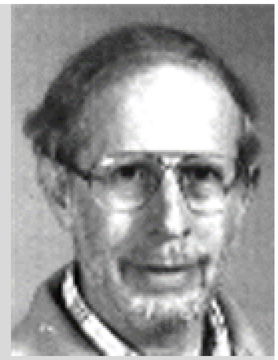
Lecture12



CSE 417

Algorithms and Complexity

Winter 2023

Lecture 12

Shortest Paths Algorithm and Minimum Spanning Trees

2/1/2023 CSE 417 1

Announcements

- Reading
  - 4.4, 4.5, 4.7
- Midterm
  - Wednesday, February 8
  - In class, closed book
  - Material through 4.7
  - Old midterm questions available
    - Note – some listed questions are out of scope
- No homework due on February 10

Assume all edges have non-negative cost

Dijkstra's Algorithm

S = { }; d[s] = 0; d[v] = infinity for v != s

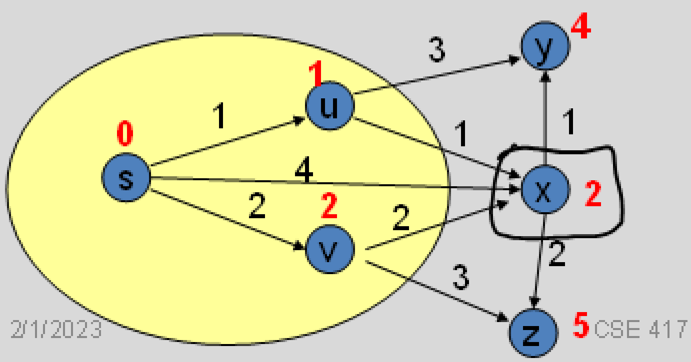
While S != V

Choose v in V-S with minimum d[v]

Add v to S

For each w in the neighborhood of v

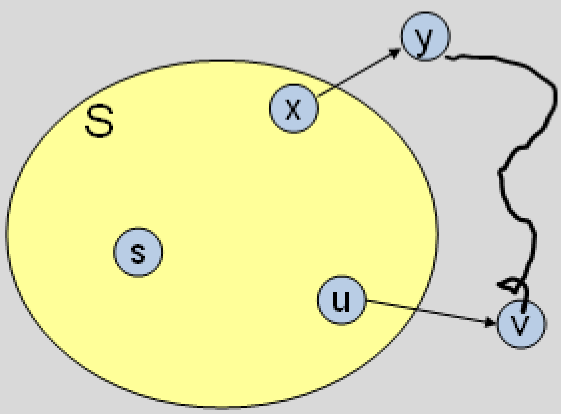
d[w] = min(d[w], d[v] + c(v, w))



Flow Jack
Bottle S.P.
d[v]+c(v,w)
+
max(d[v], c(v,p))

Correctness Proof

- Elements in S have the correct label
- Induction: when v is added to S, it has the correct distance label
  - Dist(s, v) = d[v] when v added to S



Heap - O(log n)

Dijkstra Implementation

S = { }; d[s] = 0; d[v] = infinity for v != s

While S != V

    Choose v in V-S with minimum d[v]

    Add v to S

    For each w in the neighborhood of v

        d[w] = min(d[w], d[v] + c(v, w))

DeleteMin

Update Key

m decrease Key

- Basic implementation requires Heap for tracking the distance values
- Run time O(m log n)

O(n^2) Implementation for Dense Graphs

```
FOR i := 1 TO n
       d[i] := Infinity;  visited[i] := FALSE;
d[s] := 0;

FOR i := 1 TO n
       v := -1;  dMin := Infinity;
       FOR j := 1 TO n
              IF visited[j] = FALSE AND d[j] < dMin
                     v := j; dMin := d[j];
       IF v = -1
              RETURN;
       visited[v] := TRUE;

       FOR j := 1 TO n
              IF d[v] + len[v, j] < d[j]
                     d[j] := d[v] + len[v, j];
                     prev[j] := v;
```

