

Lecture26



CSE 417

Algorithms and Complexity

Autumn 2023

Lecture 26

Network Flow Applications

Announcements

- Homework 9
- Exam practice problems on course homepage
- Final Exam: Monday, December 11, 8:30 AM
– One Hour Fifty Minutes

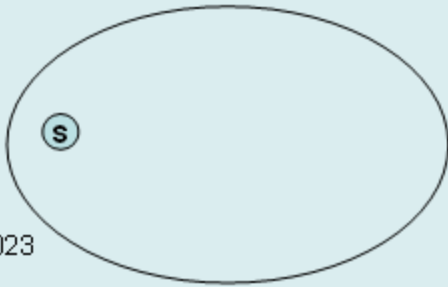
Fri, Dec 1	Net Flow Applications
Mon, Dec 4	Net Flow Applications + NP-Completeness
Wed, Dec 6	NP-Completeness
Fri, Dec 8	NP-Completeness
Mon, Dec 11	Final Exam

Outline

- ~~Network flow definitions~~
- ~~Flow examples~~
- ~~Augmenting Paths~~
- ~~Residual Graph~~
- ~~Ford Fulkerson Algorithm~~
- ~~Cuts~~
- ~~Maxflow-MinCut Theorem~~
- ~~Maxflow Algorithms~~
- Simple applications of Max Flow
- Non-simple applications of Max Flow

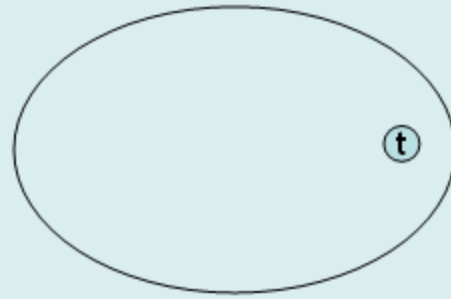
Cuts in a graph

- Cut: Partition of V into disjoint sets S , T with s in S and t in T .
- $\text{Cap}(S, T)$: sum of the capacities of edges from S to T
- Problem: Find the s - t Cut with minimum capacity



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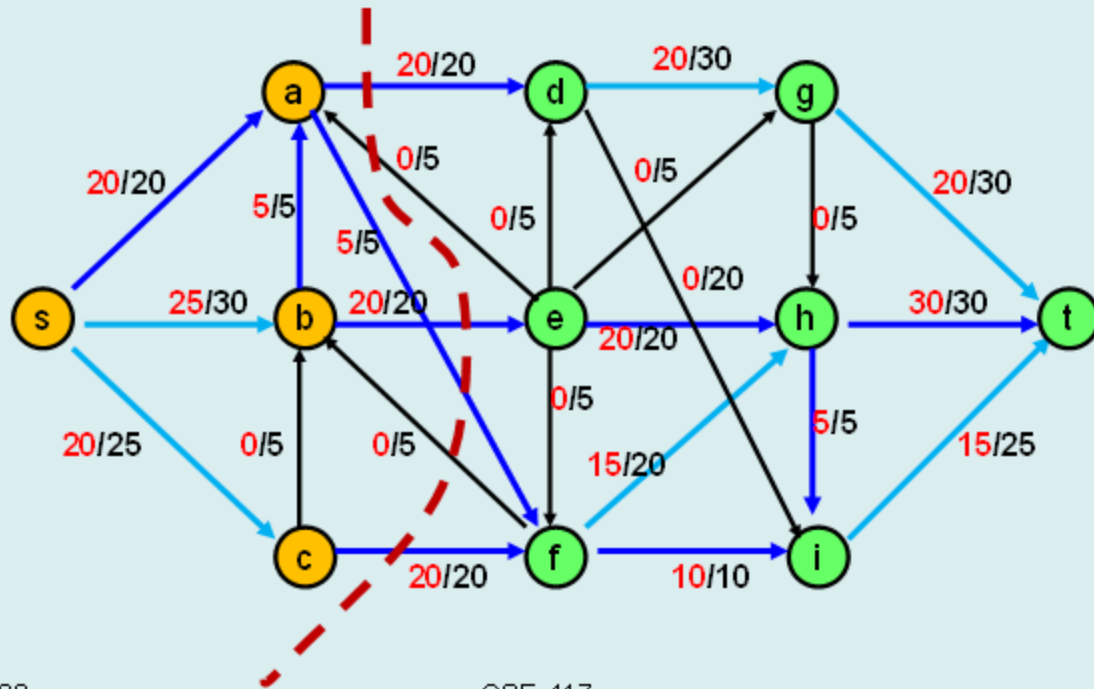
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Review

