CSE 421: Introduction to Algorithms

1: Overview

Winter 2012

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University of Washington

Computer Science & Engineering

CSE 421, Wi '12: Introduction to Algorithms

CSE Home

Administrative
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- Schedule & Reading

Course Email/BBoard
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- GoPost BBoard

Lecture Notes

Lecture: EEB 037 (schematic) MWF 1:30-2:20

		Office Hours	Location	Phon
Instructor:	Larry Ruzzo, ruzzo@cs	M 12:00-1:00	CSE	
TA:	Danushen Gnanaprasagam, nanabyte@cs	W 4:00		
TA:	Pratik Prasad, pratikp@cs			
TA:	Milda Zizyte, mzizyte@c			

Course Email: ... interest student/staff Q&A about homew... are ... probably should change their default subscription optior... GoPost to discuss homework, etc.

... for design of efficient algorithms. Methods for showing lower bounds on compu... ..., set manipulation, arithmetic, graph problems, pattern matching.

...quisites: either CSE 312 or CSE 322; either CSE 326 or CSE 332.

Credits: 3

Grading: Homework, Midterm, Final. Homework will be a mix of paper & pencil exercises and programing. Overal

Late Policy: Unless otherwise announced, homework is due at the start of class on the due date. 20% off per day the Friday due dates).

Extra Credit: Assignments may include "extra credit" sections. These will enrich your understanding of the material them for the glory, not the points, and don't start extra credit until the basics are complete.

Textbook: *Algorithm Design* by Jon Kleinberg and Eva Tardos. Addison Wesley, 2006. (Available from U Book S

http://www.cs.washington.edu/421

What you have to do

Homework (~55% of grade)

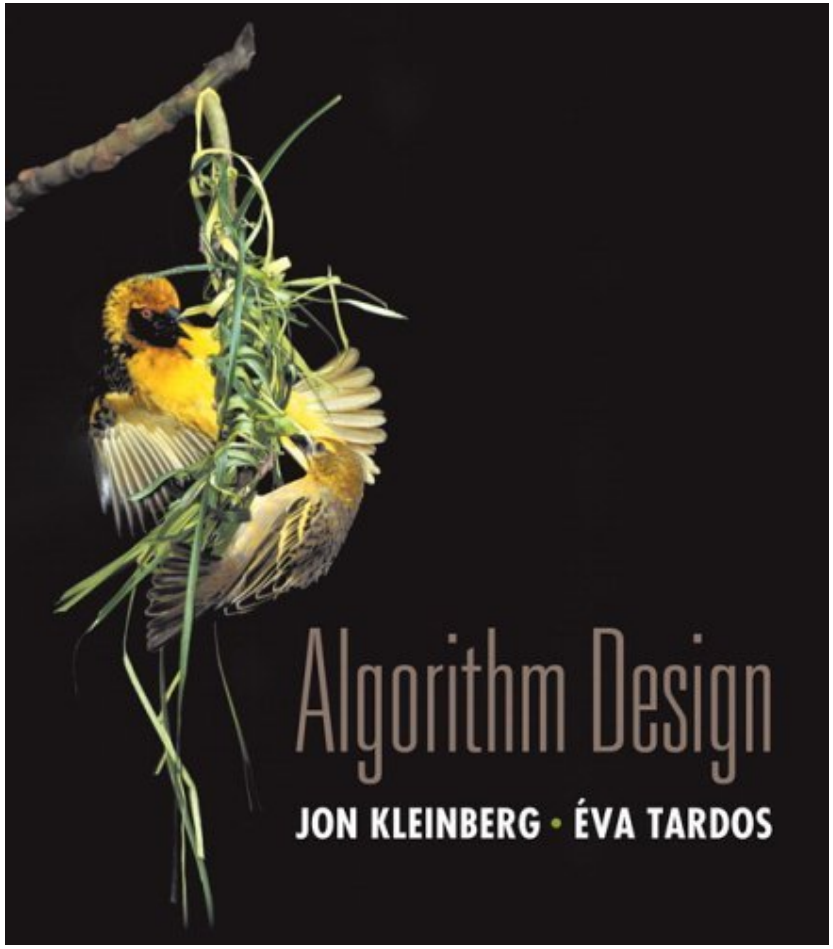
- Programming?
  - perhaps some small projects
- Written homework assignments
  - English exposition and pseudo-code
  - Analysis and argument as well as design

Midterm / Final Exam (~15% / 30%)

Late Policy:

- Papers and/or electronic turnins are due at the *start* of class on the due date; minus 20% per day thereafter

Textbook



Algorithm Design by Jon Kleinberg and Eva Tardos. Addison Wesley, 2006.

What the course is about

Design of Algorithms

- design methods
- common or important types of problems
- analysis of algorithms - efficiency
- correctness proofs

What the course is about

Complexity, NP-completeness and intractability

- solving problems in principle is not enough
  - algorithms must be *efficient*
- some problems have *no efficient solution*
- NP-complete problems
  - important & useful class of problems whose solutions (seemingly) cannot be found efficiently, but *can* be checked easily

Very Rough Division of Time

Algorithms (7 weeks)

- Analysis
- Techniques: greedy, divide&conquer, dynamic programming
- Applications: graph algorithms, flows & matchings

Complexity & NP-completeness (3 weeks)

Check online schedule page for (evolving) details

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CSE 417, Wi '06: *Approximate* Schedule

CSE Home About Us

		Due	Lecture Topic	Reading
Week 1 1/2-1/6	M		Holiday	
	W		Intro, Examples & Complexity	Ch. 1; Ch. 2
	F		Intro, Examples & Complexity	
Week 2 1/9-1/13	M		Intro, Examples & Complexity	
	W		Graph Algorithms	Ch. 3
	F		Graph Algorithms	

Complexity Example

Cryptography (e.g. RSA, SSL in browsers)

- Secret: p,q prime, say 512 bits each
- Public: n which equals p x q, 1024 bits

In principle

- *there is an algorithm* that given n will find p and q:
  try all $2^{512} > 1.3 \times 10^{154}$ possible p's: kinda slow…

In practice

- *no fast algorithm* known for this problem (on non-quantum computers)
- security of RSA depends on this fact
- ("quantum computing": strongly driven by possibility of changing this)

Algorithms versus Machines

We all know about Moore's Law and the exponential improvements in hardware...

