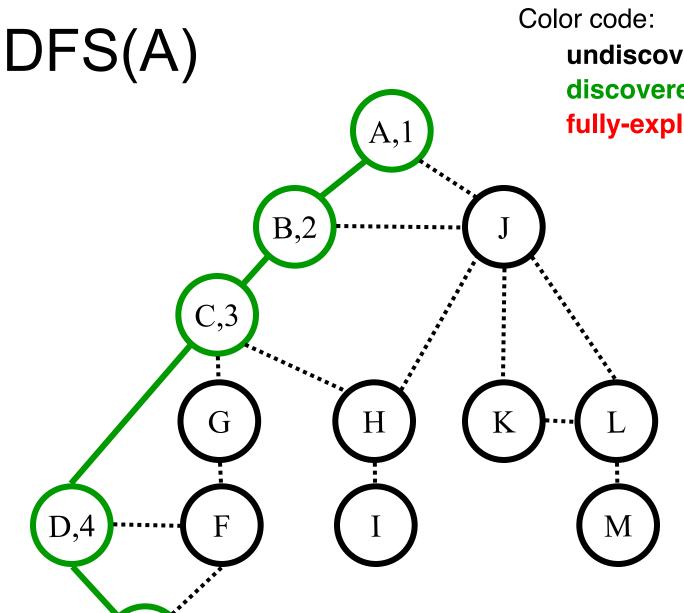
A(B,J)

 $B(\cancel{A},\cancel{C},J)$ 

C (**p**,**p**,G,H)

D(C,E,F)

st[] = {1,2,3,4}



discovered

fully-explored

Call Stack: (Edge list)

A(B,J)

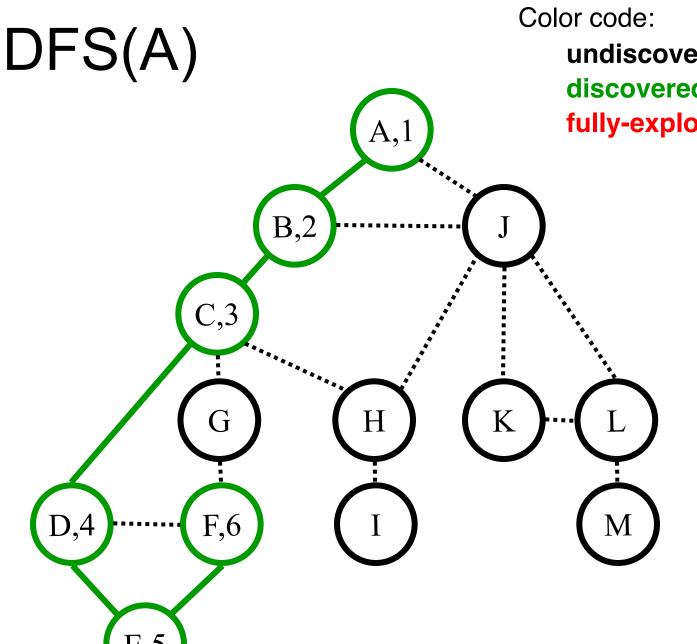
 $B(\cancel{K},\cancel{C},J)$ 

C (**B**,**D**,**G**,**H**)

D (**%**,**F**,F)

E(D,F)

st[] = {1,2,3,4,5}



discovered

fully-explored

Call Stack: (Edge list)

A(B,J)

 $B(\cancel{K},\cancel{C},J)$ 

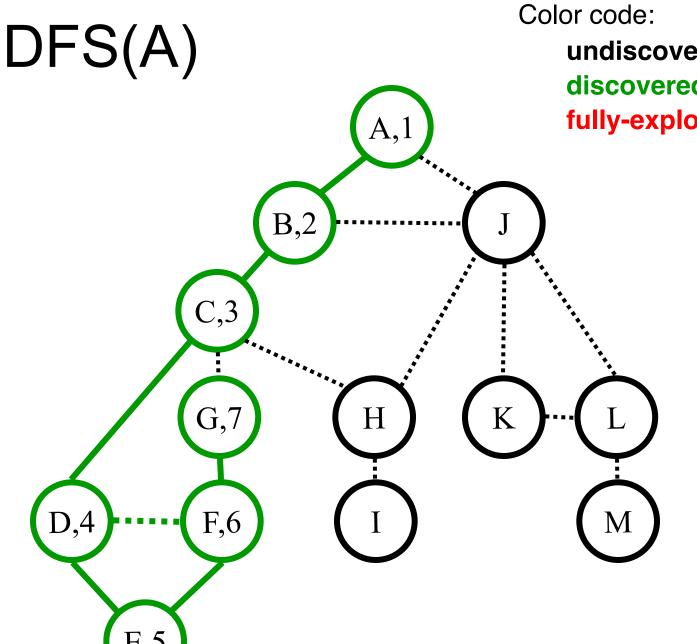
C (**p**,**p**,G,H)

D (**%**,**F**,F)

 $E(\vec{p},\vec{r})$ 

F (D,E,G)

st[] = {1,2,3,4,5, 6}



discovered

fully-explored

Call Stack: (Edge list)

