Game Playing

Why do AI researchers study game playing?

- 1. It's a good reasoning problem, formal and nontrivial.
- 2. Direct comparison with humans and other computer programs is easy.

What Kinds of Games?

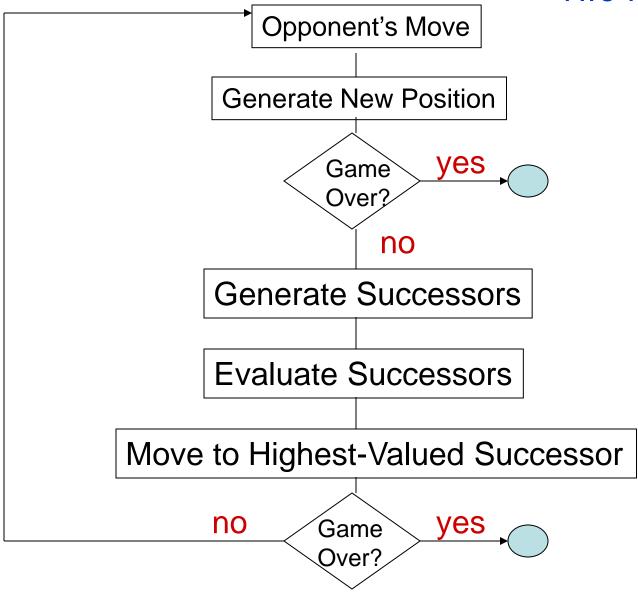
Mainly games of strategy with the following characteristics:

- 1. Sequence of moves to play
- 2. Rules that specify possible moves
- 3. Rules that specify a payment for each move
- 4. Objective is to maximize your payment

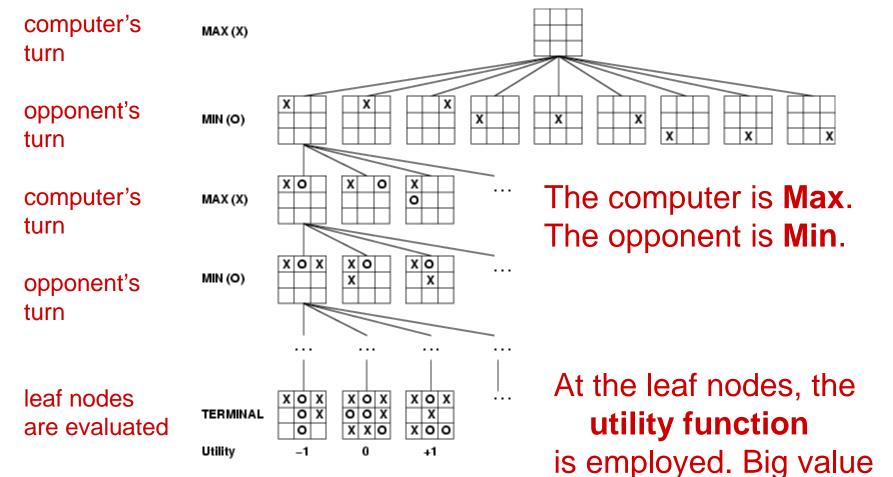
Games vs. Search Problems

- Unpredictable opponent → specifying a move for every possible opponent reply
- Time limits → unlikely to find goal, must approximate

Two-Player Game



Game Tree (2-player, Deterministic, Turns)



means good, small is bad.