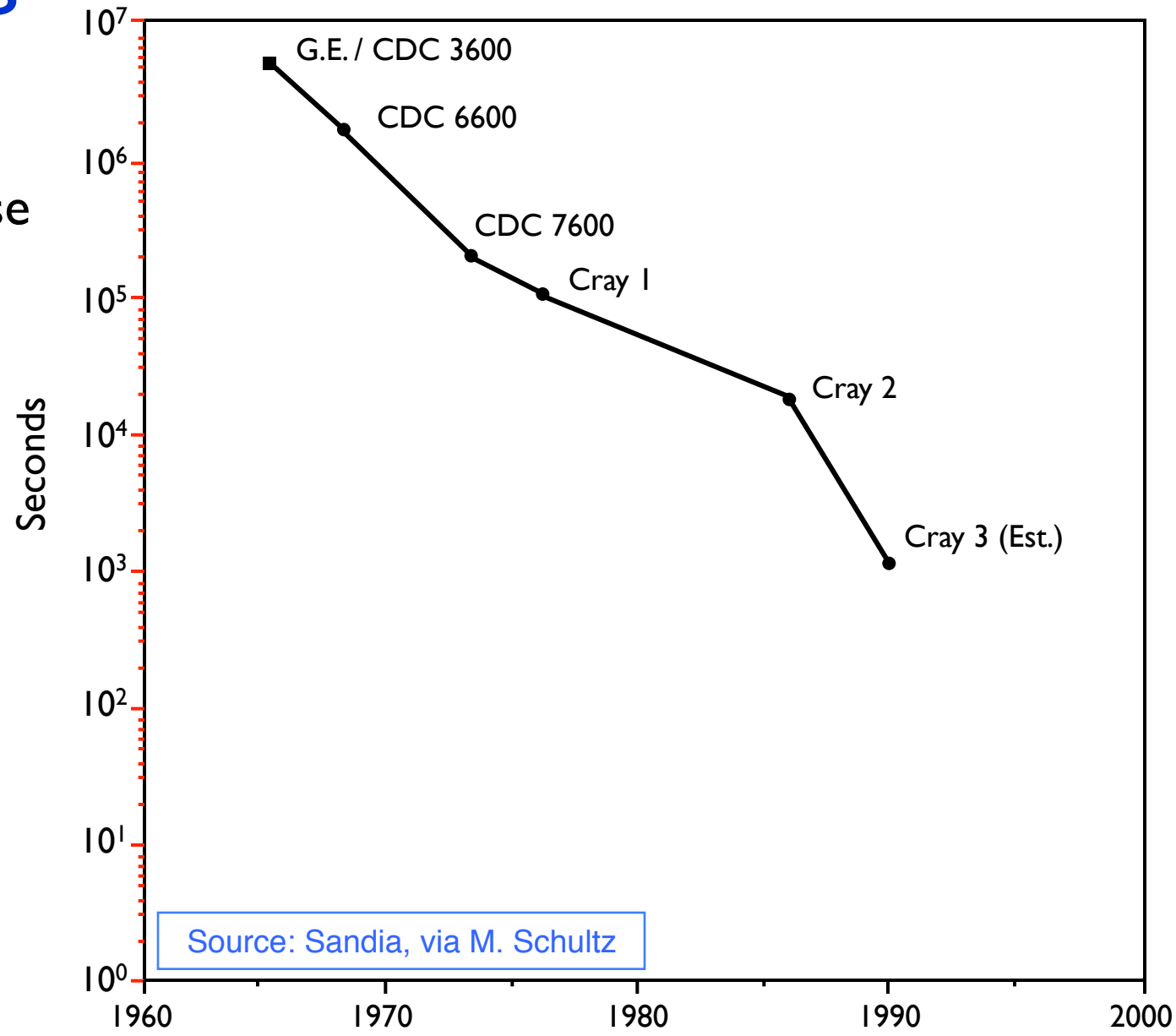
Ex: sparse linear equations over 25 years

10 orders of magnitude improvement!

Algorithms or Hardware?