A(B,J)

 $B(\cancel{K},\cancel{C},J)$ 

C (**B**,**D**,G,H)

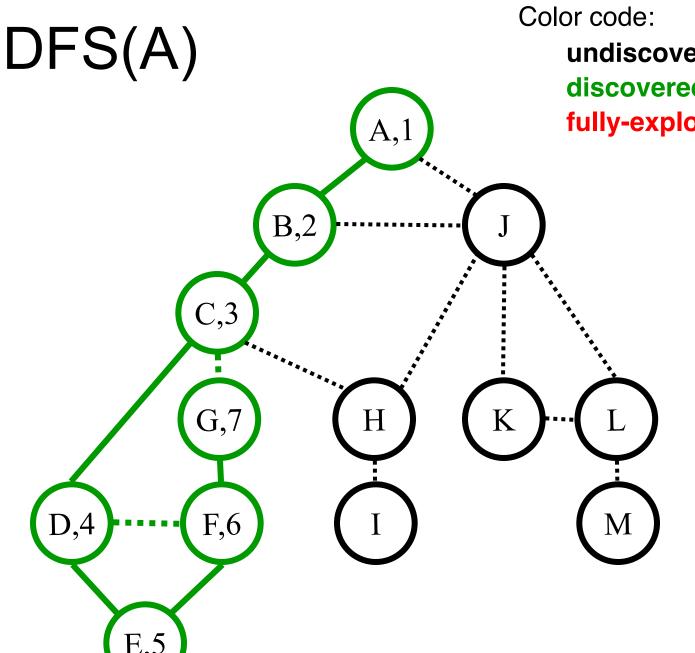
D (**∅**,**万**,F)

 $E(\vec{p},\vec{r})$ 

F (D,E,G)

G(C,F)

st[] = {1,2,3,4,5, 6,7}



discovered

fully-explored

Call Stack: (Edge list)

A(B,J)

 $B(\cancel{K},\cancel{C},J)$ 

C (**B**,**D**,G,H)

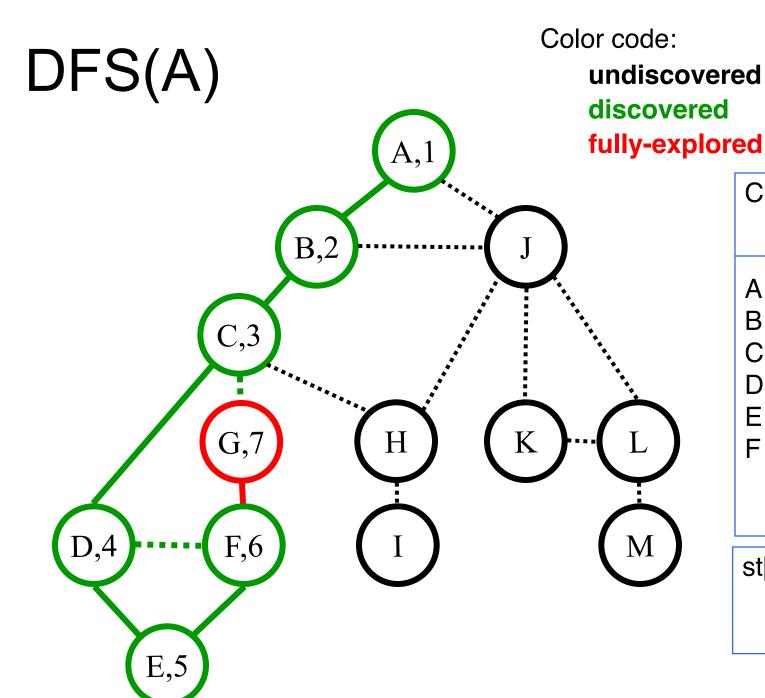
D (**∅**,**ළ**,F)

 $E(\vec{p},\vec{r})$ 

F (**D**,**E**,**G**)

G(**Ø**,**F**)

st[] = {1,2,3,4,5, 6,7}



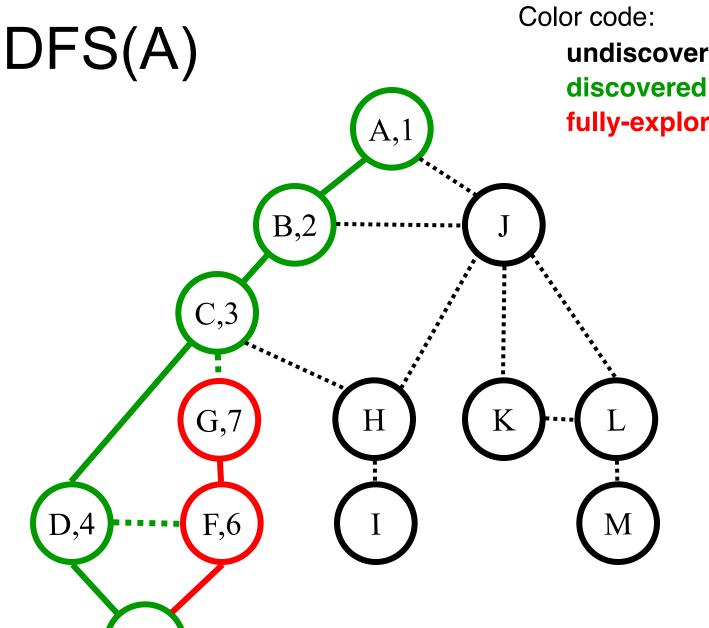
Call Stack: (Edge list)

