CSE 417: Algorithms and Computational Complexity

Lecture I: Overview

Winter 2019

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University of Washington Computer Science & Engineering

CSE 417, Wi '19: Algorithms & Computational Complexity

CSE Home

Administrative

FAQ

Schedule & Reading

Course Email/BBoard

Subscription Options Class List Archive E-mail Course Staff **Google Groups BBoard**

Lecture Notes

1: Overview & Example

Lecture Recordings

Lecture: JHN 102 (room info) MWF1:30-2:20

> **Office Hours Phone** Location

2:30-3:30 CSE 554 (206) 543-6298 <--1/11 excepted Instructor: Larry Ruzzo, ruzzo@cs

TAs: 1:00-2:00 4th floor breakout Yuqing Ai, yuqingai@cs

> Daniel Jones, dcjones@cs **TBA** Saidutt Nimmagadda, nimmas@cs TBA

Alex Okeson, amokeson@cs Tu 2:30-3:30 CSE 021

t homework. should change their

Efficient algorithms for manipulating graphs and Turing machines. Time and space complexity. NP-

http://courses.cs.washington.edu/417 Homework, Midterm, Final. Homework will be a mix of paper & pencil exercises and programing. Overall weights 55%, 15%, 30%, roughly.

> **Late Policy:** Papers and/or electronic turnins are due at the **start** of class on the due date. 10% off for up to one day late; additional 15% per day thereafter. (Day = calendar day, i.e., Sunday is later than Saturday.)

> Textbooks: Algorithm Design by Jon Kleinberg and Eva Tardos. Addison Wesley, 2006. (Available from U Book Store,

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Open "mailman.u.washington.edu/mailman/private/cse417a_wi19" in a new tab

What you'll have to do

Homework (~55% of grade)

Programming

Several small projects

Written homework assignments

English exposition and pseudo-code

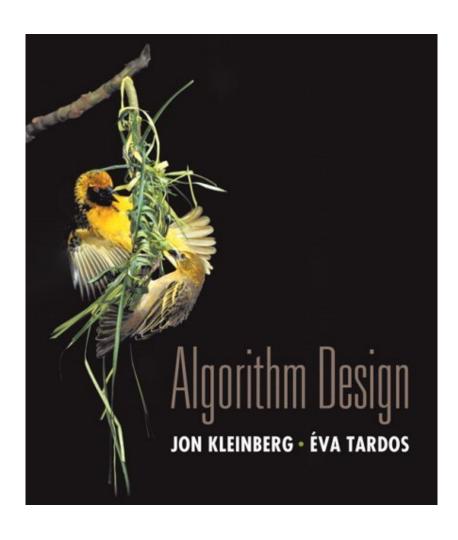
Analysis and argument as well as design

Midterm / Final Exam (~15% / 30%)

Late Policy:

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Textbook



Algorithm Design by Jon Kleinberg and Eva Tardos. Addison Wesley, 2006.

What the course is about

Design of Algorithms

design methods
common or important types of problems
analysis of algorithms - efficiency
correctness proofs

What the course is about

Complexity, NP-completeness and intractability solving problems in principle is not enough algorithms must be efficient

some problems have no efficient solution

NP-complete problems

important & useful class of problems whose solutions (seemingly) cannot be found efficiently, but *can* be checked easily

Very Rough Division of Time

Algorithms (7 weeks)

Analysis of Algorithms

Basic Algorithmic Design Techniques

Applications

Complexity & NP-completeness (3 weeks)

Check online schedule page for (evolving) details



University of Washington Computer Science & Engineering

CSE 417, Wi '06: Approximate Schedule

CSE Home

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		Due	Lecture Topic	Reading
Week 1 1/2-1/6	М		Holiday	
	W		Intro, Examples & Complexity	Ch. 1; Ch. 2
	F		Intro, Examples & Complexity	
Week 2 1/9-1/13	М		Intro, Examples & Complexity	7
	W		Graph Algorithms	Ch. 3
	F		Graph Algorithms	