25 years progress solving sparse linear systems

hardware: 4 orders of magnitude

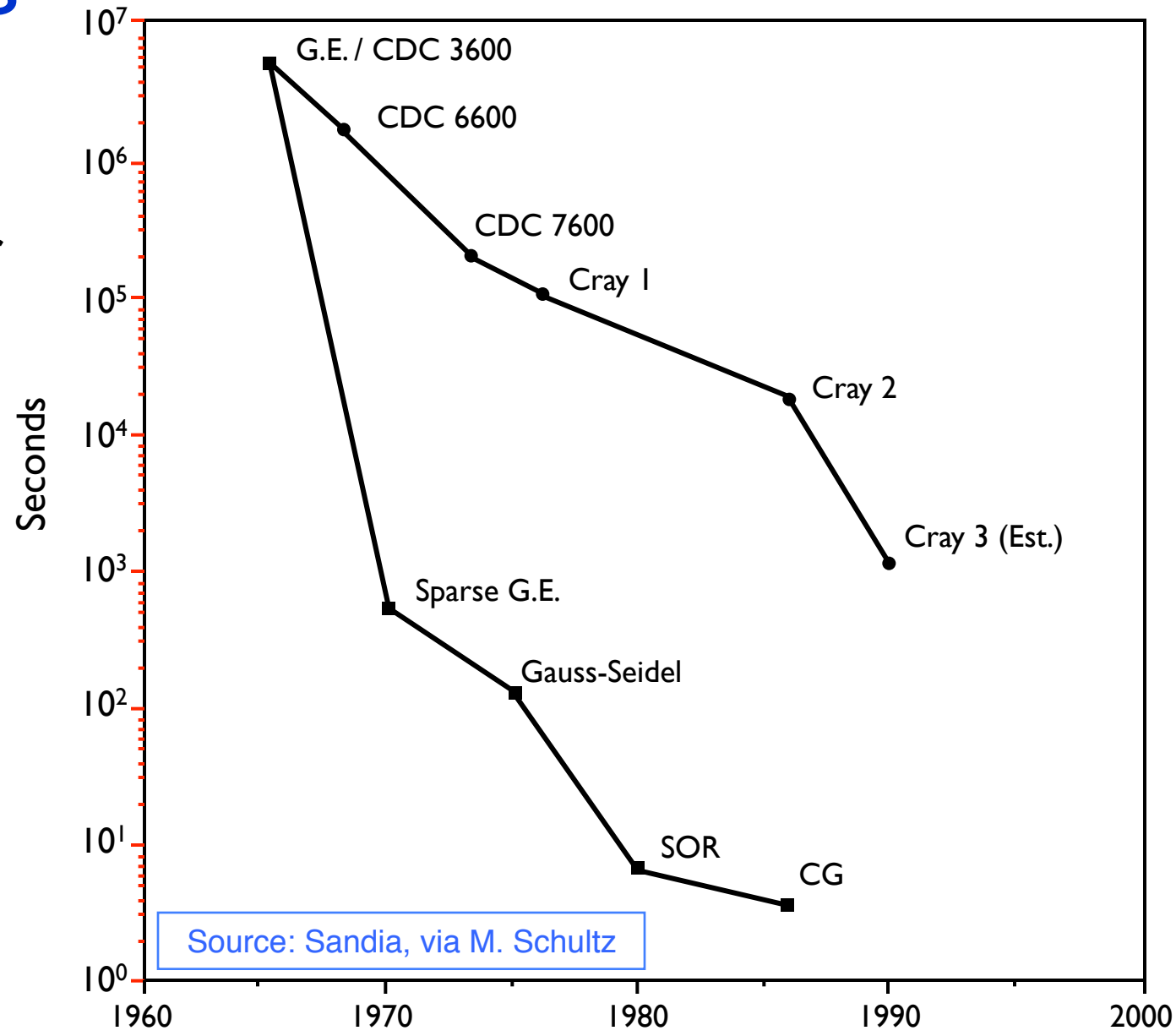


Algorithms or Hardware?