Max Flow / Min Cut



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Max Flow - Min Cut Theorem

- There exists a cut S, T such that
$$\text{Flow}(S, T) = \text{Cap}(S, T)$$
- Proof also shows that Ford Fulkerson algorithm finds a maximum flow

Ford Fulkerson Runtime

- Cost per phase \times number of phases
- Phases
 - Capacity leaving source: C
 - Add at least one unit per phase
- Cost per phase
 - Build residual graph: $O(m)$
 - Find s-t path in residual: $O(m)$

Network flow performance

- **Ford-Fulkerson algorithm**
 - $O(mC)$
- **Find the maximum capacity augmenting path**
 - $O(m^2 \log(C))$ time algorithm for network flow
- **Find the shortest augmenting path**
 - $O(m^2n)$ time algorithm for network flow
- **Find a blocking flow in the residual graph**
 - $O(mn \log n)$ time algorithm for network flow
- **Interior Point Methods**
 - $O(m + n)$

Problem Reduction

- ~~Reduce~~ ^{Convert} Problem A to Problem B
 - Convert an instance of Problem A to an instance of Problem B
 - Use a solution of Problem B to get a solution to Problem A
- Practical
 - Use a program for Problem B to solve Problem A
- Theoretical
 - Show that Problem B is at least as hard as Problem A

Problem Reduction Examples

- Reduce the problem of finding the Maximum of a set of integers to finding the Minimum of a set of integers

Find the maximum of: 8, -3, 2, 12, 1, -6

$-8, 3, -2, -12, -1, 6$
 -12
 12

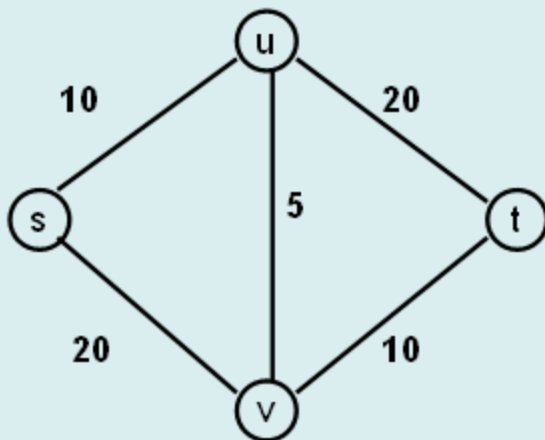
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Construct an equivalent minimization problem

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Undirected Network Flow

- Undirected graph with edge capacities
- Flow may go either direction along the edges (subject to the capacity constraints)



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Construct an equivalent flow problem

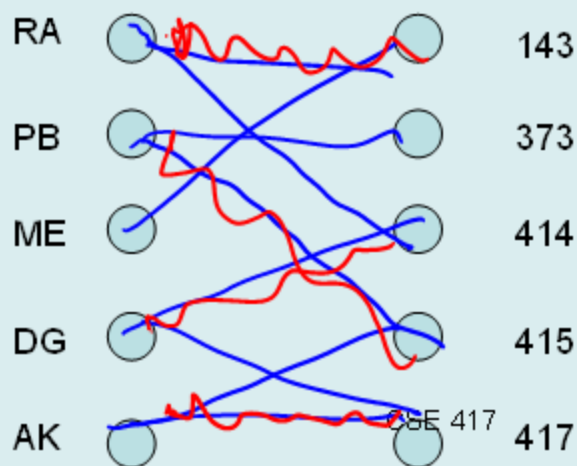
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Bipartite Matching

- A graph $G=(V,E)$ is bipartite if the vertices can be partitioned into disjoint sets X,Y
- A matching M is a subset of the edges that does not share any vertices
- Find a matching as large as possible