Complexity Example

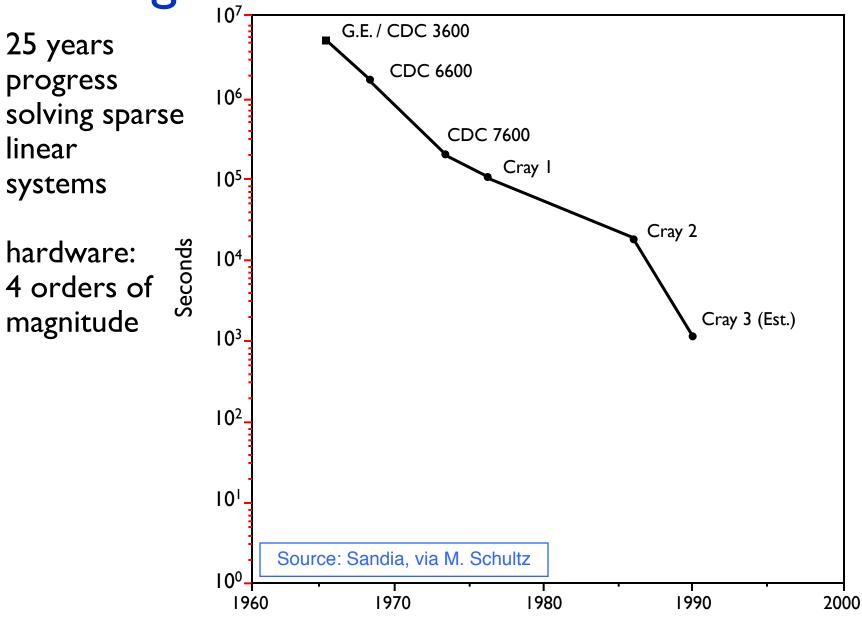
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Cryptography (e.g., RSA, SSL in browsers)
    Secret: p,q prime, say 512 bits each
    Public: n which equals p \times q, 1024 bits
In principle
    there is an algorithm that given n will find p and q:
   try all 2^{512} > 1.3 \times 10^{154} possible p's: kinda slow...
In practice
    no fast algorithm known for this problem (on non-quantum computers)
    security of RSA depends on this fact
    ("quantum computing": strongly driven by possibility of changing this)
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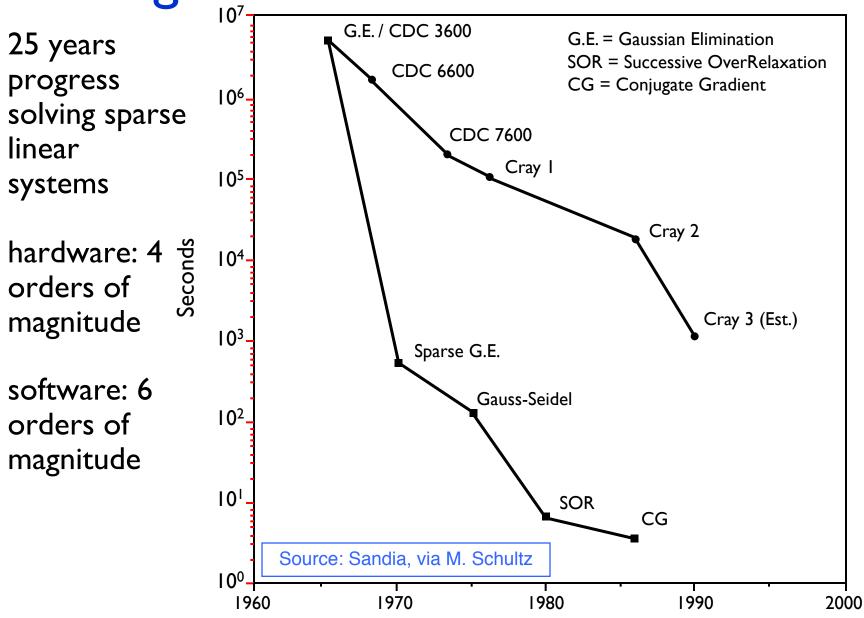
Algorithms versus Machines

We all know about Moore's Law and the exponential improvements in hardware...

