

CSE 417

Algorithms and Complexity

Autumn 2023

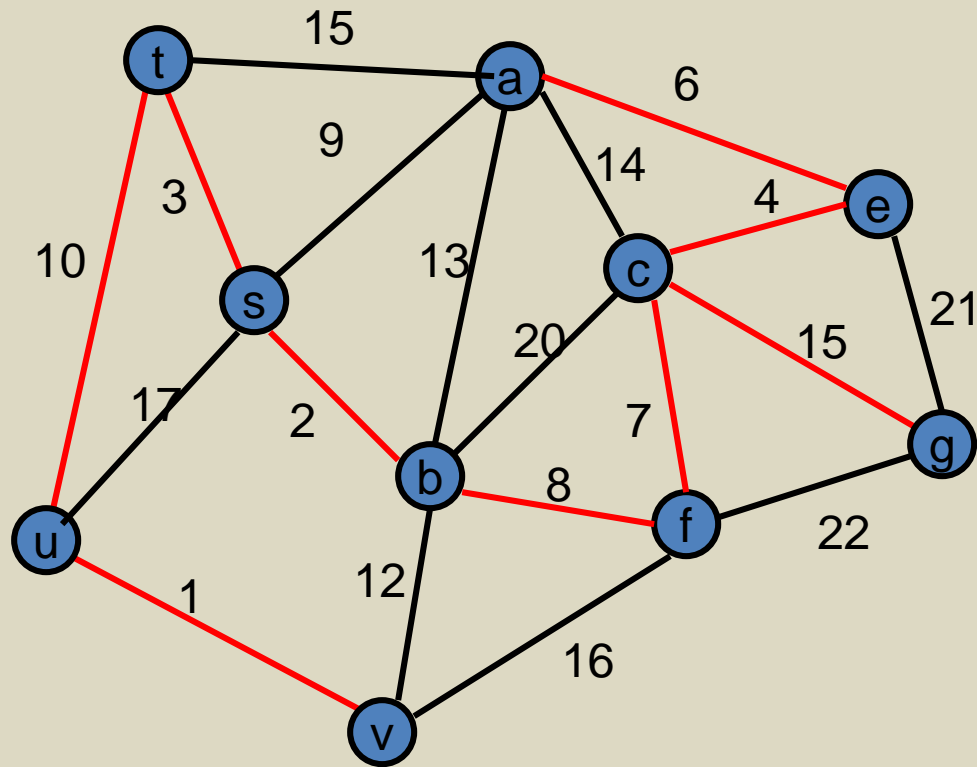
Lecture 14

Finishing Minimum Spanning Trees

Announcements

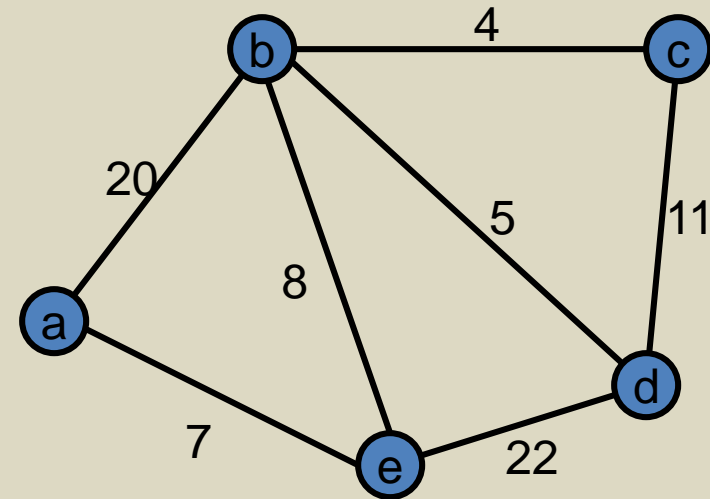
- Midterm, Monday, October 30
 - Closed book, closed notes, no calculators
 - Time limit: 50 minutes
 - Answer the problems on the exam paper.
 - If you need extra space use the overflow page (and the back of a page if necessary)
 - Problems are not of equal difficulty, if you get stuck on a problem, move on.
 - ``Justify your answer'' means give a short and convincing explanation. Depending on the situation, justifications can involve counter examples, or cite results established in the text or in lecture.
- Homework 5, Available, Due November 3
 - Homework 5 is OPTIONAL

Minimum Spanning Tree



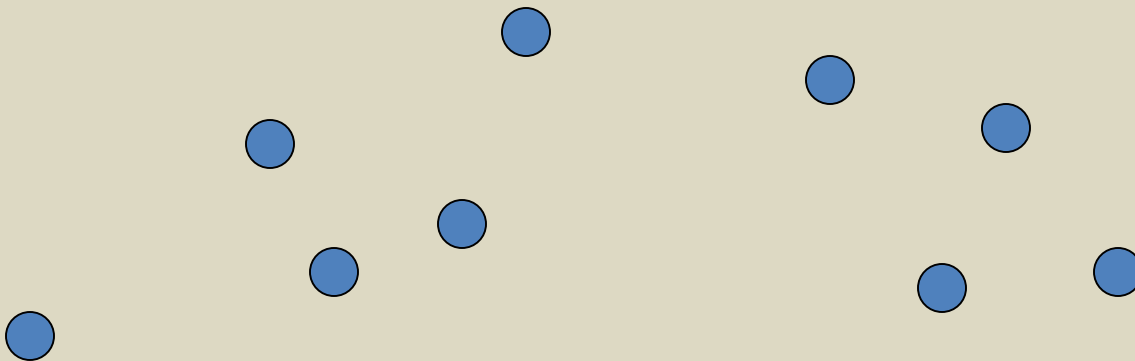
Greedy Algorithms for Minimum Spanning Tree

- Prim's Algorithm:
Extend a tree by including the cheapest out going edge
- Kruskal's Algorithm:
Add the cheapest edge that joins disjoint components