25 years progress solving sparse linear systems

hardware: 4 orders of magnitude

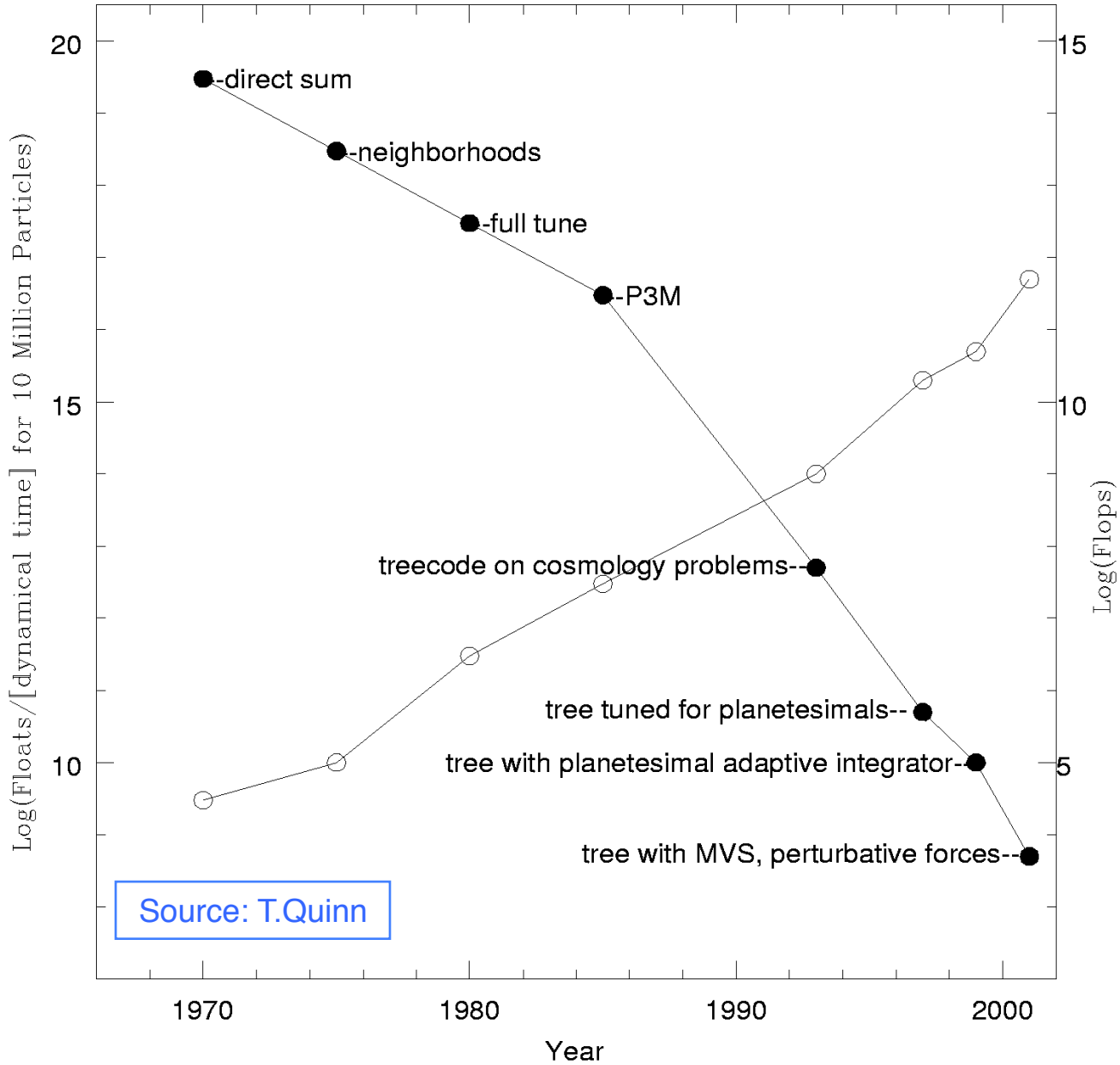
software: 6 orders of magnitude



Algorithms or Hardware?

The
N-Body
Problem:

in 30 years
10^7 hardware
10^{10} software



Algorithm: definition

Procedure to accomplish a task or solve a well-specified problem

- Well-specified: know what all possible inputs look like and what output looks like given them
- "accomplish" via simple, well-defined steps
- Ex: sorting names (via comparison)
- Ex: checking for primality (via +, -, *, /, ≤)

Goals

Correctness

- often subtle

Analysis

- often subtle

Generality, Simplicity, ‘Elegance’

Efficiency

- time, memory, network bandwidth, …

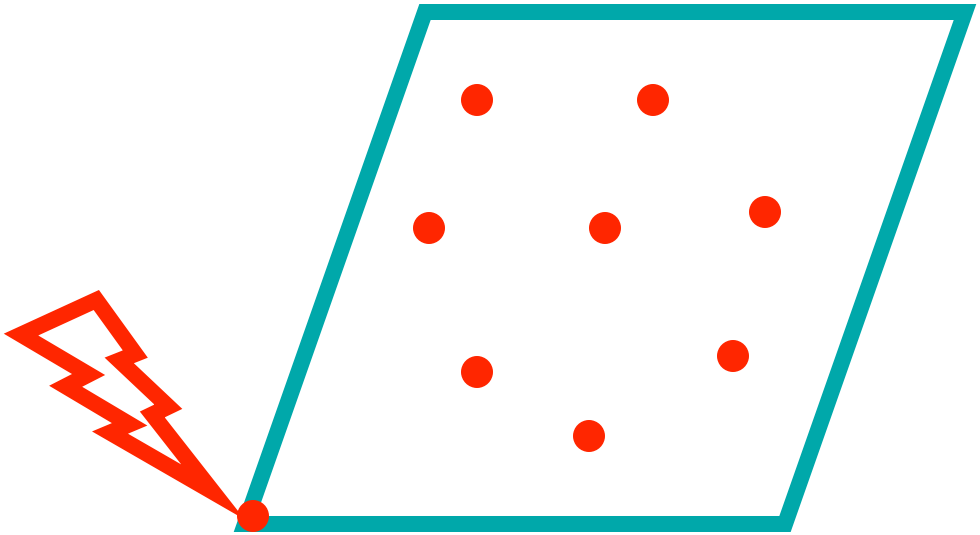
Algorithms: a sample problem

Printed circuit-board company has a robot arm that solders components to the board

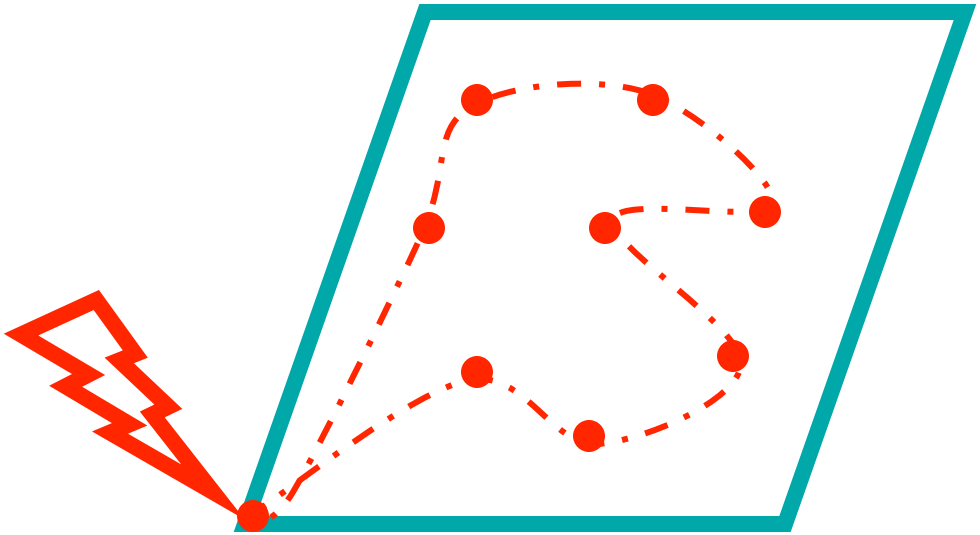
Time: proportional to total distance the arm must move from initial rest position around the board and back to the initial position

For each board design, find best order to do the soldering

Printed Circuit Board



Printed Circuit Board



A Well-defined Problem

Input: Given a set S of n points in the plane

Output: The shortest cycle tour that visits each point in the set S.