Ex: sparse linear equations over 25 years

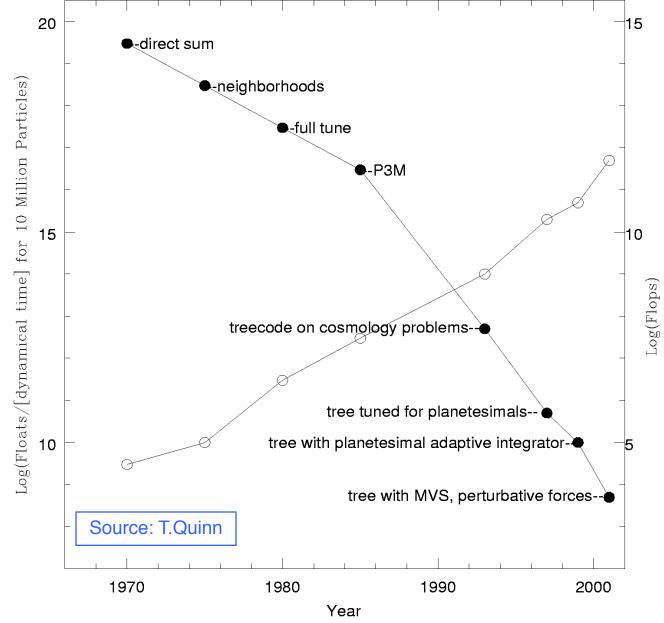
10 orders of magnitude improvement!



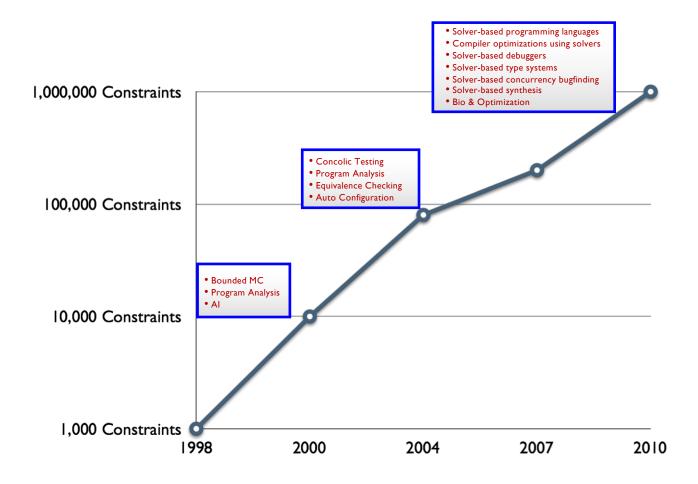


The N-Body Problem:

in 30 years 10⁷ hardware 10¹⁰ software



SAT/SMT Solvers: 1000x improvement in a dozen years



Data courtesy of Dr. Vijay Ganesh, U. Waterloo

Algorithm: definition

Procedure to accomplish a task or solve a well-specified problem

Well-specified: know what all possible inputs look like and what output looks like given them "accomplish" via simple, well-defined steps

Ex: sorting names (via comparison)

Ex: checking for primality (via +, -, *, /, \leq)

Goals

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Correctness

often subtle

Analysis

often subtle

Generality, Simplicity, 'Elegance'

Efficiency

time, memory, network bandwidth, ...
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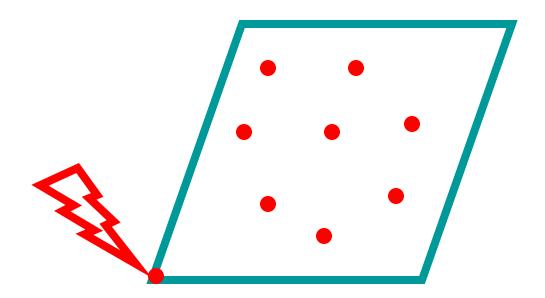
Algorithms: a sample problem

Printed circuit-board company has a robot arm that solders components to the board