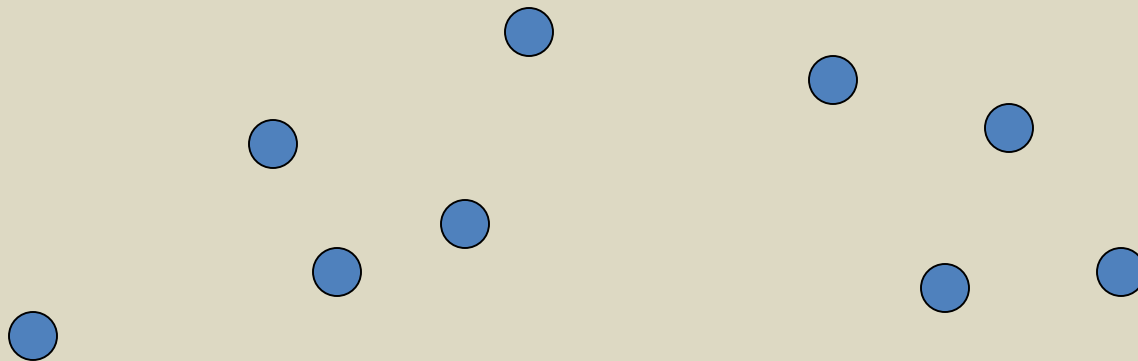
Application: Clustering

- Given a collection of points in an r -dimensional space and an integer K , divide the points into K sets that are closest together