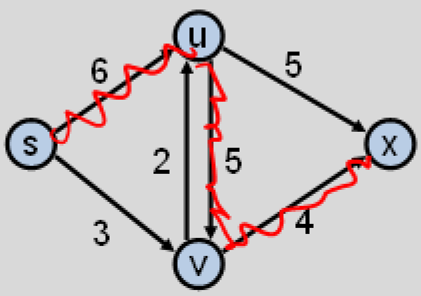
Future stuff for shortest paths

- Bellman-Ford Algorithm
  - O(nm) time
  - Handles negative cost edges
    - Identifies negative cost cycle if present
  - Dynamic programming algorithm
  - Very easy to implement

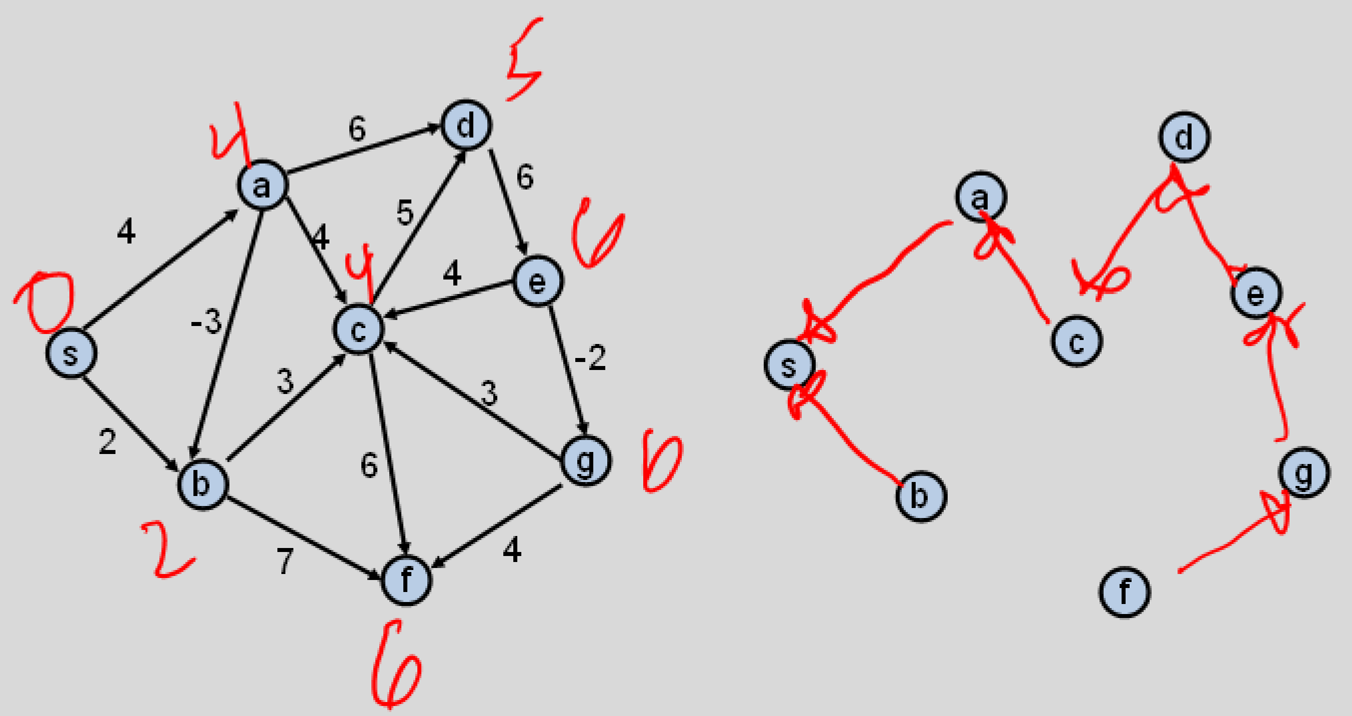
Bottleneck Shortest Path

- Define the bottleneck distance for a path to be the maximum cost edge along the path

Len P = max edge cost



Compute the bottleneck shortest paths



How do you adapt Dijkstra's algorithm to handle bottleneck distances

- Does the correctness proof still apply?