A (₱,J) B (♠,Ø,J) C (₱,₱,G,H) D (Ø,₱,F)

 $E(\mathcal{D},\mathcal{F})$ 

F (**D**,**E**,**G**)

st[] = {1,2,3,4,5, 6}



fully-explored

Call Stack: (Edge list)

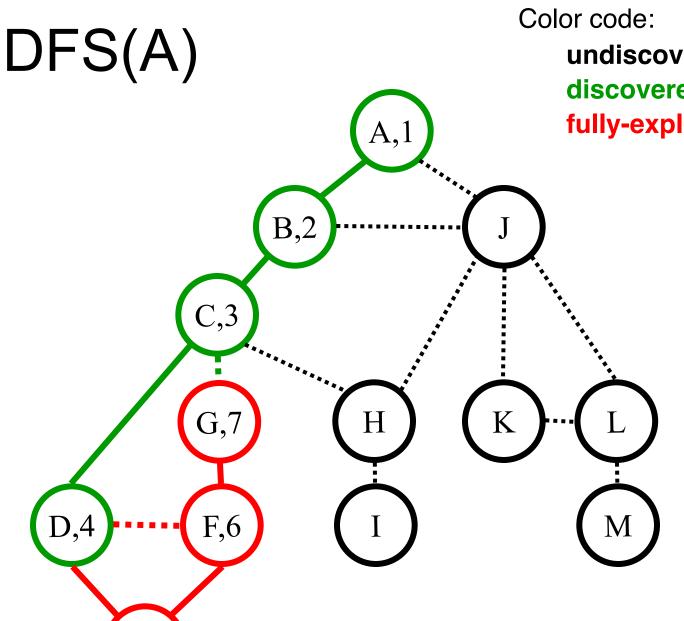
A(B,J)

B (**A**,**C**,J) C (**B**,**D**,G,H)

D (**%**,**F**,F)

 $E(\mathcal{D},\mathcal{F})$ 

st[] = {1,2,3,4,5}



discovered

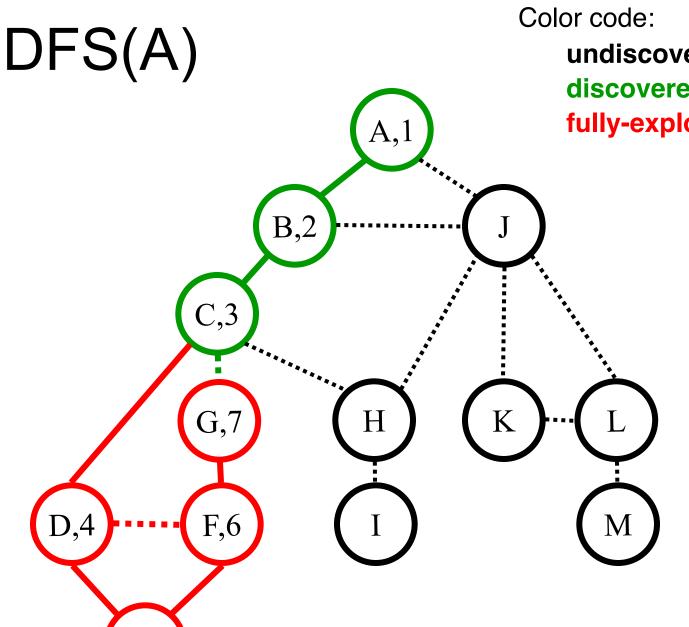
fully-explored

Call Stack: (Edge list)

A(B,J)

B (**A**,**C**,J) C (**B**,**D**,G,H) D (**C**,**E**,**F**)

st[] = {1,2,3,4}



discovered

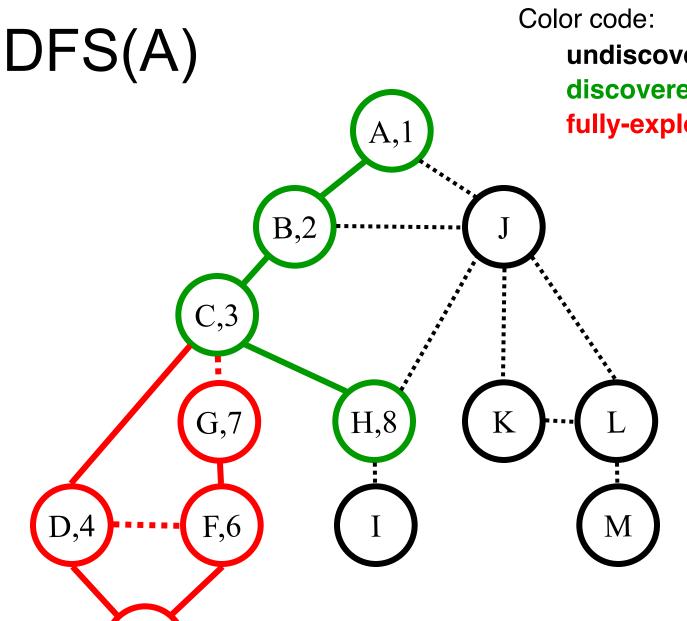
fully-explored

Call Stack: (Edge list)