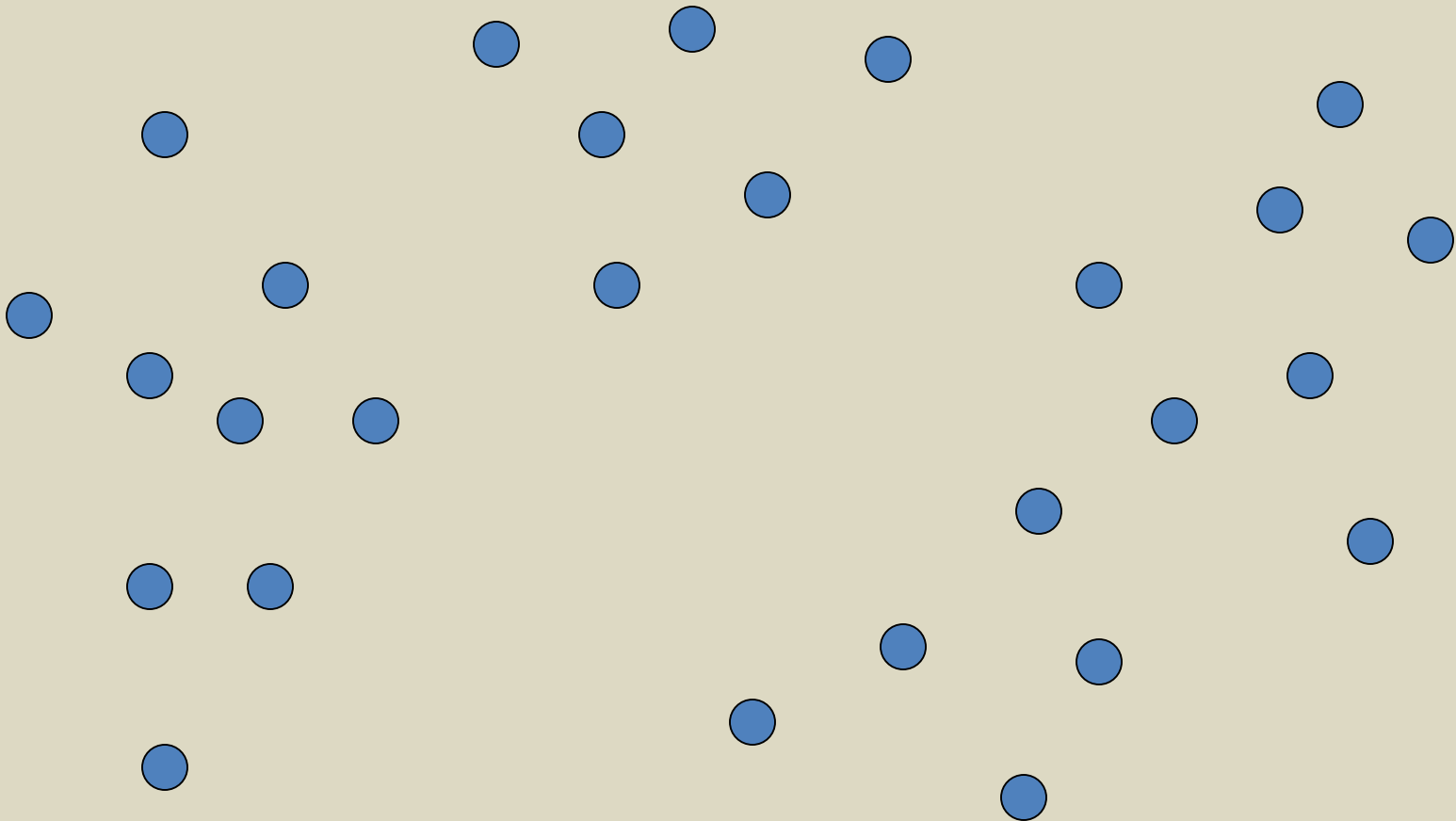


Distance clustering

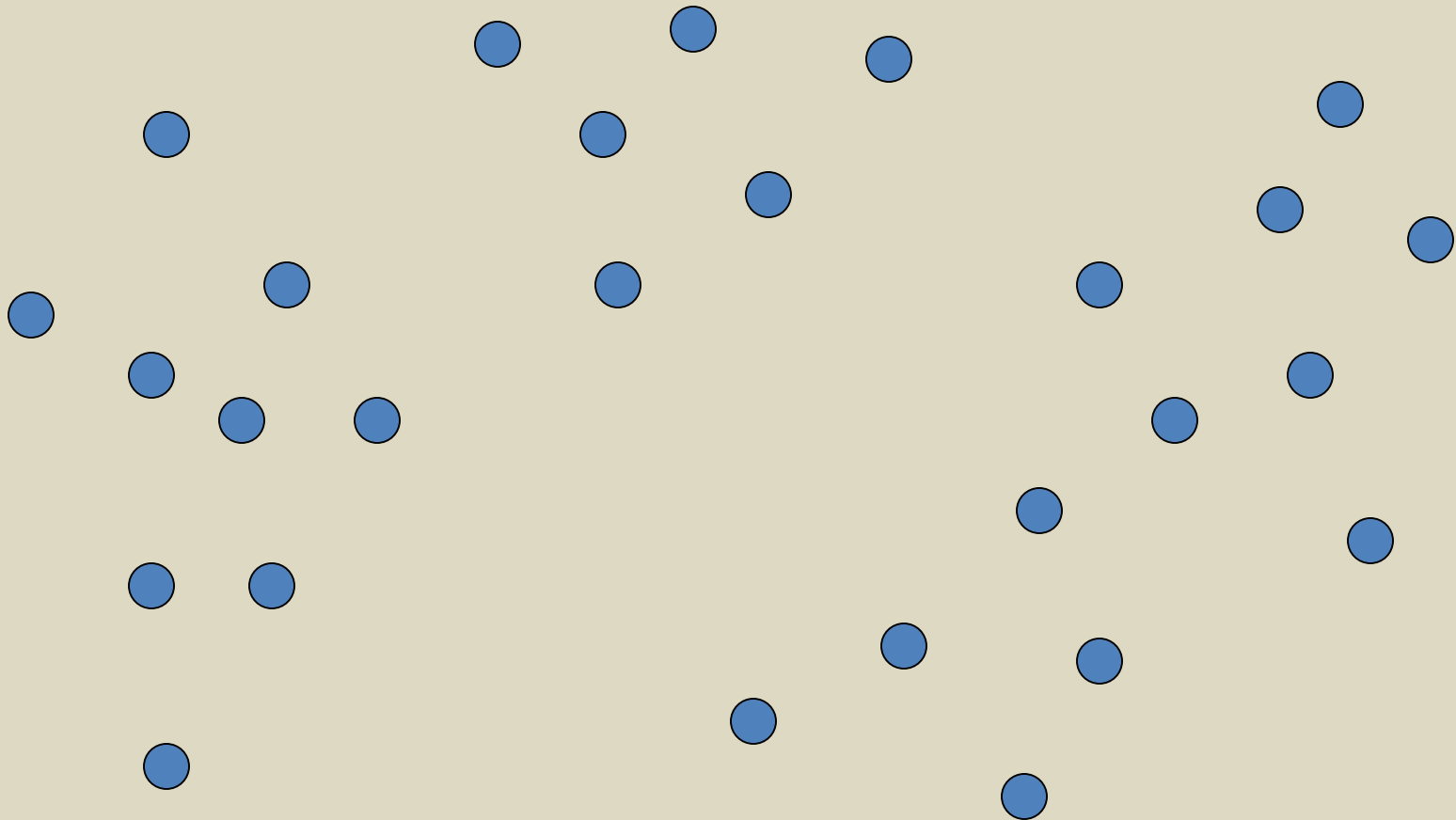
- Divide the data set into K subsets to maximize the distance between any pair of sets
 - $\text{dist}(S_1, S_2) = \min \{ \text{dist}(x, y) \mid x \text{ in } S_1, y \text{ in } S_2 \}$