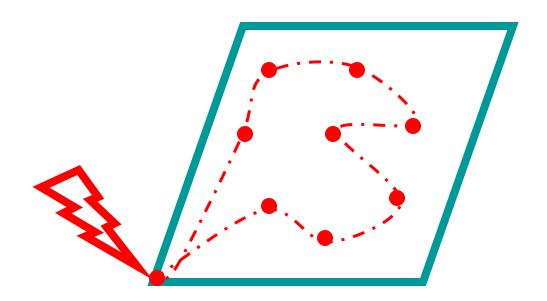
Time: proportional to total distance the arm must move from initial rest position around the board and back to the initial position

For each board design, find best order to do the soldering

Printed Circuit Board



Printed Circuit Board



A Well-defined Problem

Input: Given a set *S* of *n* points in the plane Output: The shortest cycle tour that visits each point in the set *S* once.

Better known as "TSP"

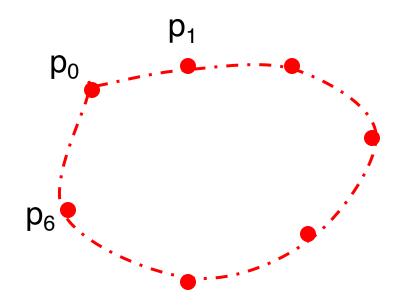
How might you solve it?

Nearest Neighbor Heuristic

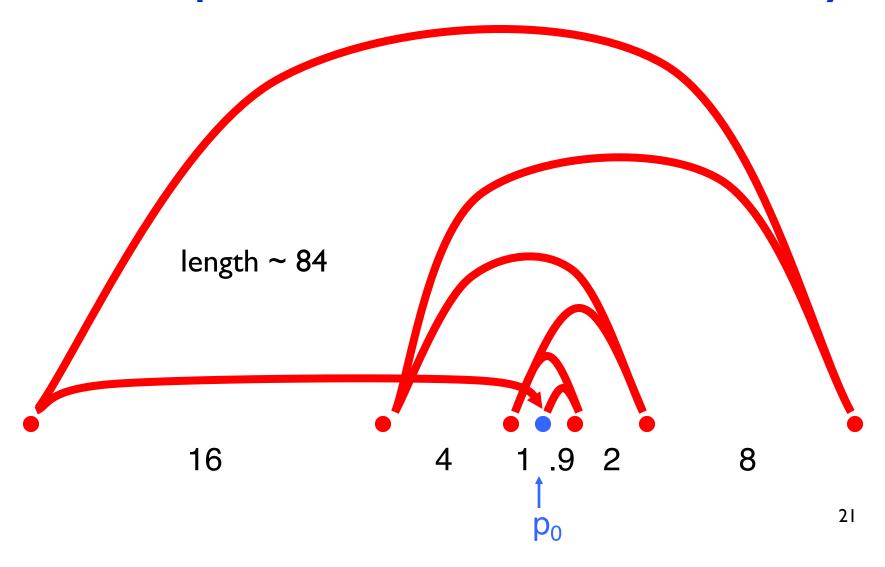
Start at some point p_0 Walk first to its nearest neighbor p_1 heuristic: A rule of thumb, simplification, or educated guess that reduces or limits the search for solutions in domains that are difficult and poorly understood. May be good, but usually *not* guaranteed to give the best or fastest solution.

Repeatedly walk to the nearest unvisited neighbor p_2 , then p_3 ,... until all points have been visited. Then walk back to p_0