A(B,J)

B (**A**,**C**,J) C (**B**,**D**,G,H)

st[] = {1,2,3}



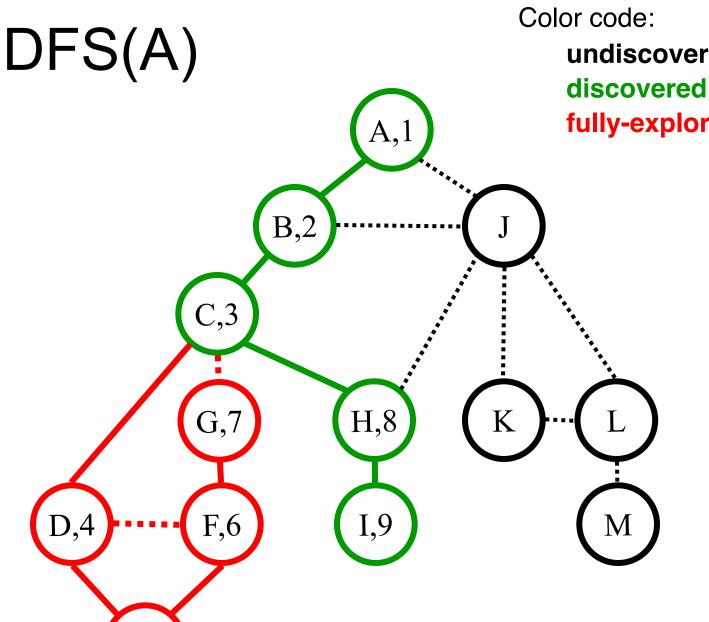
discovered

fully-explored

Call Stack: (Edge list)

A(B,J)B (**A**,**C**,**J**) C (**B**,**D**,**C**,**H**) H(C,I,J)

st[] = {1,2,3,8}

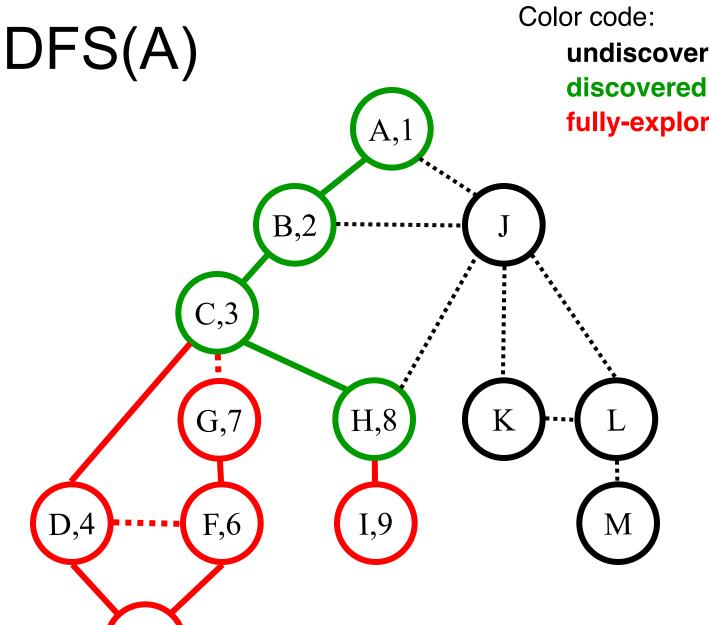


fully-explored

Call Stack: (Edge list)

A(B,J) $B(\cancel{R},\cancel{Q},J)$  $C(\cancel{R},\cancel{Q},\cancel{M},\cancel{M})$ H (**%**,**/**,**J**) I (H)