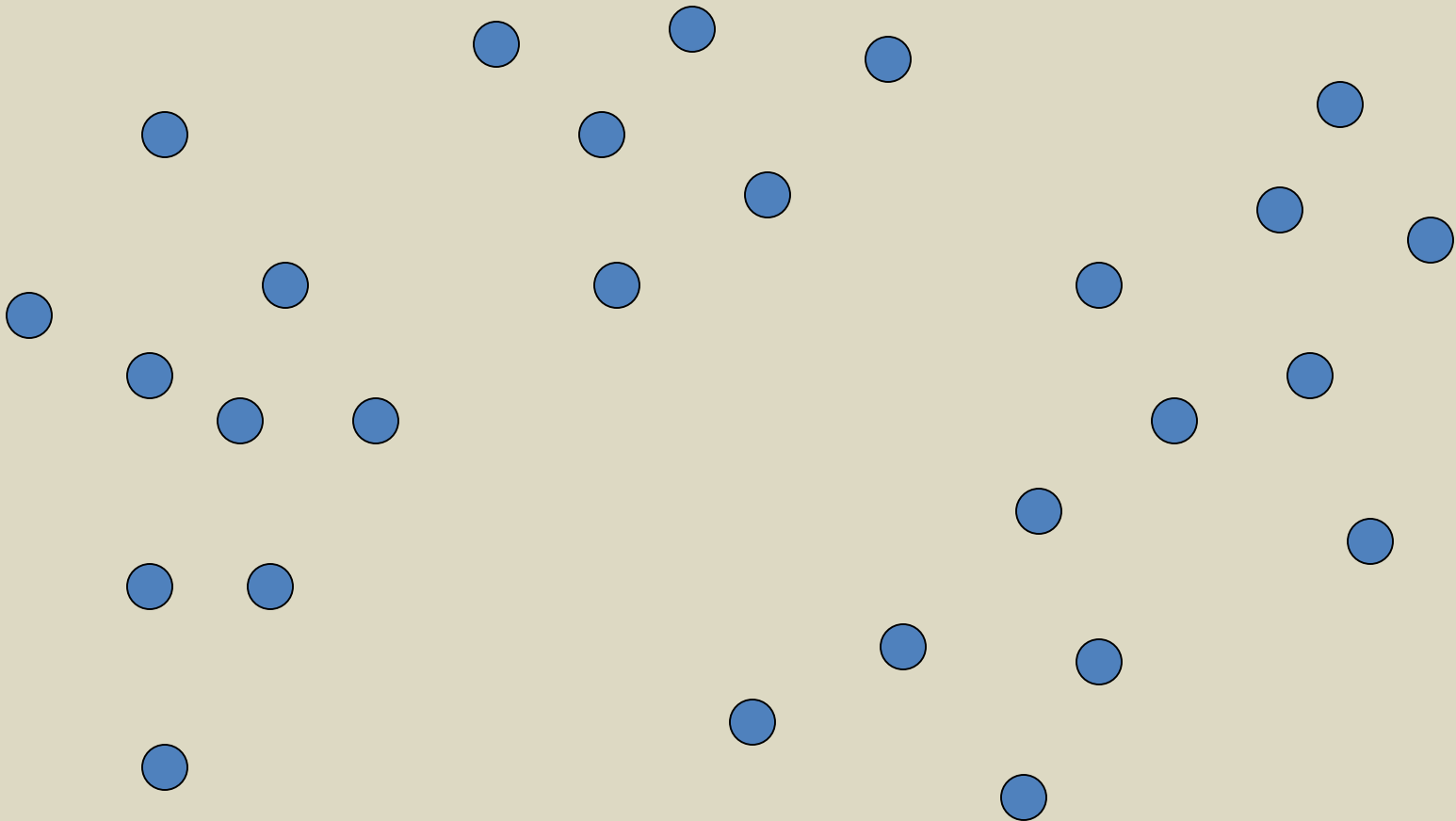
Divide into 2 clusters



Divide into 3 clusters



Divide into 4 clusters



Distance Clustering Algorithm

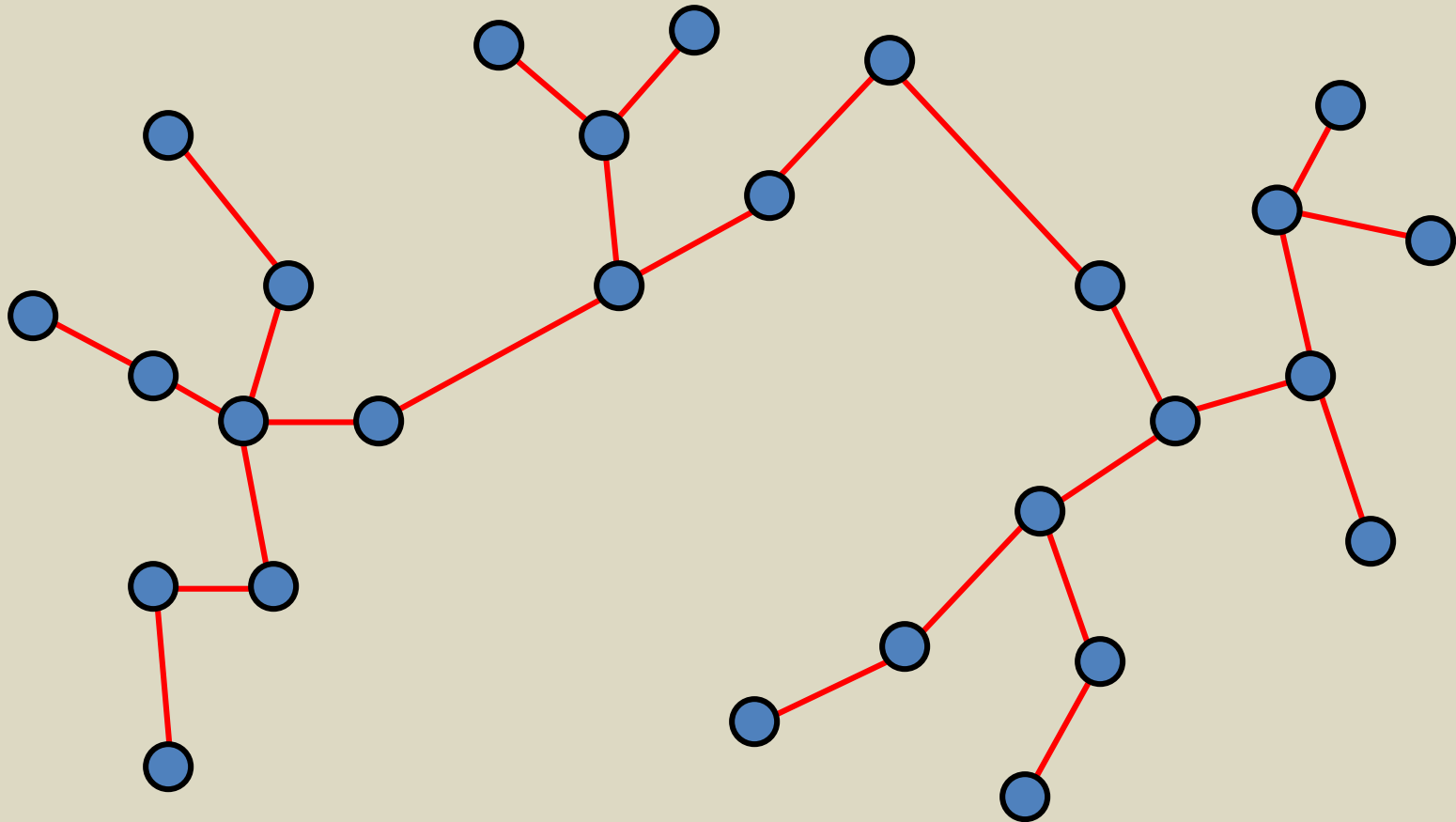
Let $C = \{\{v_1\}, \{v_2\}, \dots, \{v_n\}\}$; $T = \{ \}$

while $|C| > K$

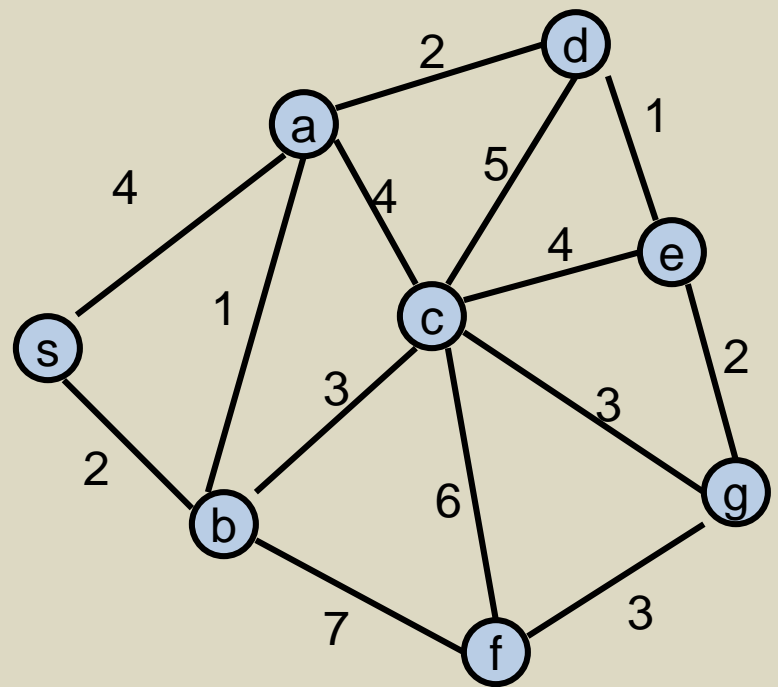
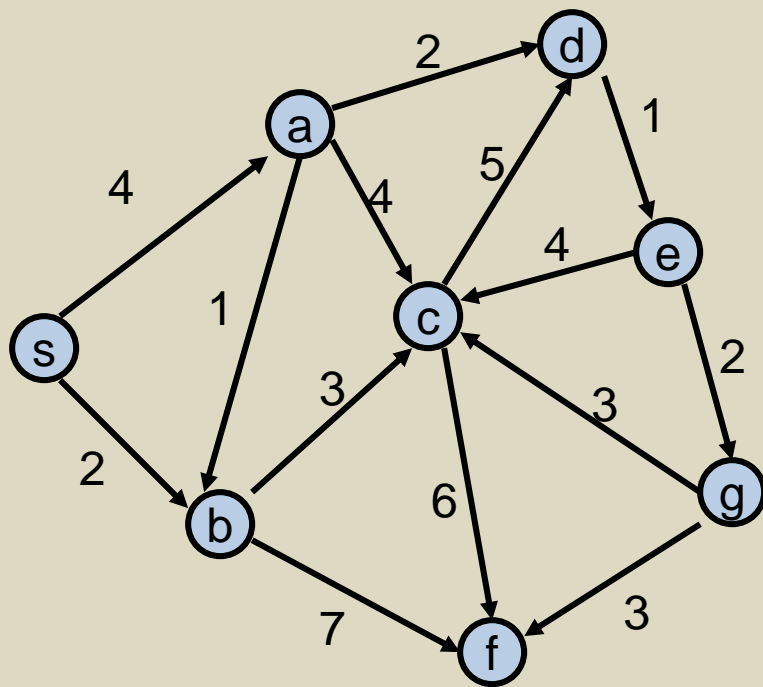
Let $e = (u, v)$ with u in C_i and v in C_j be the minimum cost edge joining distinct sets in C