Dijkstra's Algorithm for Bottleneck Shortest Paths

S = { }; d[s] = negative infinity; d[v] = infinity for v != s

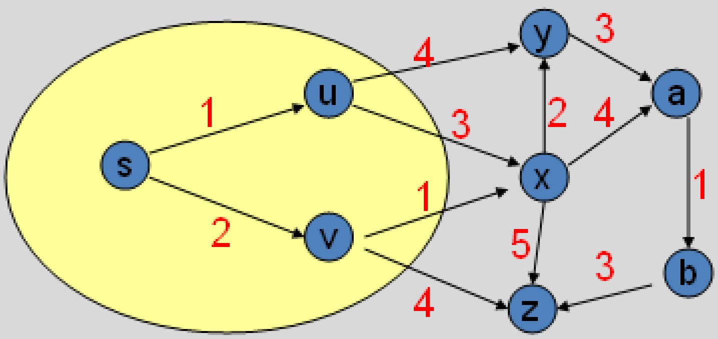
While S != V

Choose v in V-S with minimum d[v]

Add v to S

For each w in the neighborhood of v

d[w] = min(d[w], max(d[v], c(v, w)))

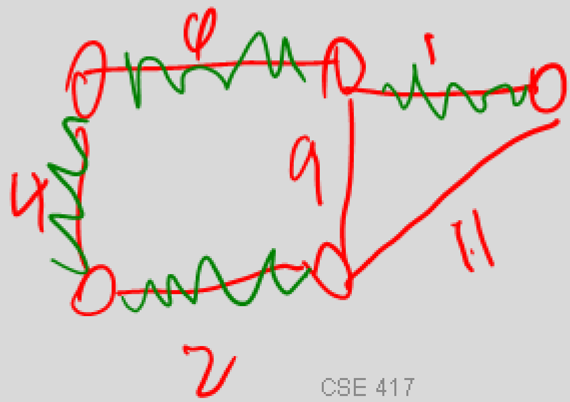


Minimum Spanning Tree

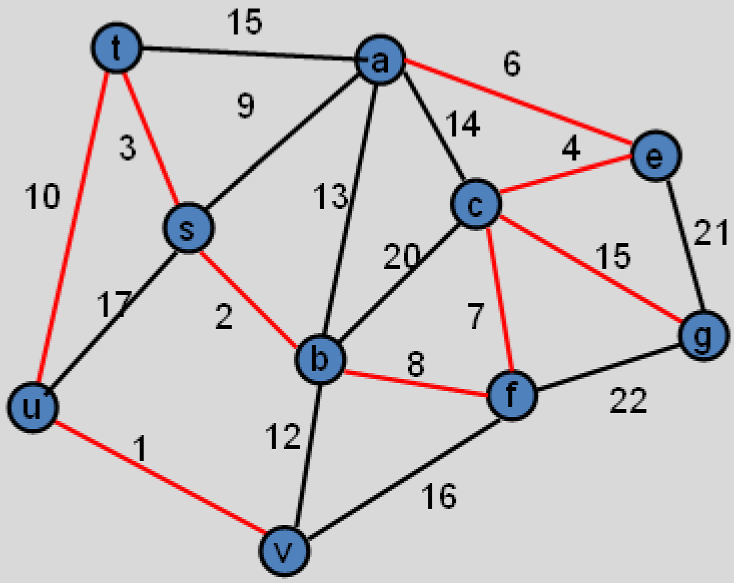
- Introduce Problem
- Demonstrate three different greedy algorithms
- Provide proofs that the algorithms work