st[] = {1,2,3,8,9}

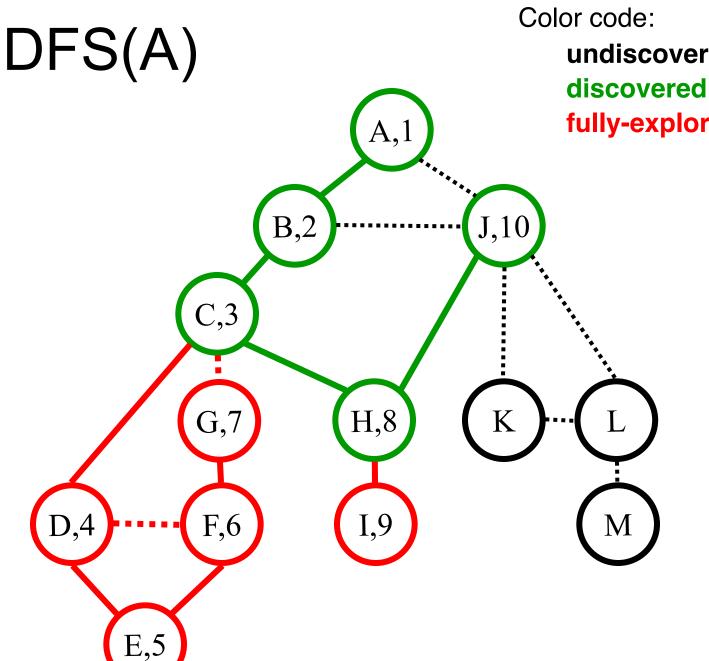


fully-explored

Call Stack: (Edge list)

A(B,J)B (**A**,**C**,**J**) C (**B**,**D**,**C**,**H**) H (**C**,**Y**,**J**)

st[] = {1,2,3,8}

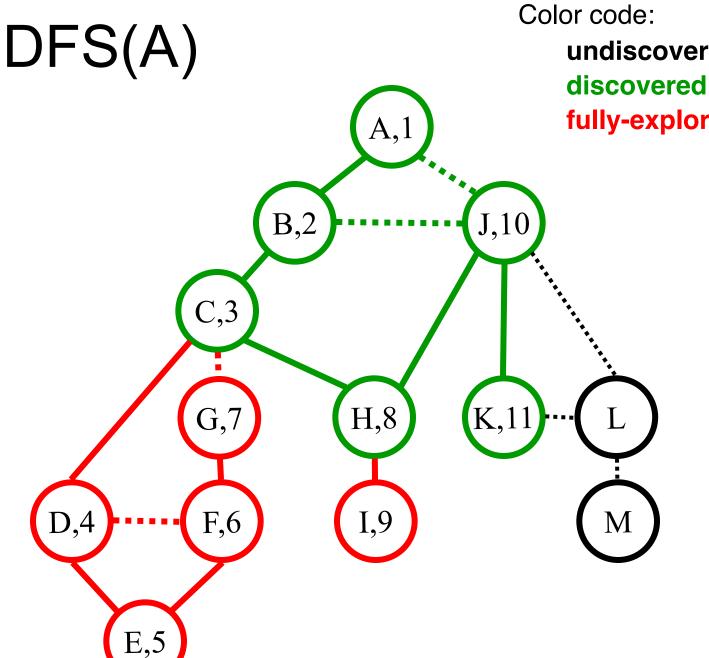


fully-explored

Call Stack: (Edge list)

A(B,J)B (**A**,**C**,J) C (**B**,**D**,**C**,**H**) H (2,1,1) J(A,B,H,K,L)

st[] = {1,2,3,8, 10}



fully-explored

Call Stack: (Edge list)

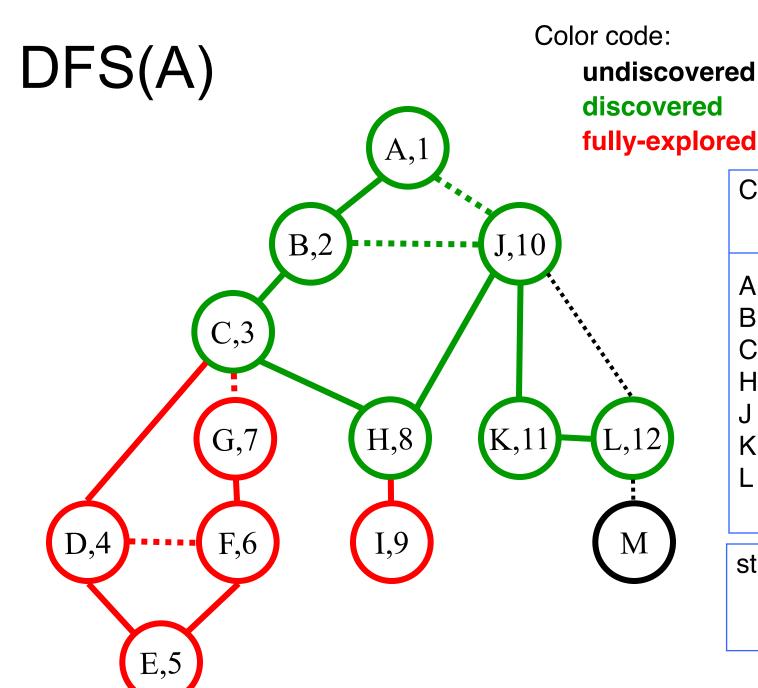
A(B,J)B (**A**,**C**,J) C (**B**,**D**,**C**,**H**)

 $H(\mathcal{C},\mathcal{V},\mathcal{Y})$ 

J(A,B,H,K,L)

**K** (J,L)

st[] = {1,2,3,8,10 ,11}



Call Stack: (Edge list)

A(B,J)

 $B(\mathcal{A},\mathcal{C},J)$ 

C(B,D,C,H)