Application

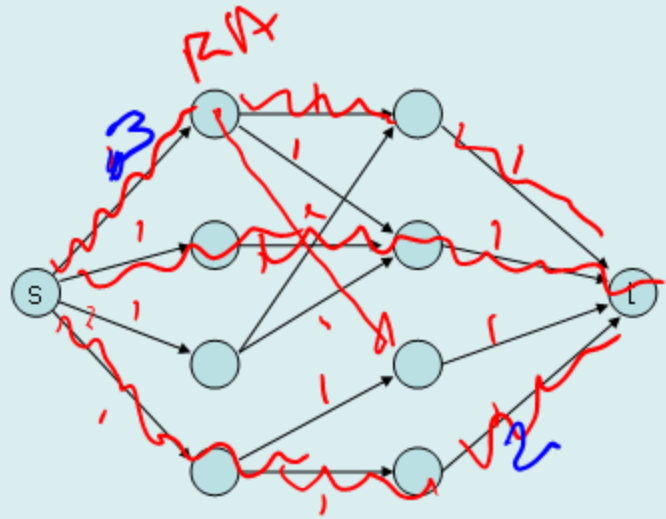
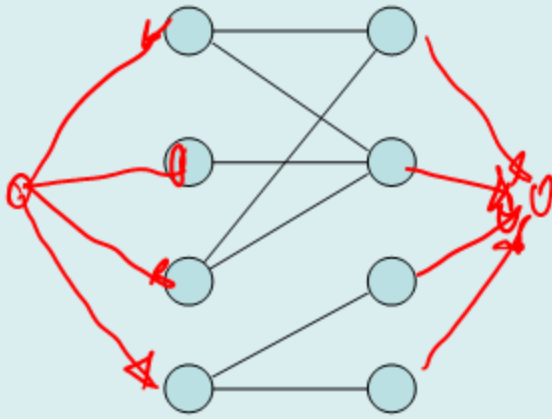
- A collection of teachers
- A collection of courses
- And a graph showing which teachers can teach which courses



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Converting Matching to Network Flow



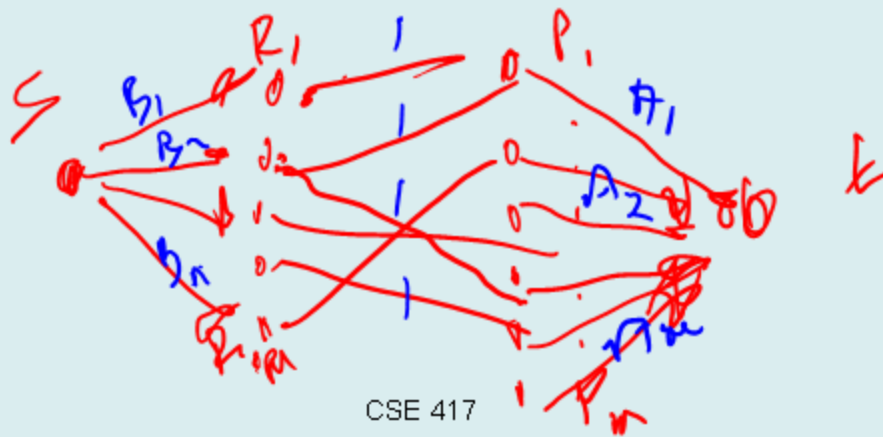
Multi-source network flow

- Multi-source network flow
 - Sources s_1, s_2, \dots, s_k
 - Sinks t_1, t_2, \dots, t_j
- Solve with Single source network flow



Resource Allocation: Assignment of reviewers

- A set of papers P_1, \dots, P_n
- A set of reviewers R_1, \dots, R_m
- Paper P_i requires A_i reviewers
- Reviewer R_j can review B_j papers
- For each reviewer R_j , there is a list of paper L_{j1}, \dots, L_{jk} that R_j is qualified to review



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Baseball elimination

- Can the Dinosaurs win the league?
- Remaining games:
 - AB, AC, AD, AD, AD, BC, BC, BC, BD, CD