Minimum Spanning Tree Definitions

- G=(V,E) is an UNDIRECTED graph
- Weights associated with the edges
- Find a spanning tree of minimum weight
  - If not connected, complain

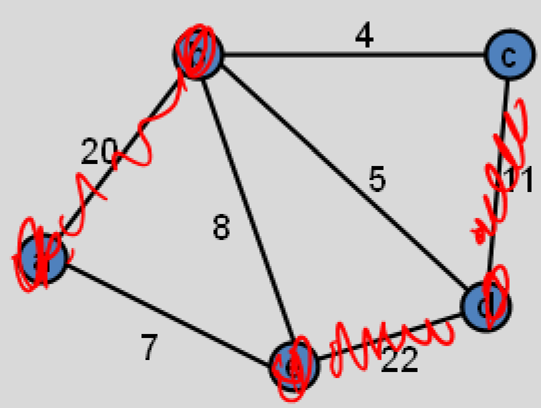


Minimum Spanning Tree



Greedy Algorithms for Minimum Spanning Tree

- Extend a tree by including the cheapest out going edge
- Add the cheapest edge that joins disjoint components
- Delete the most expensive edge that does not disconnect the graph

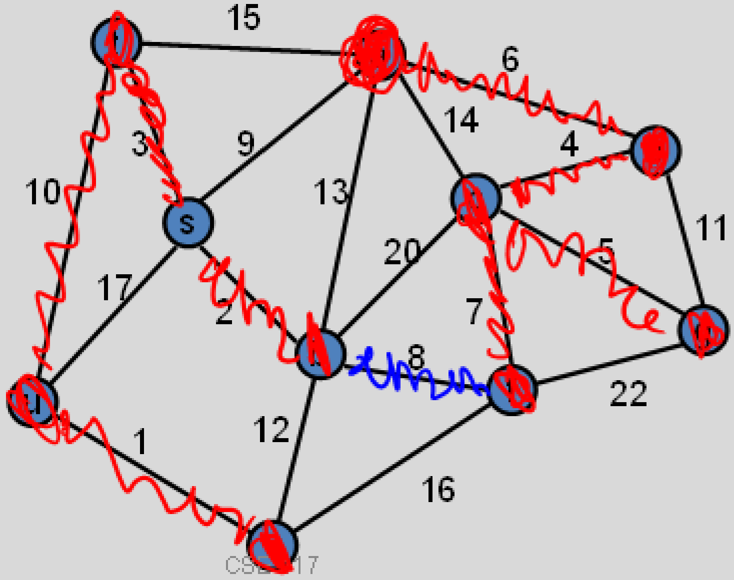


Greedy Algorithm 1
Prim's Algorithm

- Extend a tree by including the cheapest out going edge

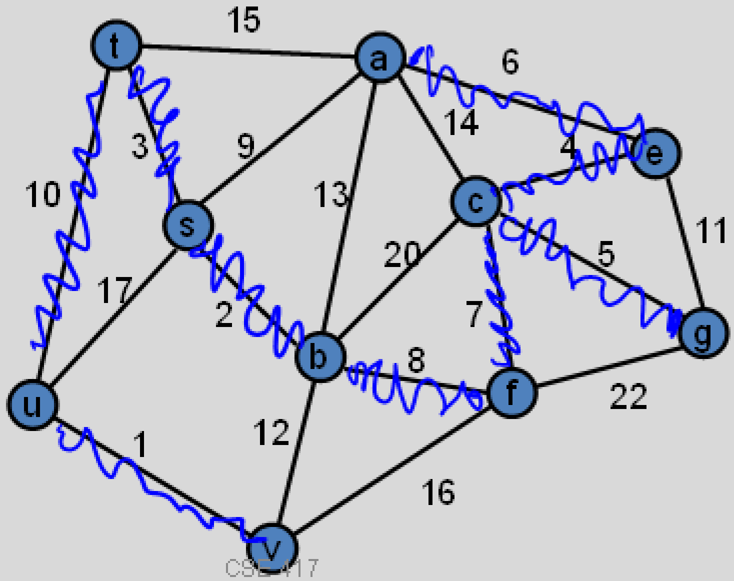
Construct the MST with Prim's algorithm starting from vertex a