Better known as "TSP"

How might you solve it?

Nearest Neighbor Heuristic

heuristic:
A rule of thumb, simplification, or educated guess that reduces or limits the search for solutions in domains that are difficult and poorly understood. May be good, but usually *not* guaranteed to give the best or fastest solution.

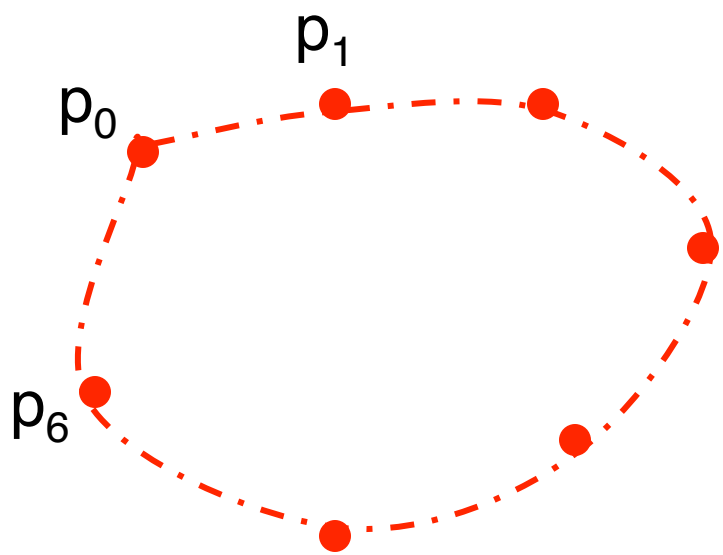
Start at some point p_0

Walk first to its nearest neighbor p_1

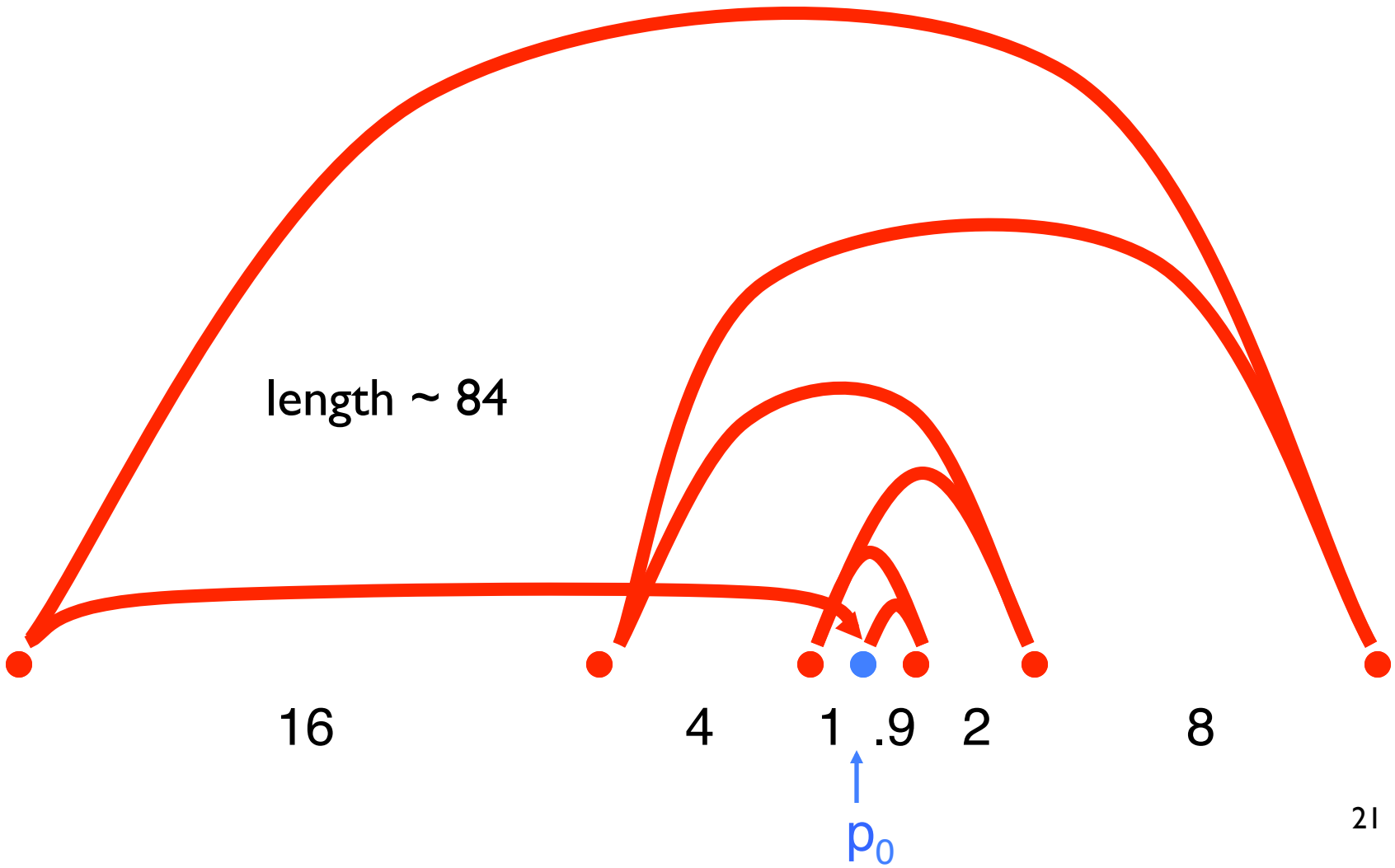
Repeatedly walk to the nearest unvisited neighbor p_2, then p_3,… until all points have been visited