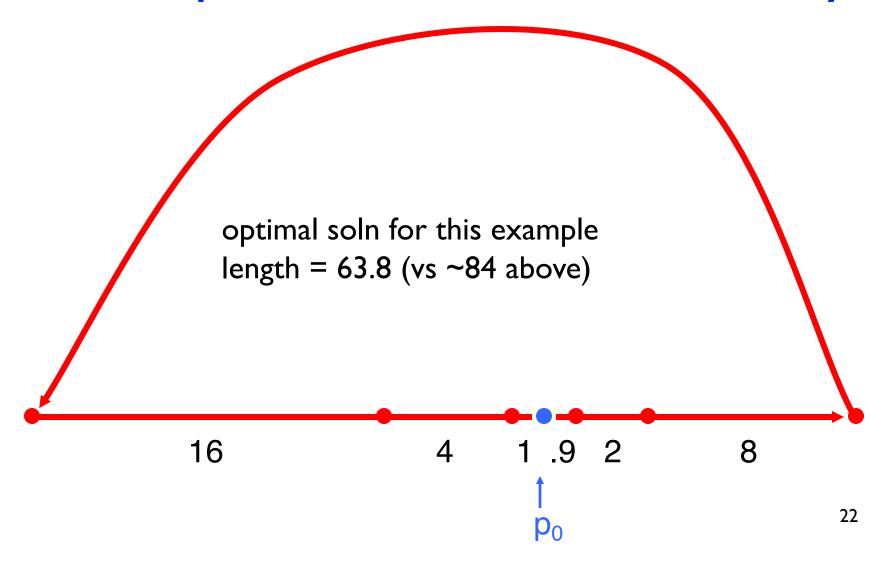
Nearest Neighbor Heuristic



An input where NN works badly



An input where NN works badly

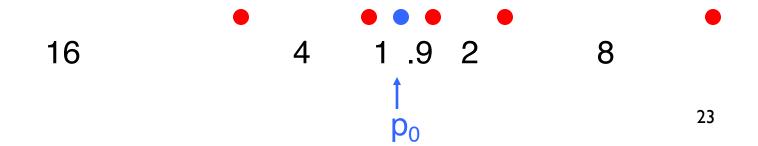


Revised idea - Closest pairs first

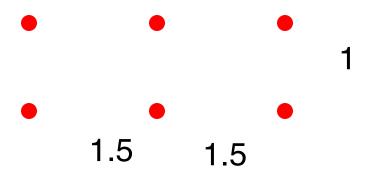
Repeatedly join the closest pair of points

(s.t. result can still be part of a single loop in the end. I.e., join endpoints, but not points in middle, of path segments already created.)

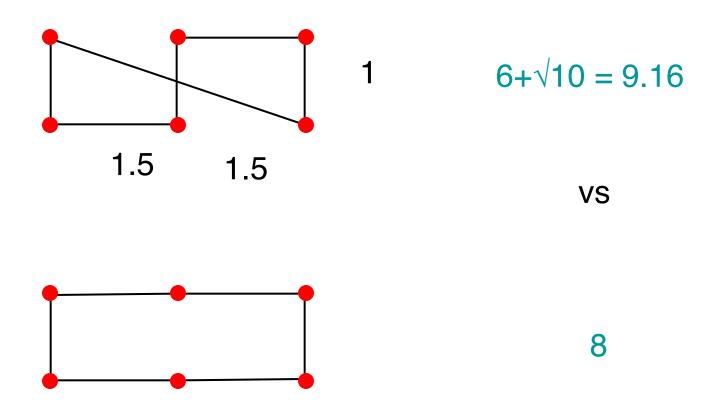




Another bad example



Another bad example



Something that works

"Brute Force Search":

For each of the n! = n(n-1)(n-2)...I orderings of the points, check the length of the cycle you get Keep the best one

Two Notes

The two incorrect algorithms were "greedy"

Often very natural & tempting ideas

They make choices that look great "locally" (and never reconsider them)

When greed works, the algorithms are typically efficient

BUT: often does not work - you get boxed in

Our correct alg avoids this, but is incredibly slow

20! is so large that checking one billion orderings per second would take 2.4 billion seconds (around 70 years!)

And growing: n! ~ $\sqrt{2} \pi n$ · $(n/e)^n$ ~ $2^{O(n \log n)}$

The Morals of the Story

Algorithms are important

Many performance gains outstrip Moore's law

Simple problems can be hard

Factoring, TSP

Simple ideas don't always work

Nearest neighbor, closest pair heuristics

Simple algorithms can be very slow

Brute-force factoring, TSP

For some problems, even the best algorithms are slow

Course Goals:

formalize these ideas, and develop more sophisticated approaches