Replace C_i and C_j by $C_i \cup C_j$

K-clustering

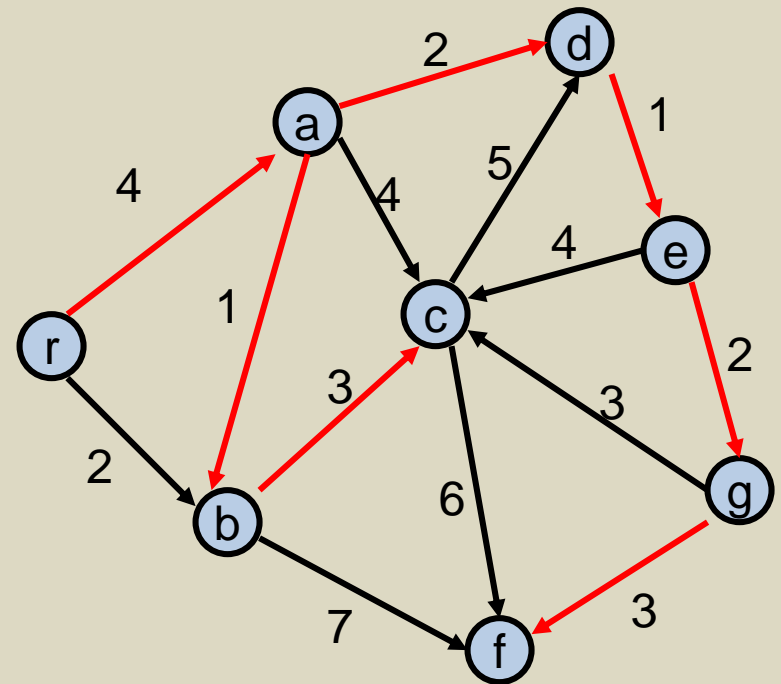
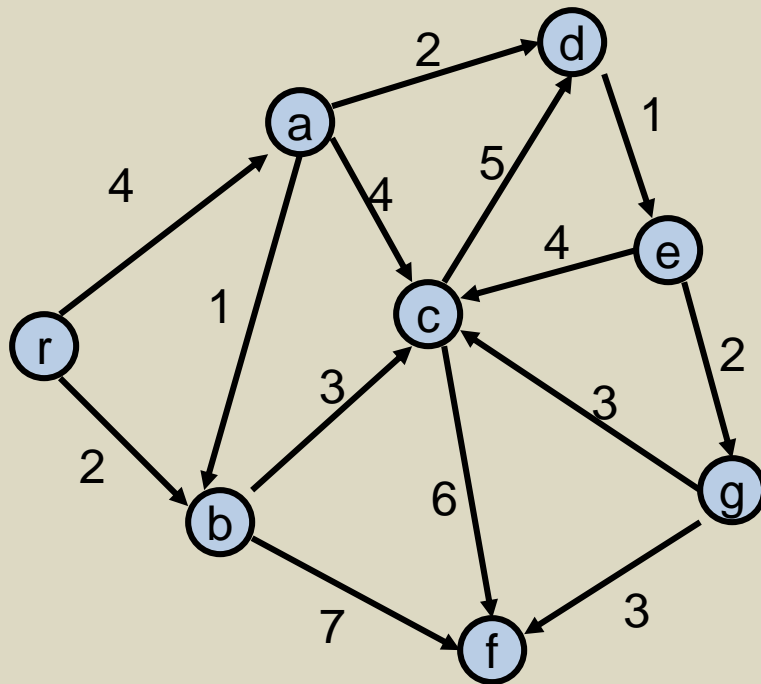


Shortest paths in directed graphs vs undirected graphs



What about the minimum spanning tree of a directed graph?

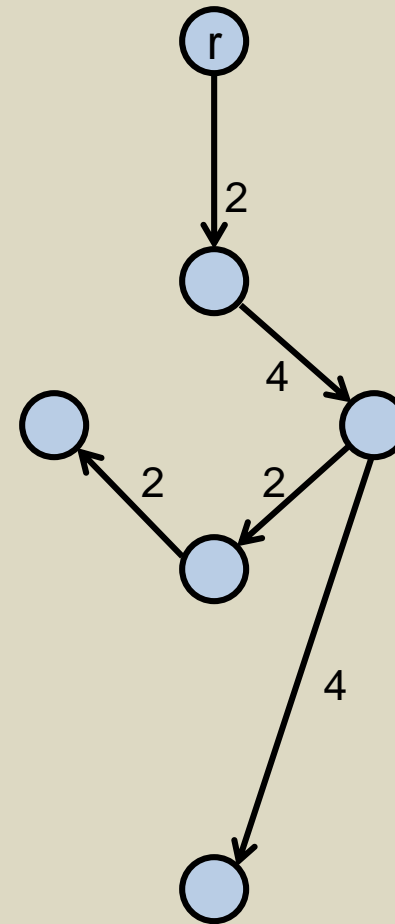
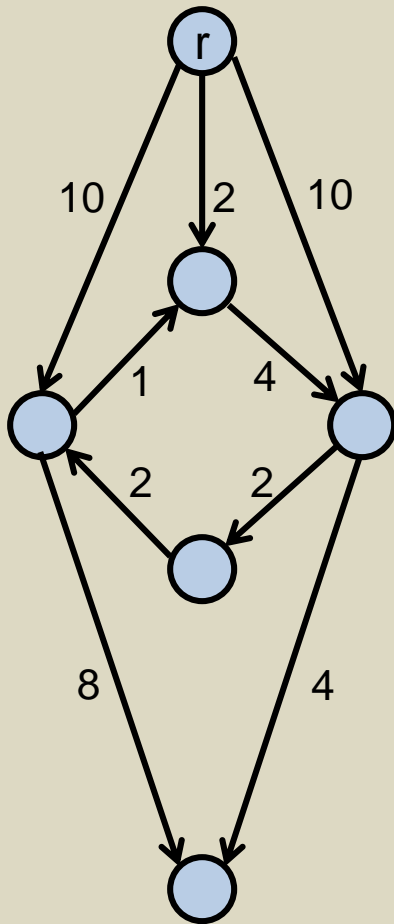
- Must specify the root r
- Branching: Out tree with root r