Then walk back to p_0

Nearest Neighbor Heuristic



An input where it works badly



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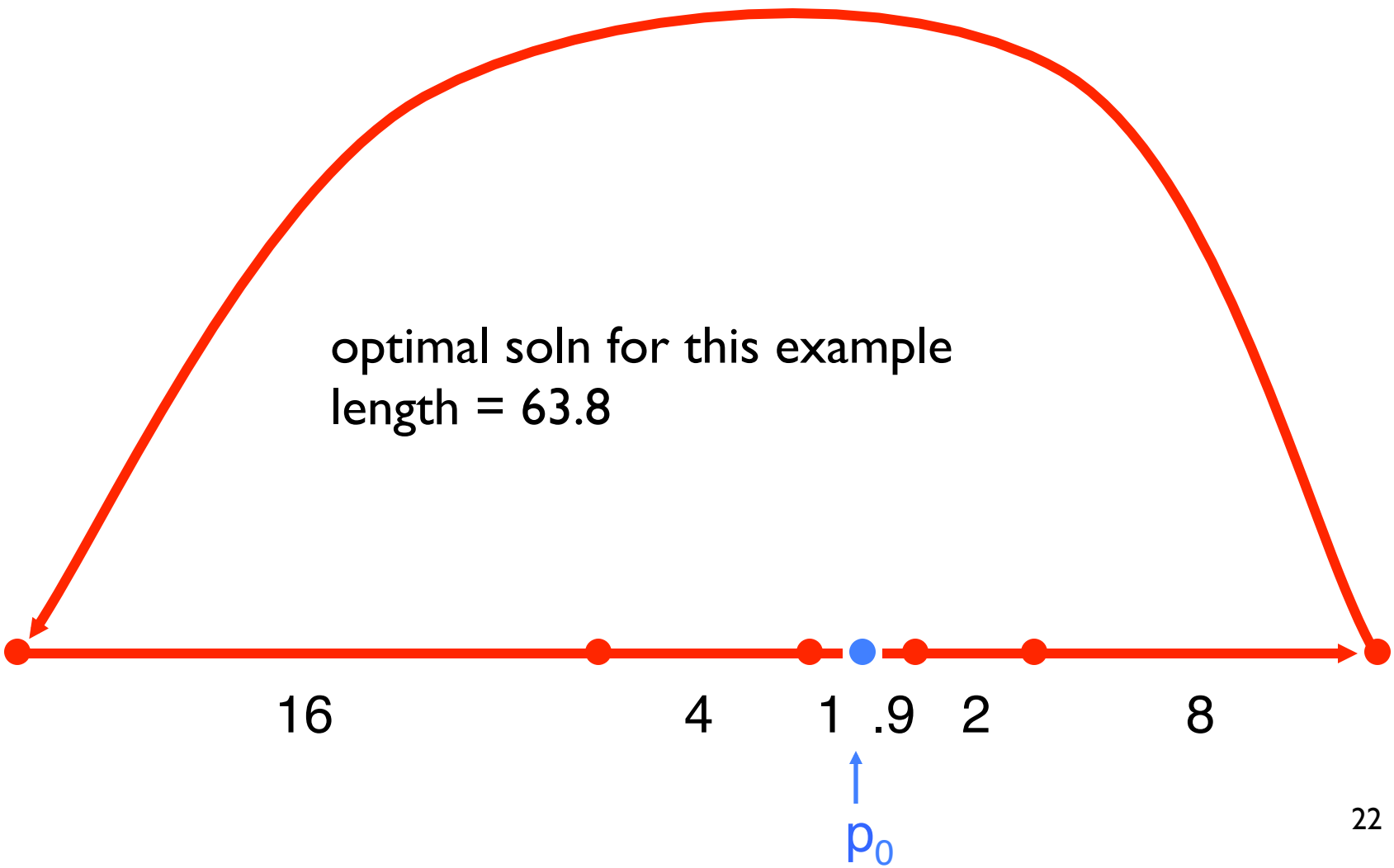
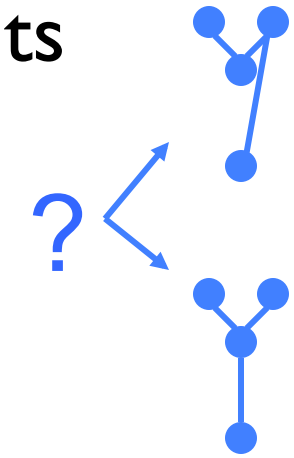