	W	L
Ants	4	2
Bees	4	2
Cockroaches	3	3
Dinosaurs	6	5

A team **wins** the league if it has strictly more wins than any other team at the end of the season
 A team **ties** for first place if no team has more wins, and there is some other team with the same number of wins

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Baseball elimination

- Can the Fruit Flies win or tie the league?
- Remaining games:
 - AC, AD, AD, AD, AF,
 - BC, BC, BC, BC, BC,
 - BD, BE, BE, BE, BE,
 - BF, CE, CE, CE, CF,
 - CF, DE, DF, EF, EF

	W	L	
Ants	17	12	2
Bees	16	7	3
Cockroaches	16	7	3
Dinosaurs	14	13	5
Earthworms	14	10	5
Fruit Flies	12	15	

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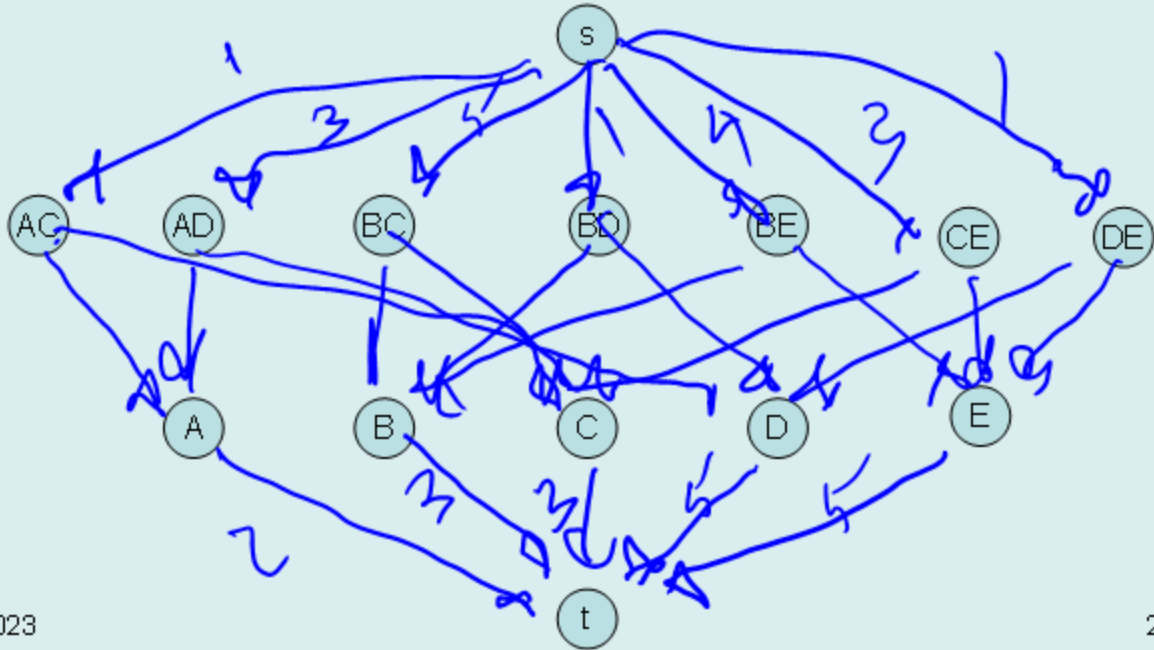
Assume Fruit Flies win remaining games

- Fruit Flies are tied for first place if no team wins more than 19 games
- Allowable wins
 - Ants (2)
 - Bees (3)
 - Cockroaches (3)
 - Dinosaurs (5)
 - Earthworms (5)
- 18 games to play
 - AC, AD, AD, AD, BC, BC, BC, BC, BC, BD, BE, BE, BE, BE, CE, CE, CE, DE

	W	L
Ants	17	13
Bees	16	8
Cockroaches	16	9
Dinosaurs	14	14
Earthworms	14	12
Fruit Flies	19	15

Remaining games

AC, AD, AD, AD, BC, BC, BC, BC, BC, BD, BE, BE, BE, BE, CE, CE, CE, DE



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Minimum Cut Applications

- Image Segmentation
- Open Pit Mining / Task Selection Problem
- Reduction to Min Cut problem