Assume all vertices reachable from r

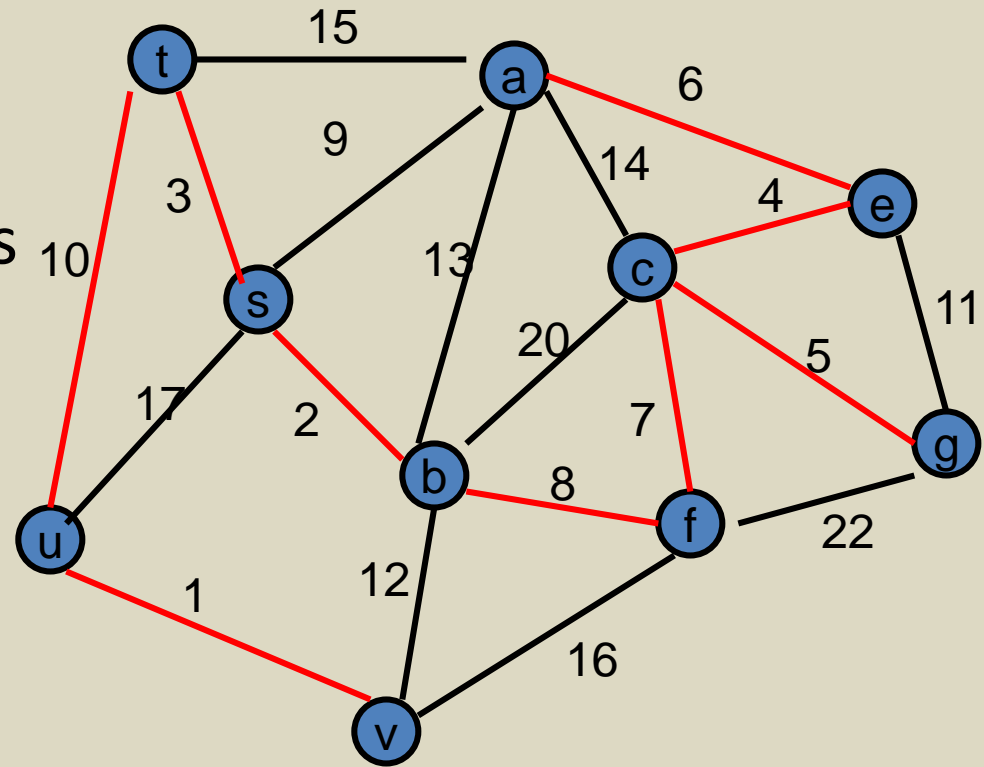
Also called an arborescence

Finding a minimum branching



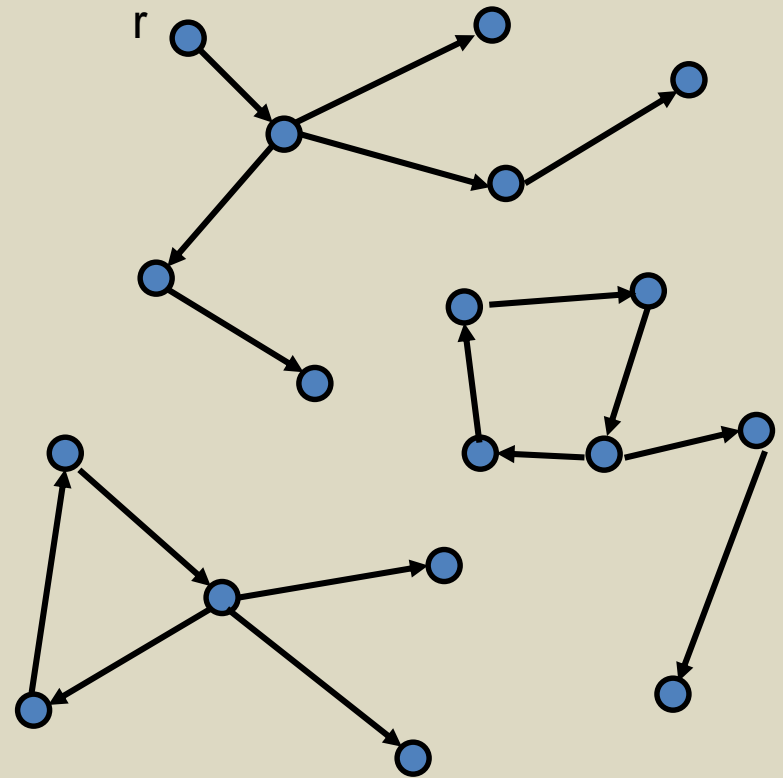
Another MST Algorithm

- Choose minimum cost edge into each vertex
- Merge into components
- Repeat until done



Idea for branching algorithm

- Select minimum cost edge going into each vertex
- If graph is a branching then done
- Otherwise collapse cycles and repeat

