# CSE 4I7: Algorithms and <br> Computational Complexity 

# Lecture I: Overview 

Winter 2019
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## University of Washington <br> Computer Science \& Engineering

CSE 417, Wi '19: Algorithms \& Computational Complexity

Administrative

## FAQ

Schedule \& Reading
Course Email/BBoard
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Class List Archive
E-mail Course Staff
Google Groups BBoard

## Lecture Notes

1: Overview \& Example

## Lecture Recordings

Lecture: JHN 102 (room info) MWF 1:30-2:20

Instructor: Larry Ruzzo, ruzzo@cs
IAs: Yuqing Ai, yuqingai@cs
Daniel Jones, dcjones@cs
TBA
Office Hours Location Phone

F 2:30-3:30 CSE $554 \quad$ (206) 543-6298 <-1/11 excepted
Tu 1:00-2:00 4th floor breakout

Saidutt Nimmagadda, nimmas@cs TBA
Alex Okeson, amokeson@cs Tu 2:30-3:30 CSE 021
Course Email: cse417a_wi19@uw.edu. Staff announcements and lectures, etc. The instructor and WAs are subscribed $t n+2 \cdot$ default subscription options. Messages are

## t homework,

 should change theirdefault subscription options. Messages ara
Discussion Board: Also f-
Catalog Dom-
courses.CS.Was
using: Homework, Midterm, Final. Homework will be a mix of paper \& pencil exercises and programing. Overall weights $55 \%, 15 \%, 30 \%$, roughly.
Late Policy: Papers and/or electronic turning are due at the start of class on the due date. $10 \%$ off for up to one day late; additional $15 \%$ per day thereafter. (Day = calendar day, ie., Sunday is later than Saturday.)
Textbooks: Algorithm Design by Jon Kleinberg and Eva Tardos. Addison Wesley, 2006. (Available from U Book Store,
Open "mailman.u.washington.edu/mailman/private/cse417a_wi19" in a new tab

## What you'll have to do

## Homework Programming

Several small projects
Written homework assignments
English exposition and pseudo-code
Analysis and argument as well as design
Midterm / Final Exam (~15\% / 30\%)
Late Policy:
Papers and/or electronic turnins are due at the start of class on the due date. I0\% off for one day late; $15 \%$ per day thereafter.

## Textbook



Algorithm Design by Jon Kleinberg and Eva Tardos. Addison Wesley, 2006.

## What the course is about

## Design of Algorithms

 design methodscommon or important types of problems analysis of algorithms - efficiency correctness proofs

## What the course is about

Complexity, NP-completeness and intractability
solving problems in principle is not enough algorithms must be efficient
some problems have no efficient solution
NP-complete problems
important \& useful class of problems whose solutions (seemingly) cannot be found efficiently, but can be checked easily

## Very Rough Division of Time

Algorithms (7 weeks)
Analysis of Algorithms
Basic Algorithmic Design Techniques
Applications
Complexity \& NP-completeness (3 weeks)

Check online schedule page for (evolving) details


CSE 417, Wi '06: Approximate Schedule
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|  |  | Due | Lecture Topic | Reading |
| :---: | :--- | :--- | :--- | :--- |
| Week 1 <br> $1 / 2-1 / 6$ | M |  | Holiday |  |
|  | W |  | Intro, Examples \& Complexity | Ch. 1; Ch. 2 |
|  | F |  | Intro, Examples \& Complexity |  |
| Week 2 <br> $\mathbf{1 / 9 - 1 / 1 3}$ | M |  | Intro, Examples \& Complexity |  |
|  | W |  | Graph Algorithms | Ch. 3 |
|  | F |  | Graph Algorithms |  |

## Complexity Example

Cryptography (e.g., RSA, SSL in browsers)
Secret: p,q prime, say 512 bits each
Public: n which equals $\mathrm{p} \times \mathrm{q}, 1024$ bits
In principle
there is an algorithm that given n will find p and q :
try all $2^{512}>1.3 \times 10^{154}$ possible p's: kinda slow...
In practice
no fast algorithm known for this problem (on non-quantum computers)
security of RSA depends on this fact
("quantum computing": strongly driven by possibility of changing this)

## Algorithms versus Machines

We all know about Moore's Law and the exponential improvements in hardware...