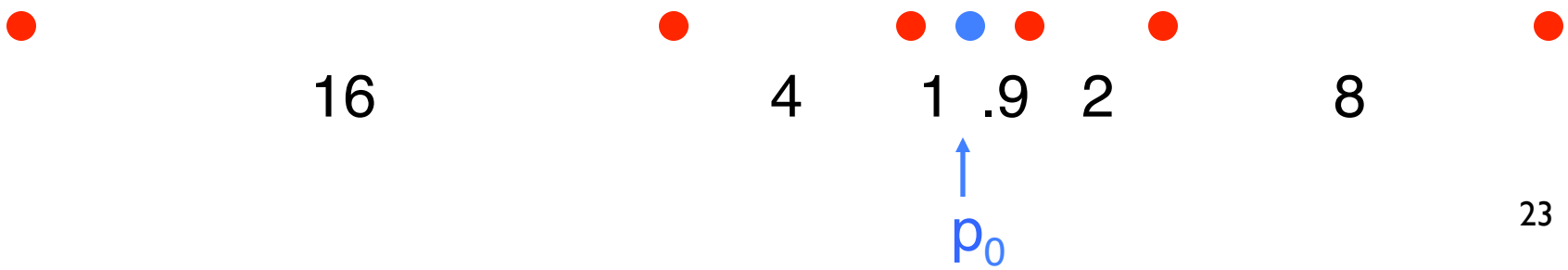
Revised idea - Closest pairs first

Repeatedly join the closest pair of points

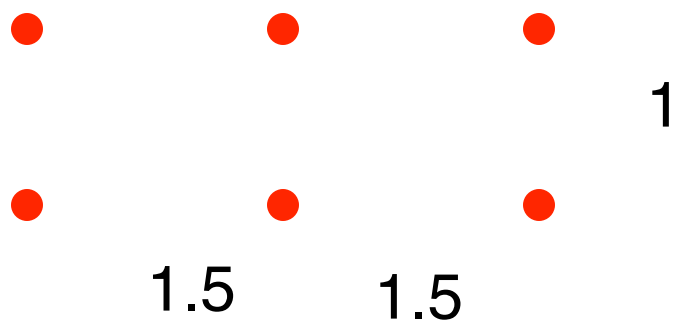
(s.t. result can still be part of a single loop in the end. I.e., join endpoints, but not points in middle, of path segments already created.)



How does this work on our bad example?



Another bad example



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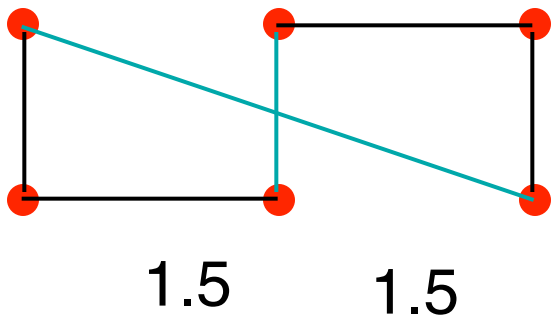
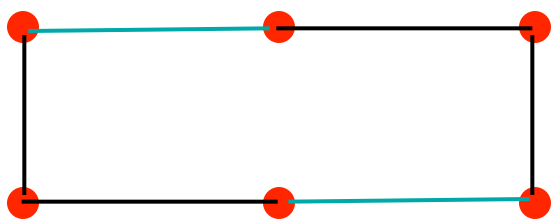


$6+\sqrt{10} = 9.16$

vs



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Something that works

"Brute Force Search":

For each of the $n! = n(n-1)(n-2)\ldots 1$ orderings of the points, check the length of the cycle you get

Keep the best one

Two Notes

The two *in*correct algorithms were greedy

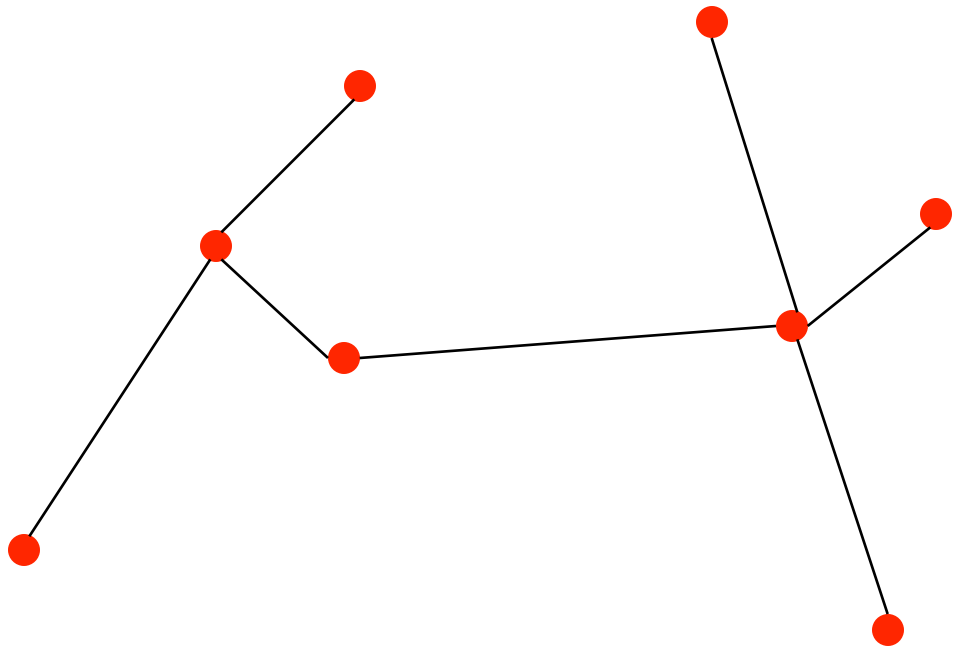
- Often very natural & tempting ideas
- They make choices that look great "locally" (and never reconsider them)
- When greed works, the algorithms are typically efficient
- BUT: often does not work - you get boxed in