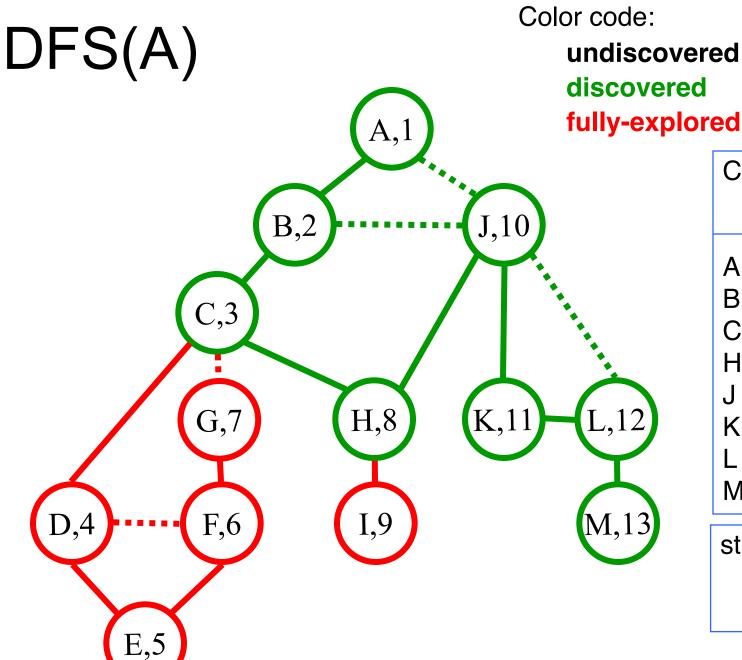
 $H(\mathcal{C},\mathcal{Y},\mathcal{Y})$ 

J(A,B,H,K,L)

K (J,L)

L(J,K,M)

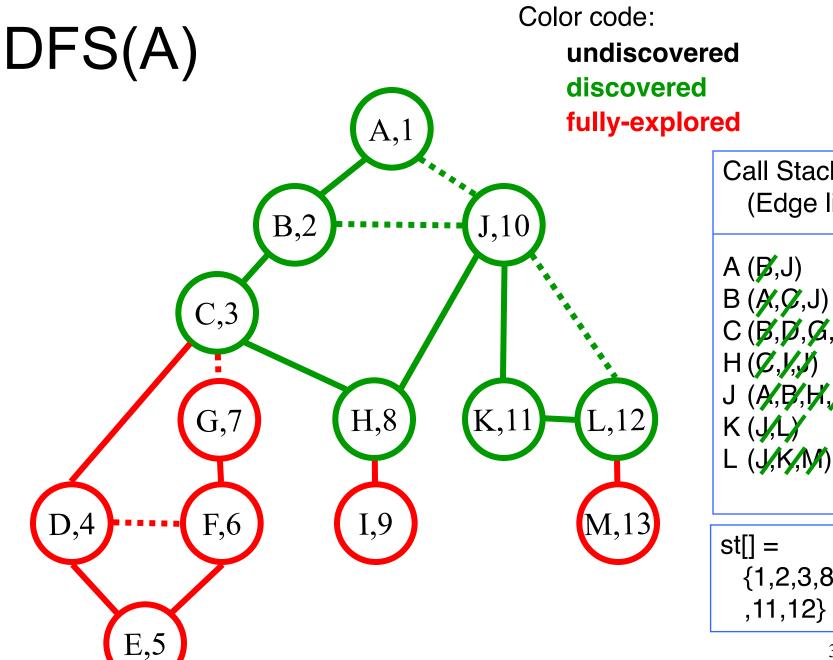
st[] = {1,2,3,8,10 ,11,12}



Call Stack: (Edge list)

A (\$\beta, \beta, \beta

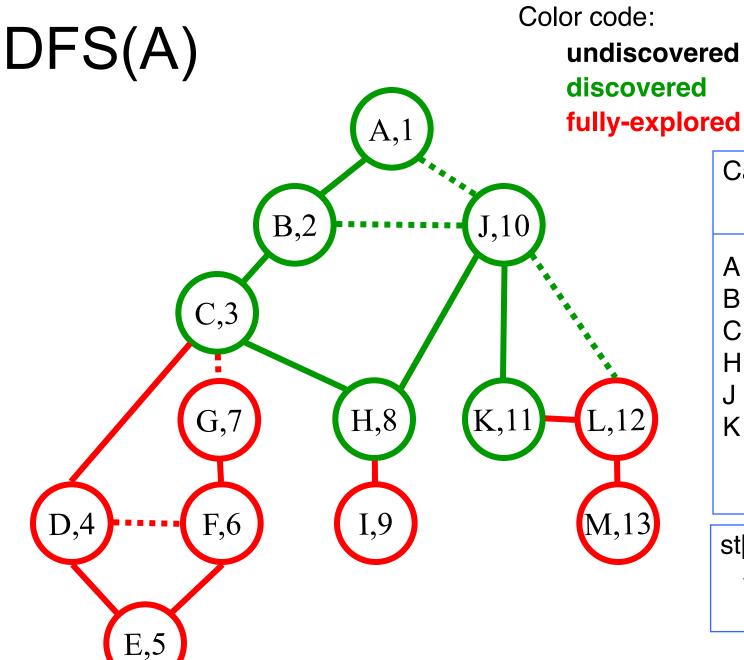
st[] = {1,2,3,8,10 ,11,12,13}



Call Stack: (Edge list)