Label the edges in order of insertion



Greedy Algorithm 2
Kruskal's Algorithm

- Add the cheapest edge that joins disjoint components

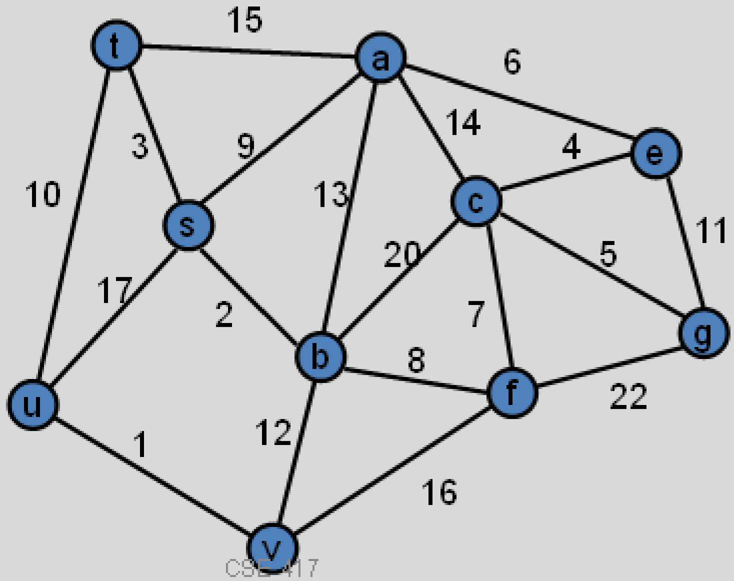


Construct the MST with Kruskal's algorithm

Label the edges in order of insertion

Greedy Algorithm 3
Reverse-Delete Algorithm

- Delete the most expensive edge that does not disconnect the graph



Construct the MST with the reverse-delete algorithm

Label the edges in order of removal

Prim

Dijkstra's Algorithm for Minimum Spanning Trees

S = { }; d[s] = 0; d[v] = infinity for v != s

While S != V

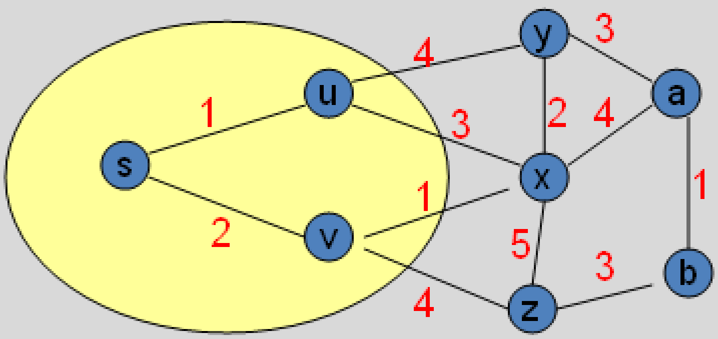
Choose v in V-S with minimum d[v]