Our correct alg avoids this, but is incredibly slow

- 20! is so large that checking one billion orderings per second would take 2.4 billion seconds (around 70 years!)
- And *growing*: $n! \sim \sqrt{2 \pi n} \cdot (n/e)^n \sim 2^{O(n \log n)}$

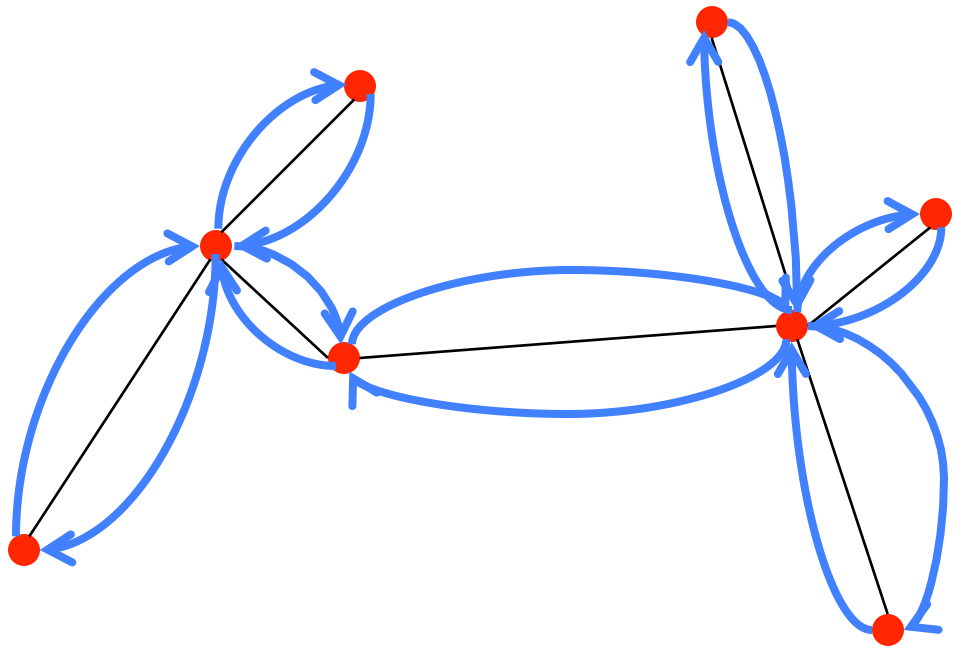
Something that “works” (differently)

1. Find Min Spanning Tree



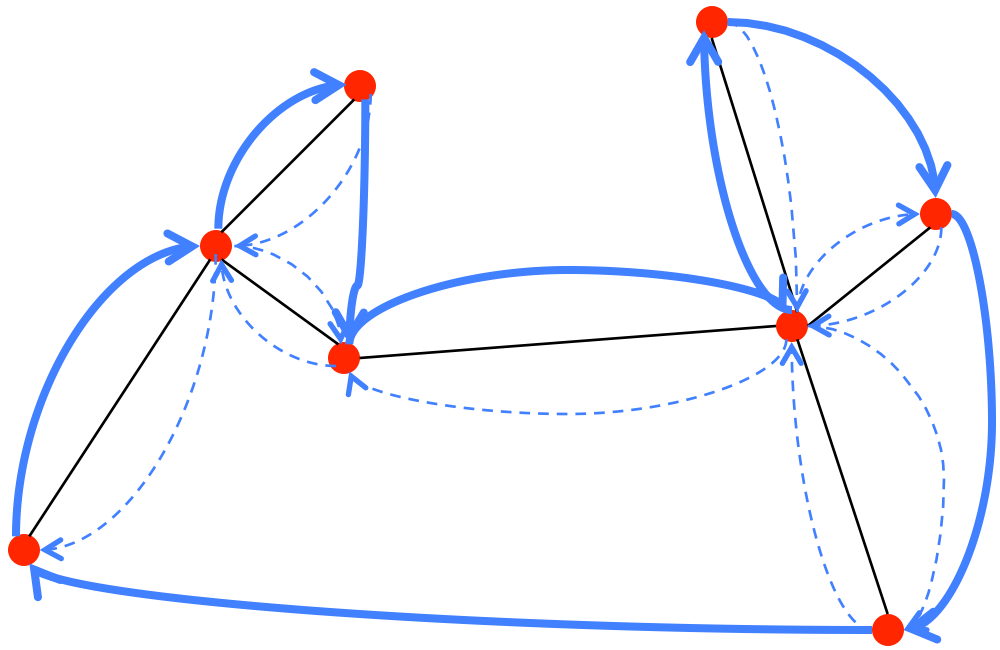
Something that “works” (differently)

2. Walk around it



Something that “works” (differently)

3. Take shortcuts (instead of revisiting)



Something that "works" (differently): Guaranteed Approximation

Does it seem wacky?

Maybe, but it's *always* within a factor of 2 of the best tour!

deleting one edge from best tour gives *a* spanning tree, so *Min* spanning tree < best tour

best tour ≤ wacky tour ≤ 2 * MST < 2 * best

triangle inequality

The Morals of the Story

Algorithms are important

- Many performance gains outstrip Moore's law

Simple problems can be hard

- Factoring, TSP

Simple ideas don't always work

- Nearest neighbor, closest pair heuristics

Simple algorithms can be very slow

- Brute-force factoring, TSP

And: for some problems, even the *best* algorithms are slow

The Morals of the Story

Algorithms are important

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Changing your objective can be good

- Guaranteed approximation for TSP

And: for some problems, even the *best* algorithms are slow