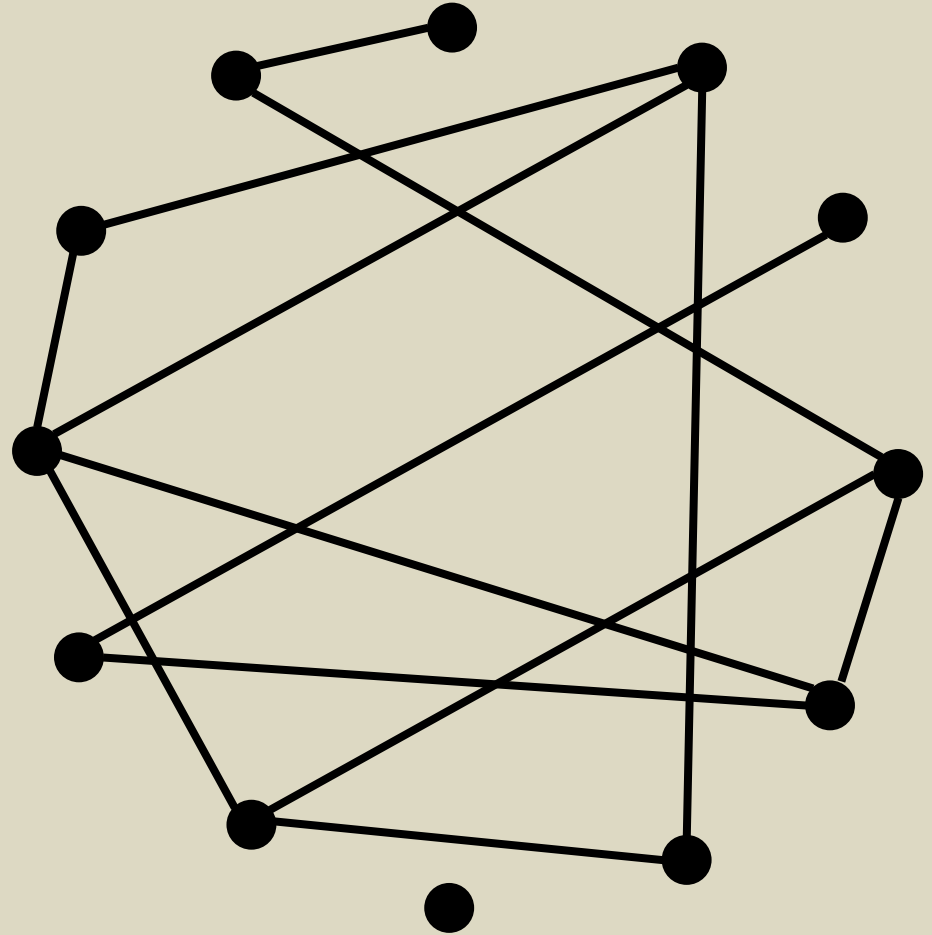


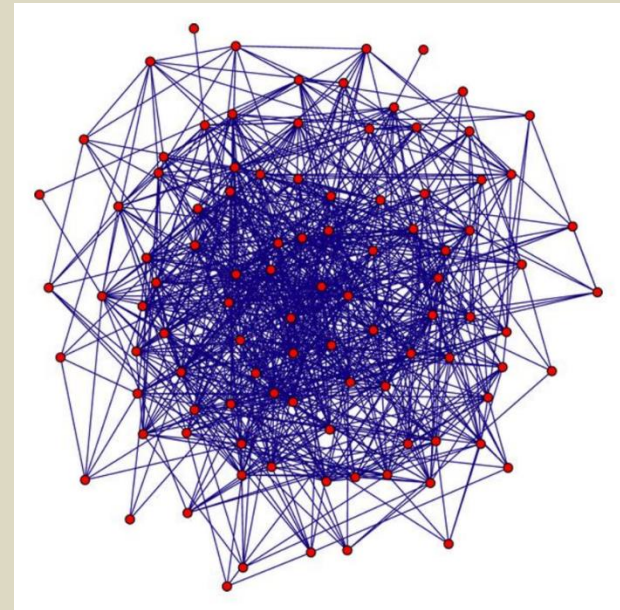
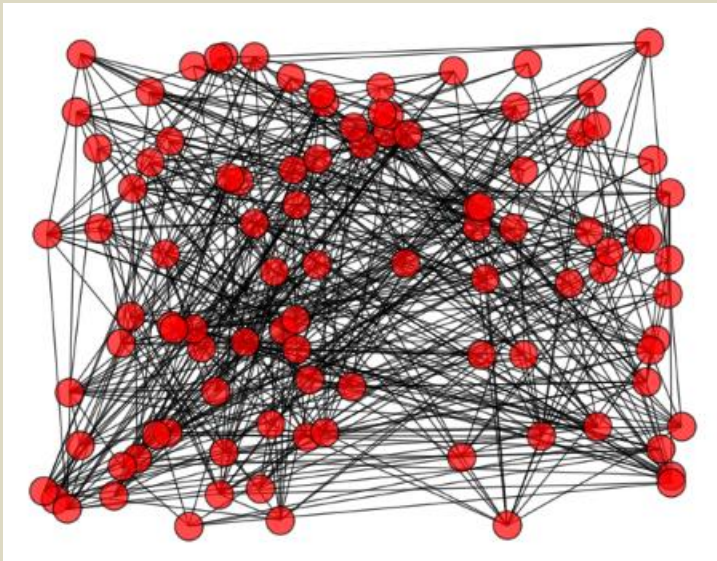
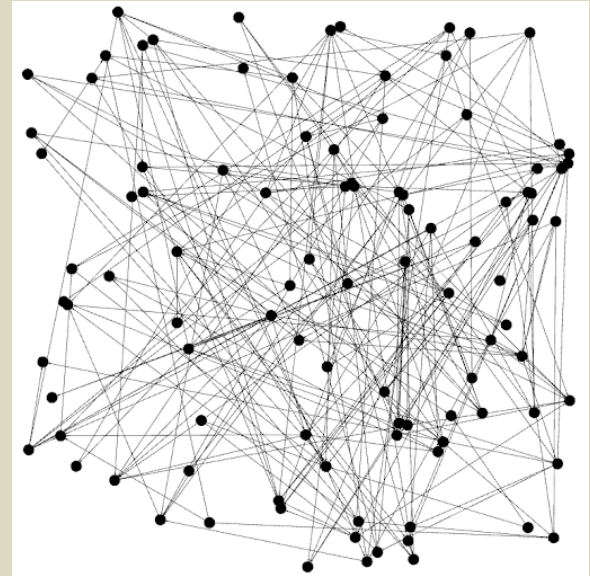
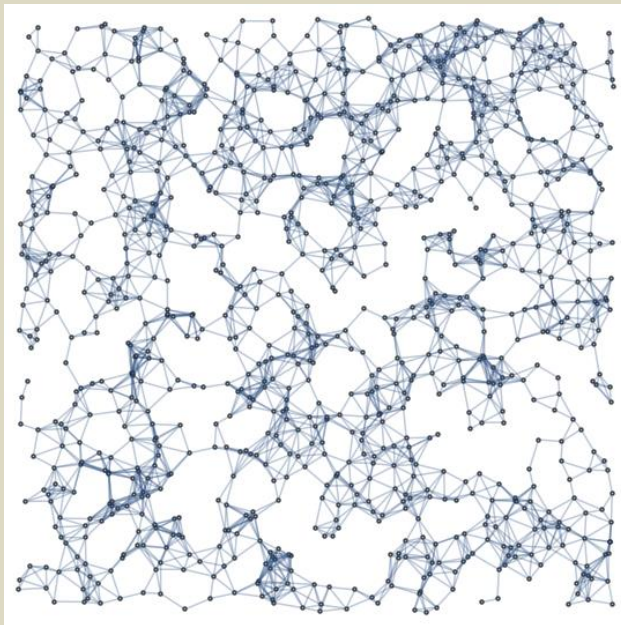
Homework 5: Create a program for coloring random graphs