Ex: sparse linear equations over 25 years

IO orders of magnitude improvement!

## Algorithms or Hardware?



## Algorithms or Hardware?

## 25 years progress

 solving sparse linear systemshardware: 4 orders of magnitude software: 6 orders of magnitude


## Algorithms or Hardware?

## The N-Body Problem:

in 30 years
$10^{7}$ hardware $10^{10}$ software


## Algorithms or Hardware?

SAT/SMT Solvers: 1000x improvement in a dozen years


## Algorithm: definition

Procedure to accomplish a task or solve a well-specified problem

Well-specified: know what all possible inputs look like and what output looks like given them
"accomplish" via simple, well-defined steps
Ex: sorting names (via comparison)
Ex: checking for primality (via,,+- ,,,$\leq$ )

## Goals

Correctness
often subtle
Analysis
often subtle
Generality, Simplicity, 'Elegance'
Efficiency
time, memory, network bandwidth, ...

## Algorithms: a sample problem

Printed circuit-board company has a robot arm that solders components to the board

Time: proportional to total distance the arm must move from initial rest position around the board and back to the initial position

For each board design, find best order to do the soldering

## Printed Circuit Board



## Printed Circuit Board



## A Well-defined Problem

Input: Given a set $S$ of $n$ points in the plane Output: The shortest cycle tour that visits each point in the set $S$ once.

Better known as "TSP"

How might you solve it?

## Nearest Neighbor Heuristic

Start at some point $\mathrm{P}_{0}$
Walk first to its
nearest neighbor $\mathrm{P}_{\mathrm{I}}$
heuristic: A rule of thumb, simplification, or educated guess that reduces or limits the search for solutions in domains that are difficult and poorly understood. May be good, but usually not guaranteed to give the best or fastest solution.

Repeatedly walk to the nearest unvisited neighbor $P_{2}$, then $P_{3}, \ldots$ until all points have been visited Then walk back to $\mathrm{P}_{0}$

## Nearest Neighbor Heuristic



An input where NN works badly


## An input where NN works badly



## Revised idea - Closest pairs first

Repeatedly join the closest pair of points (s.t. result can still be part of a single loop in the end. I.e., join endpoints, but not points in middle, of path segments already created.)


How does this work on our bad example?

$$
16
$$

4


8

## Another bad example

 $\begin{array}{llllll}\bullet & & \bullet & & \bullet & \\ \bullet & & & & & \\ & 1.5 & & & \bullet & \\ & & & & \end{array}$
## Another bad example



## Something that works

"Brute Force Search":
For each of the $n!=n(n-I)(n-2) \ldots I$ orderings of the points, check the length of the cycle you get
Keep the best one

## Two Notes

The two incorrect algorithms were "greedy" Often very natural \& tempting ideas
They make choices that look great "locally" (and never reconsider them)
When greed works, the algorithms are typically efficient BUT: often does not work - you get boxed in
Our correct alg avoids this, but is incredibly slow 20 ! is so large that checking one billion orderings per second would take 2.4 billion seconds (around 70 years!)
And growing: $n!\sim \sqrt{2 \pi n} \cdot(n / e)^{n} \sim 2^{O}(\mathrm{n} \log \mathrm{n})$

## The Morals of the Story

Algorithms are important
Many performance gains outstrip Moore's law
Simple problems can be hard Factoring, TSP
Simple ideas don't always work
Nearest neighbor, closest pair heuristics
Simple algorithms can be very slow
Brute-force factoring, TSP
For some problems, even the best algorithms are slow
Course Goals:
formalize these ideas, and
develop more sophisticated approaches

https://xkcd.com/