S, T is a cut if S, T is a partition of the vertices with s in S and t in T

The capacity of an S, T cut is the sum of the capacities of all edges going from S to T

Image Segmentation

- Separate foreground from background



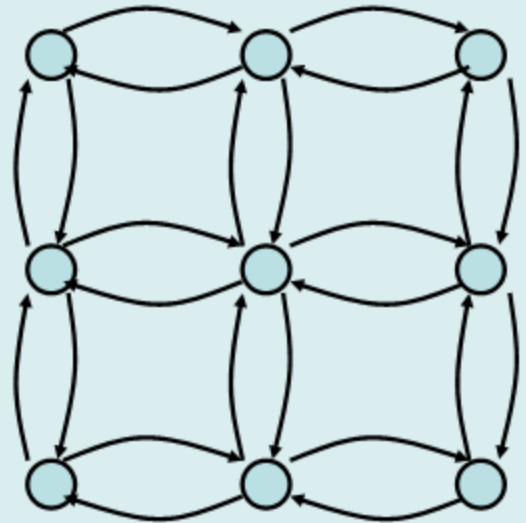
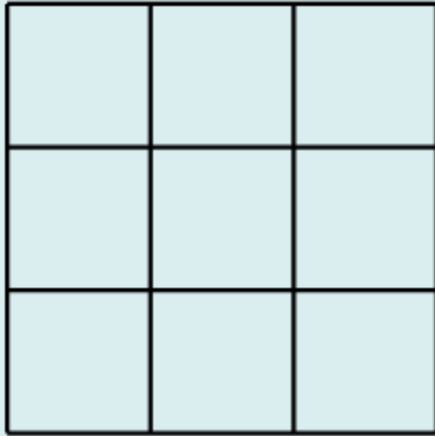


Image analysis

- a_i : value of assigning pixel i to the foreground
- b_j : value of assigning pixel j to the background
- p_{ij} : penalty for assigning i to the foreground, j to the background or vice versa
- A : foreground, B : background
- $Q(A, B) = \sum_{\{i \text{ in } A\}} a_i + \sum_{\{j \text{ in } B\}} b_j - \sum_{\{(i,j) \text{ in } E, i \text{ in } A, j \text{ in } B\}} p_{ij}$

Pixel graph to flow graph

s



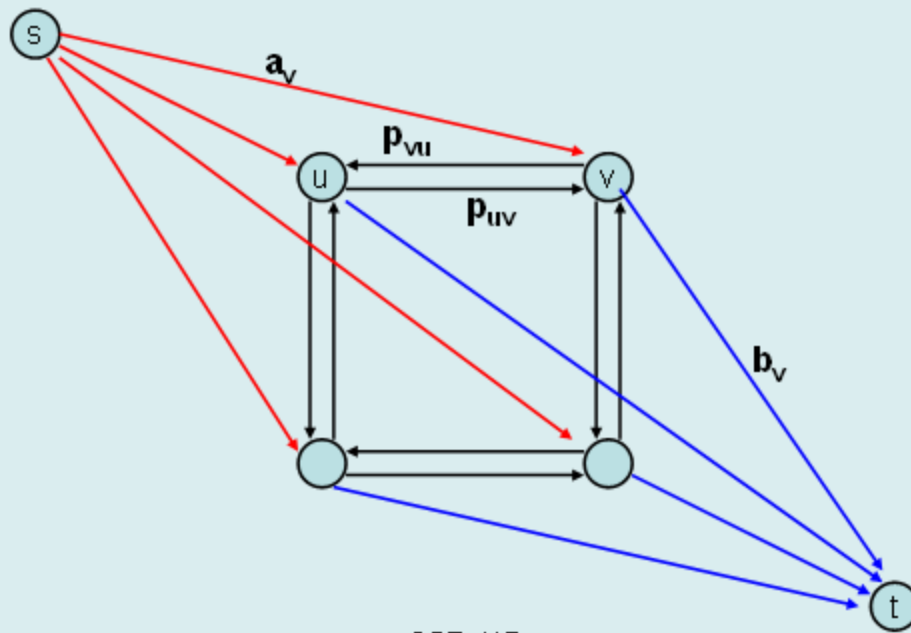
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Mincut Construction



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