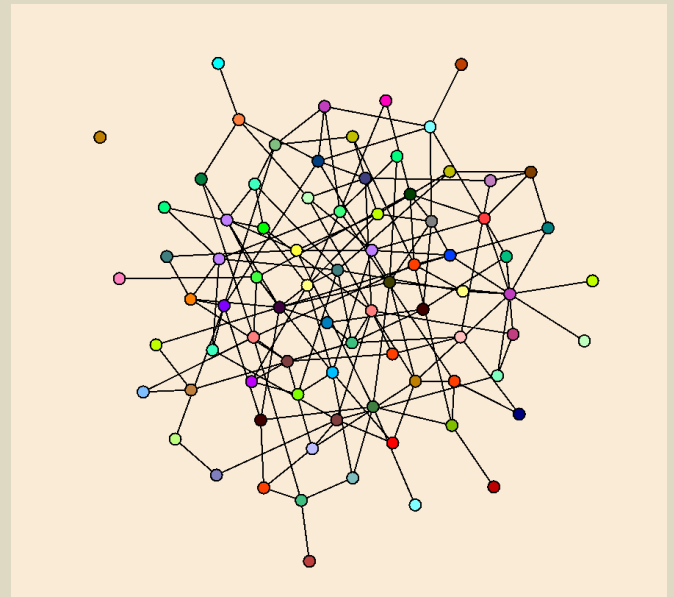
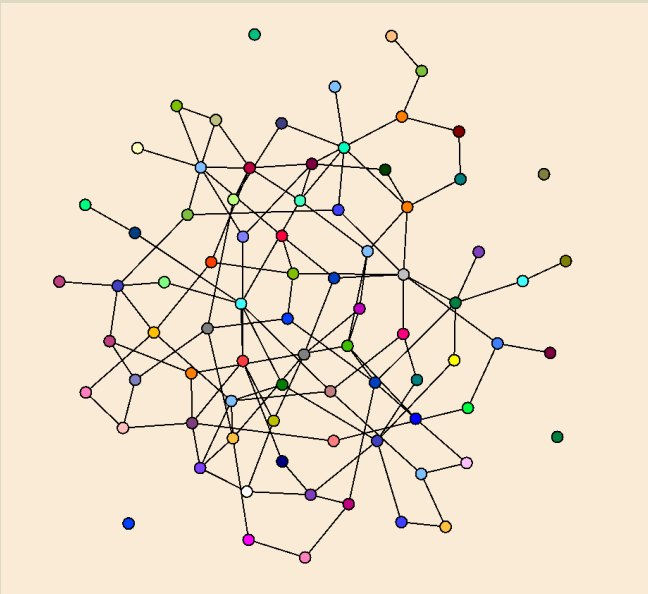
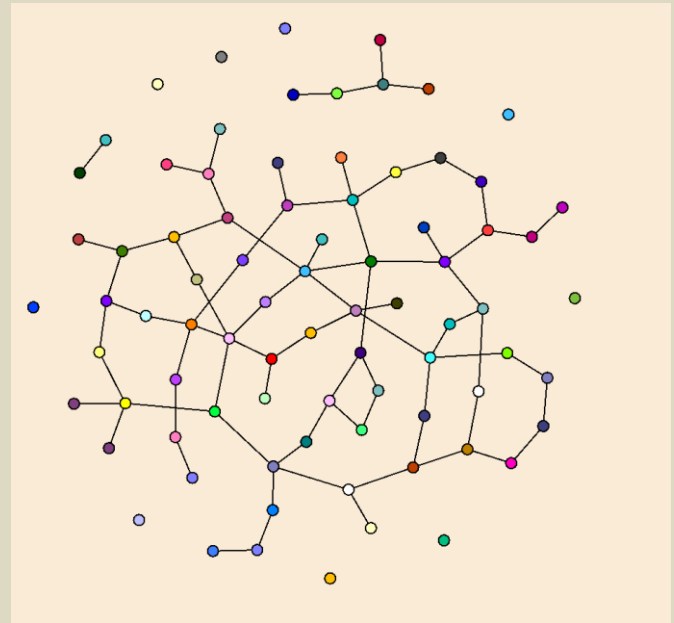
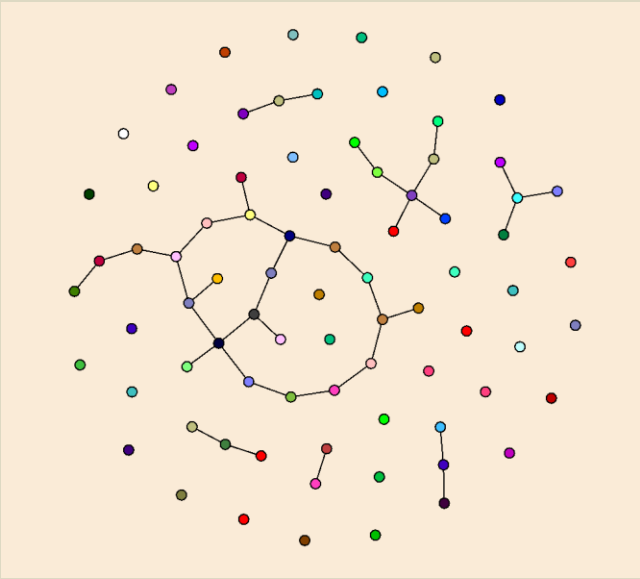
- Optional – this assignment is Just for fun!
- Problem 1: Generate Random Graphs
- Problem 2: Greedy coloring algorithm – first available color
- Problem 3: Low degree first and high degree first heuristics
- Problem 4: Can you find a better coloring algorithm

Random Graphs

- What is a random graph?
- Choose edges at random
- Interesting model of certain phenomena
- Mathematical study
- Useful inputs for graph algorithms







Model of Random Graphs

- Undirected Graphs

- Random Graph with n vertices and m edges, G_m
- Random Graph with n vertices where each edge has probability p , G_p
- Models are similar when $p = 2m / (n * (n - 1))$

```
for (int i = 0; i < n - 1; i++)  
    for (int j = i + 1; j < n; j++)  
        if (random.NextDouble() < p)  
            AddEdge(i, j);
```

Coloring Random Graphs

- Chromatic index of a graph G , $\chi(G)$ – minimum number of colors needed to color G
- Mathematical question, given a graph g chosen at random from $G_p(n)$, what is the expected value of $\chi(g)$
- Graph coloring is NP complete – suggesting that it may be hard to determine $\chi(g)$
- There is a fairly large gap between the heuristic results and the theoretical value
- There are a number of ties with theoretical physics