

CSE 417 Algorithms and Complexity

Winter 2023
Lecture 16
Divide and Conquer and Recurrences

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Divide and Conquer

- Recurrences, Sections 5.1 and 5.2
- Algorithms
 - Median (Selection)
 - Fast Matrix Multiplication
 - Counting Inversions (5.3)
 - Multiplication (5.5)

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Divide and Conquer : Merge Sort

```
Array MSort(Array a, int n){  
  if (n <= 1) return a;  
  return Merge(MSort(a[0 .. n/2], n/2), MSort(a[n/2+1 .. n-1], n/2);  
}
```

$T(n) = 2T(n/2) + n$; $T(1) = 1$;

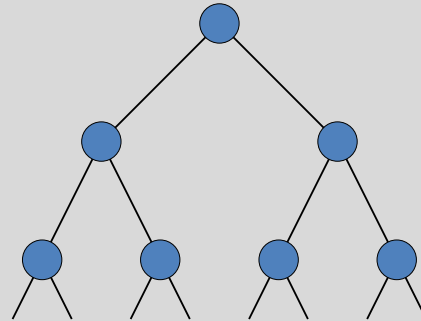
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Unrolling the recurrence



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A better mergesort (?)

- Divide into 3 subarrays and recursively sort
- Apply 3-way merge

What is the recurrence?
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Unroll recurrence for $T(n) = 3T(n/3) + n$

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$$T(n) = aT(n/b) + f(n)$$

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$$T(n) = T(n/2) + cn$$

Where does this recurrence arise?

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Solving the recurrence exactly

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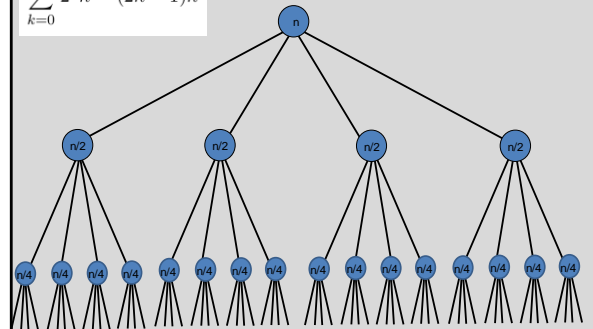
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Total Work

$$\sum_{k=0}^{\log n} 2^k n = (2n - 1)n$$

$$T(n) = 4T(n/2) + n$$



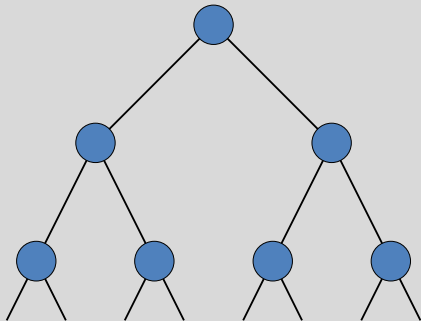
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$$T(n) = 2T(n/2) + n^2$$



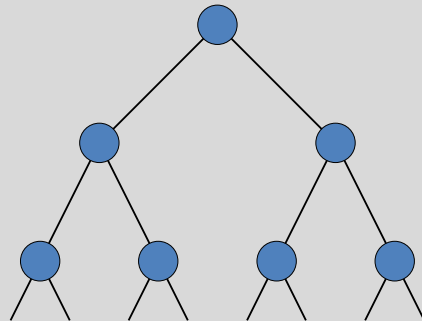
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$$T(n) = 2T(n/2) + n^{1/2}$$



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Recurrences

- Three basic behaviors
 - Dominated by initial case
 - Dominated by base case
 - All cases equal – we care about the depth

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What you really need to know about recurrences

- Work per level changes geometrically with the level
- Geometrically increasing ($x > 1$)
 - The bottom level wins
- Geometrically decreasing ($x < 1$)
 - The top level wins
- Balanced ($x = 1$)
 - Equal contribution

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Classify the following recurrences (Increasing, Decreasing, Balanced)

- $T(n) = n + 5T(n/8)$
- $T(n) = n + 9T(n/8)$
- $T(n) = n^2 + 4T(n/2)$
- $T(n) = n^3 + 7T(n/2)$
- $T(n) = n^{1/2} + 3T(n/4)$

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Recursive Matrix Multiplication

Multiply 2 x 2 Matrices:

$$\begin{vmatrix} r & s \\ t & u \end{vmatrix} = \begin{vmatrix} a & b \\ c & d \end{vmatrix} \begin{vmatrix} e & g \\ f & h \end{vmatrix}$$

$$\begin{aligned} r &= ae + bf \\ s &= ag + bh \\ t &= ce + df \\ u &= cg + dh \end{aligned}$$

A $N \times N$ matrix can be viewed as a 2×2 matrix with entries that are $(N/2) \times (N/2)$ matrices.

The recursive matrix multiplication algorithm recursively multiplies the $(N/2) \times (N/2)$ matrices and combines them using the equations for multiplying 2×2 matrices

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Recursive Matrix Multiplication

- How many recursive calls are made at each level?
- How much work in combining the results?
- What is the recurrence?

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What is the run time for the recursive Matrix Multiplication Algorithm?

- Recurrence:

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Strassen's Algorithm

Multiply 2 x 2 Matrices:

$$\begin{array}{|c|c|} \hline r & s \\ \hline t & u \\ \hline \end{array} = \begin{array}{|c|c|} \hline a & b \\ \hline c & d \\ \hline \end{array} \begin{array}{|c|c|} \hline e & g \\ \hline f & h \\ \hline \end{array}$$

$$r = p_1 + p_2 - p_4 + p_6$$

$$s = p_4 + p_5$$

$$t = p_6 + p_7$$

$$u = p_2 - p_3 + p_5 - p_7$$

Where:

$$p_1 = (b - d)(f + h)$$

$$p_2 = (a + d)(e + h)$$

$$p_3 = (a - c)(e + g)$$

$$p_4 = (a + b)h$$

$$p_5 = a(g - h)$$

$$p_6 = d(f - e)$$

$$p_7 = (c + d)e$$

From AHU 1974

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Recurrence for Strassen's Algorithms

- $T(n) = 7 T(n/2) + cn^2$
- What is the runtime?

$$\log_2 7 = 2.8073549221$$

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BFPRT Recurrence

$$T(n) \leq T(3n/4) + T(n/5) + 20n$$

What bound do you expect?

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