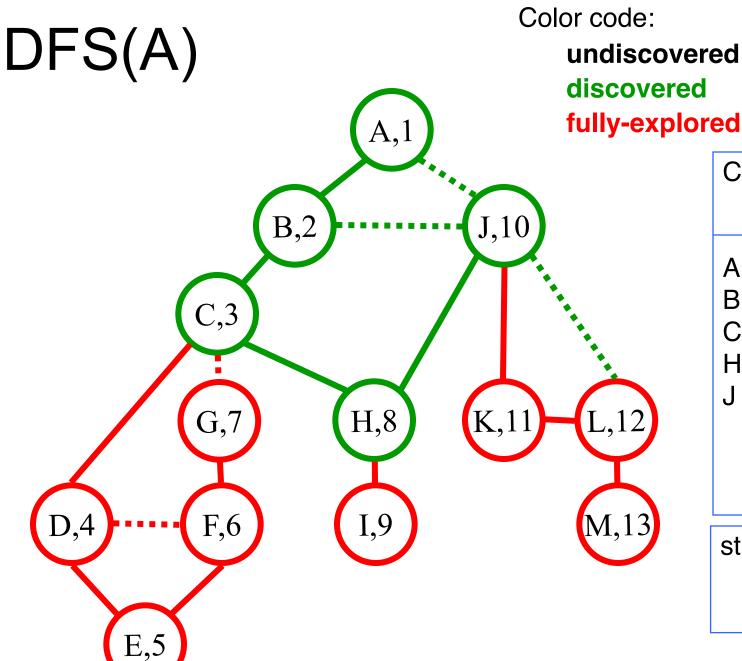
 $B(\cancel{K},\cancel{C},J)$ C(B,D,C,H) $H(\mathcal{C},\mathcal{V},\mathcal{Y})$ J(A,B,H,K,L)K (J,L)

st[] = {1,2,3,8,10 ,11,12}



Call Stack: (Edge list)

A (P,J) B (A,C,J) C (P,D,C,H) H (C,Y,Y) J (A,P,H,K,L) K (J,L)



undiscovered discovered

> Call Stack: (Edge list)

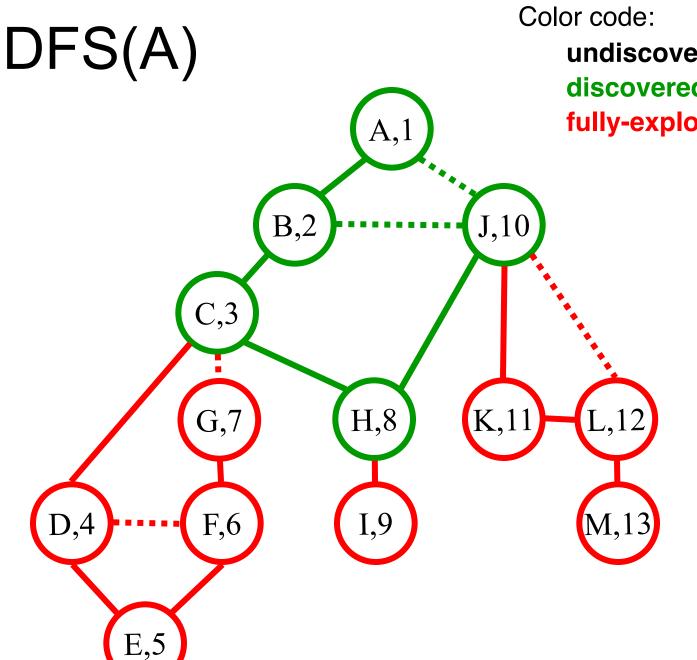
A(B,J)

B (**A**,**C**,**J**) C (**B**,**D**,**C**,**H**)

H (Z, Y, Y)

J(A,B,H,K,L)

st[] = {1,2,3,8, 10}



discovered

fully-explored

Call Stack: (Edge list)

