

Announcements

· Course website:

//courses.cs.washington.edu/courses/cse417/23au/

- Homework Due Friday
- · Office Hours:

Richard Anderson, Monday 2-3 pm, Thursday, 4-5 pm Megh Bhalerao, Friday, 4:30-6:30 pm Tieman Kennedy, Wednesday, 9-10 am, Friday, 9-10 am Yigao Li, Tuesday, 11:30am-12:30 pm, Friday, 1-2 pm Kaiyuan Liu, Tuesday, 3-4 pm, Thursday, 2-3 pm Sravani Nanduri, Monday, 4:30-5:30 pm, Friday, 11:30am-12:30pm Albert Weng, Monday, 3:30-4:30 pm, Wednesday, 3:30-4:30 pm

Theory of Algorithms

- · What is expertise?
- · How do experts differ from novices?

Introduction of five problems

- Show the types of problems wewill be considering in the class
- Examples of important types of problems
- Similar looking problems with very different characteristics
- Problems
 - Scheduling
 - Weighted Scheduling
 - Bipartite Matching

 - Maximum Independent SetCompetitive Facility Location

What is a problem?

- Instance
- Solution
- · Constraints on solution
- Measure of value

Problem: Scheduling

- · Suppose that you own a banquet hall
- · You have a series of requests for use of the hall: $(s_1, f_1), (s_2, f_2), \dots$

· Find a set of requests as large as possible with no overlap

What is the largest solution?					

Greedy Algorithm

- Test elements one at a time if they can be members of the solution
- If an element is not ruled out by earlier choices, add it to the solution
- Many possible choices for ordering (length, start time, end time)
- For this problem, considering the jobs by increasing end time works

Suppose we add values?

- (s_i, f_i, v_i), start time, finish time, payment
- · Maximize value of elements in the solution

5				<u> </u>
1	2		4	1
	3	1		6

Greedy Algorithms

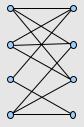
- · Earliest finish time
- Maximum value
- Give counter examples to show these algorithms don't find the maximum value solution

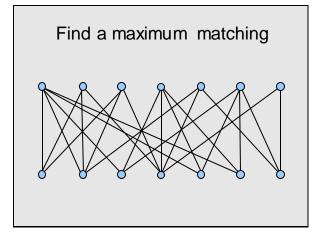
Dynamic Programming

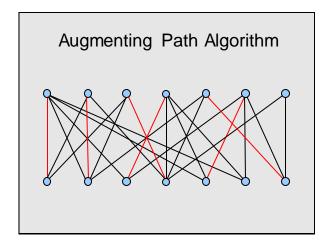
- Requests R₁, R₂, R₃, . . .
- Assume requests are in increasing order of finish time (f₁ < f₂ < f₃ . . .)
- Opt_i is the maximum value solution of $\{R_1, R_2, \ldots, R_i\}$ containing R_i
- Opt_i = Max{ j | $f_i < s_i$ }[Opt_j + v_i]

Matching (Combinatorial Optimization)

- Given a bipartite graph G=(U,V,E), find a subset of the edges M of maximum size with no common endpoints.
- · Application:
 - U: Professors
 - V: Courses
 - (u,v) in E if Prof. u can teach course v

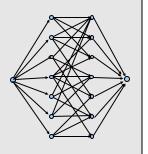






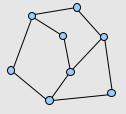
Reduction to network flow

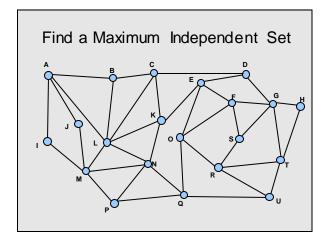
- · More general problem
- Send flow from source to sink
- Flow subject to capacities at edges
- Flow conserved at vertices
- Can solve matching as a flow problem

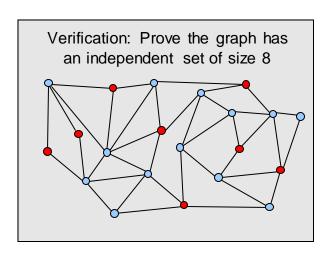


Maximum Independent Set

- Given an undirected graph G=(V,E), find a set I of vertices such that there are no edges between vertices of I
- Find a set I as large as possible







Key characteristic

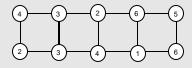
- Hard to find a solution
- Easy to verify a solution once you have one
- · Other problems like this
 - Hamiltonian circuit
 - Clique
 - Subset sum
 - Graph coloring

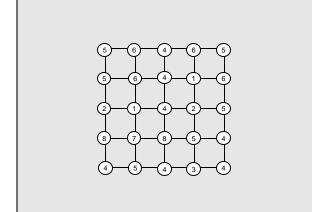
NP-Completeness

- · Theory of Hard Problems
- A large number of problems are known to be equivalent
- · Very elegant theory

Are there even harder problems?

- · Simple game:
 - Players alternate selecting nodes in a graph
 - · Score points associated with node
 - · Remove nodes neighbors
 - When neither can move, player with most points wins





Competitive Facility Location

- · Choose location for a facility
 - Value associated with placement
 - Restriction on placing facilities too close together
- Competitive
 - Different companies place facilities
 - E.g., KFC and McDonald's

Complexity theory

- These problems are P-Space complete instead of NP-Complete
 - Appear to be much harder
 - No obvious certificate
 - G has a Maximum Independent Set of size 10
 - Player 1 winsby at least 10 points

Summary

- Scheduling
- Weighted Scheduling
- Bipartite Matching
- Maximum Independent Set
- Competitive Scheduling