Add v to S

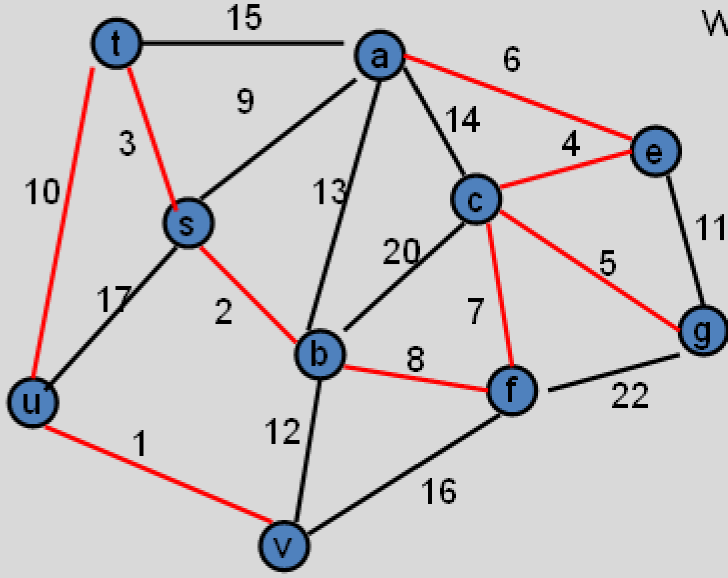
For each w in the neighborhood of v

d[w] = min(d[w], c(v, w))

Runtime m log n



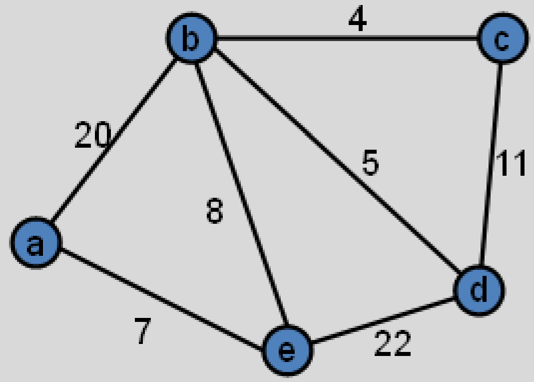
Minimum Spanning Tree



Undirected Graph G=(V,E) with edge weights

Greedy Algorithms for Minimum Spanning Tree

- [Prim] Extend a tree by including the cheapest out going edge
- [Kruskal] Add the cheapest edge that joins disjoint components
- [ReverseDelete] Delete the most expensive edge that does not disconnect the graph

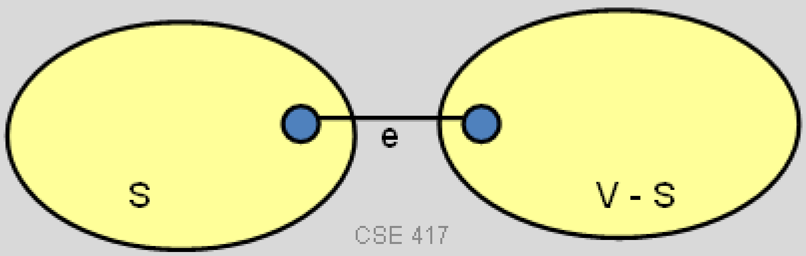


Why do the greedy algorithms work?

- For simplicity, assume all edge costs are distinct

Edge inclusion lemma

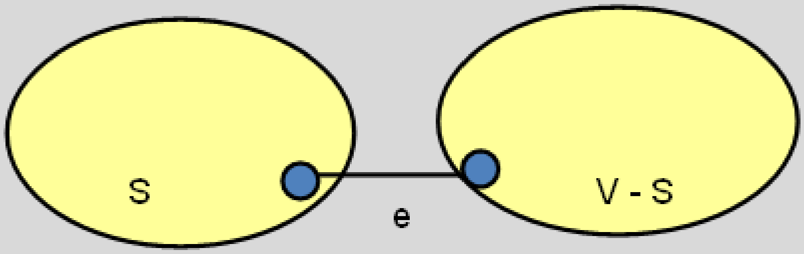
- Let S be a subset of V, and suppose e = (u, v) is the minimum cost edge of E, with u in S and v in V-S
- e is in every minimum spanning tree of G
  - Or equivalently, if e is not in T, then T is not a minimum spanning tree



e is the minimum cost edge between S and V-S

Proof

- Suppose T is a spanning tree that does not contain e
- Add e to T, this creates a cycle
- The cycle must have some edge $e_1 = (u_1, v_1)$ with u_1 in S and v_1 in V-S



- $T_1 = T - \{e_1\} + \{e\}$ is a spanning tree with lower cost
- Hence, T is not a minimum spanning tree