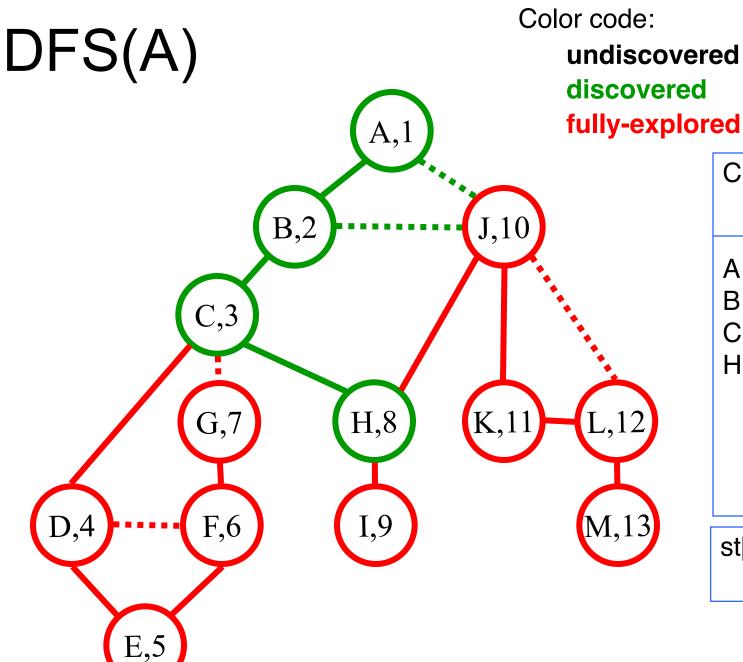
A(B,J)

B (**A**,**C**,**J**) C (**B**,**D**,**C**,**H**)

H (Z, Y, Y)

J (A,B,H,K,K)

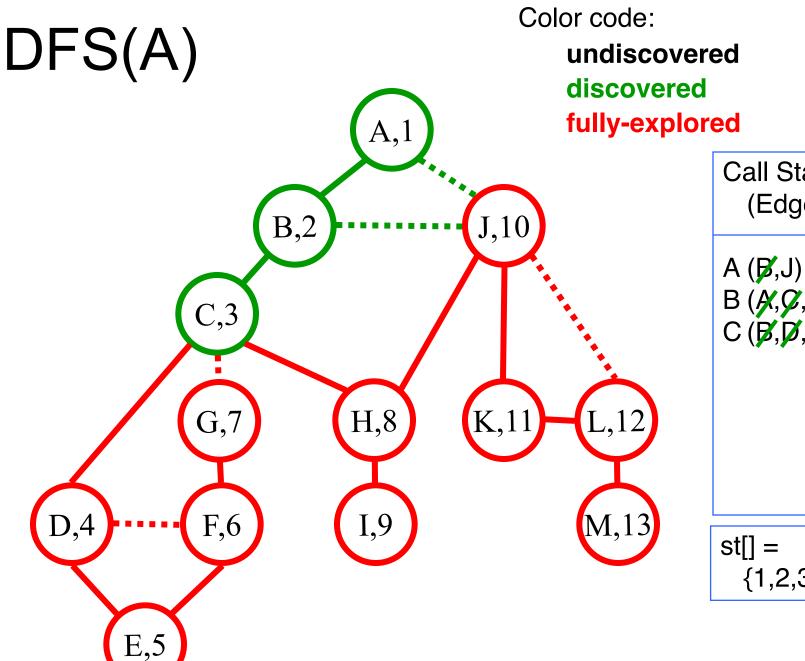
st[] = {1,2,3,8, 10}



Call Stack: (Edge list)

A (**B**',J) B (**A**',**C**',J) C (**B**',**D**',**C**',**H**') H (**C**',**Y**,**Y**')

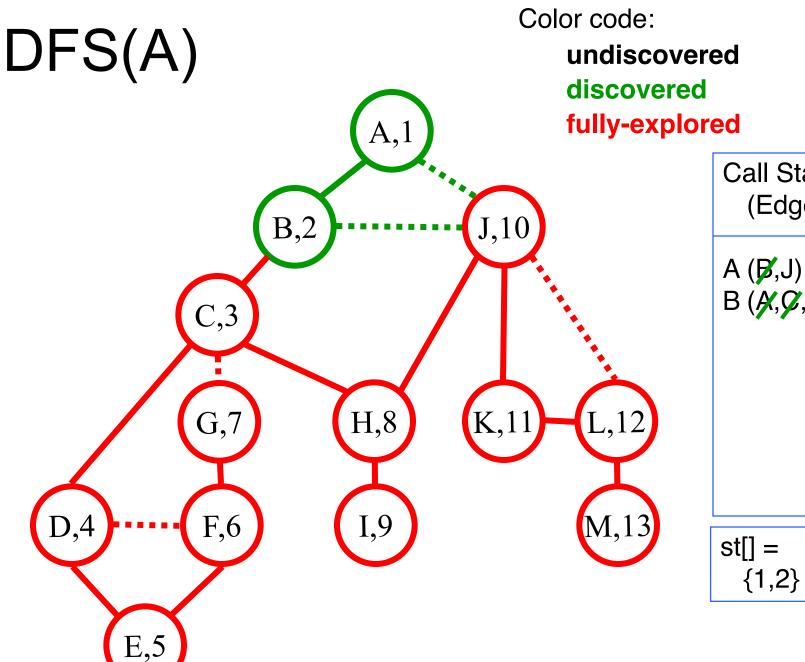
st[] = {1,2,3,8}



Call Stack: (Edge list)

B (A, Ø, J) C (B, D, Ø, H)

{1,2,3}



Call Stack: (Edge list)

B (A,Q,J)