Mini-Max Terminology

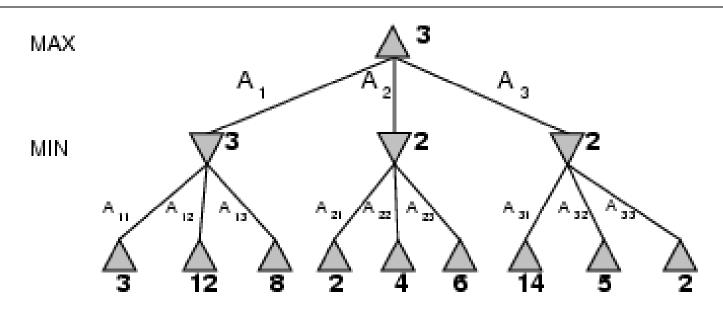
- utility function: the function applied to leaf nodes
- backed-up value
 - of a max-position: the value of its largest successor
 - of a min-position: the value of its smallest successor
- minimax procedure: search down several levels; at the bottom level apply the utility function, back-up values all the way up to the root node, and that node selects the move.

Minimax

- Perfect play for deterministic games
- Idea: choose move to position with highest minimax value

= best achievable payoff against best play

• E.g., 2-level game:



Minimax Strategy

- Why do we take the min value every other level of the tree?
- These nodes represent the opponent's choice of move.
- The computer assumes that the human will choose that move that is of least value to the computer.

Minimax algorithm

How?

function MINIMAX-DECISION(state) returns an action

```
v \leftarrow \text{MAX-VALUE}(state)
```

return the action in SUCCESSORS(state) with value v

function MAX-VALUE(state) returns a utility value

if TERMINAL-TEST(state) then return UTILITY(state)

```
v \leftarrow -\infty
```

for a, s in SUCCESSORS(state) do

```
v \leftarrow Max(v, Min-Value(s))
```

return v

function MIN-VALUE(state) returns a utility value

```
if TERMINAL-TEST(state) then return UTILITY(state)
```

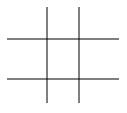
```
v \leftarrow \infty
```

for a, s in SUCCESSORS(state) do

 $v \leftarrow MIN(v, MAX-VALUE(s))$

return v

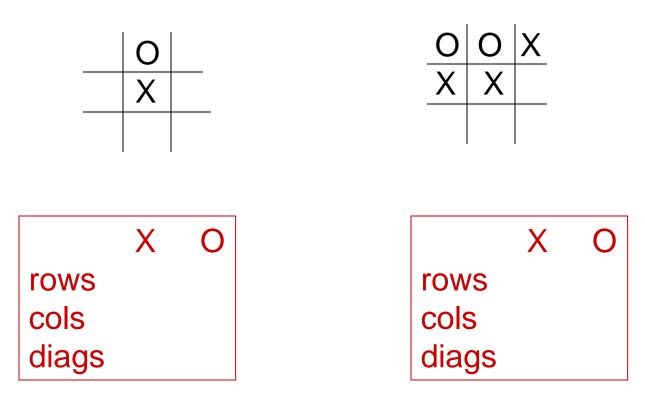
Tic Tac Toe



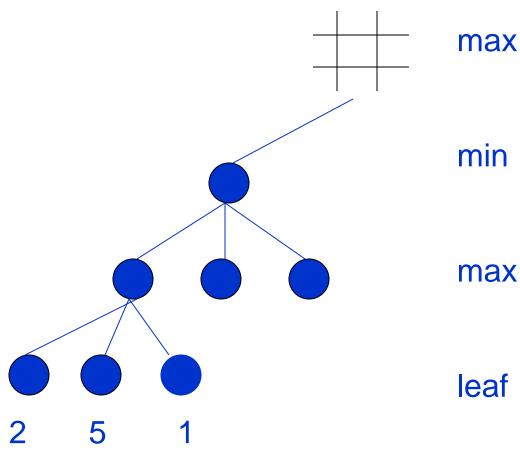
- Let p be a position in the game
- Define the utility function f(p) by
 - f(p) =
 - largest positive number if p is a win for computer
 - smallest negative number if p is a win for opponent
 - RCDC RCDO
 - where RCDC is number of rows, columns and diagonals in which computer could still win
 - and RCDO is number of rows, columns and diagonals in which opponent could still win.

Sample Evaluations

• X = Computer; O = Opponent



Minimax is done depth-first



max

Properties of Minimax

- <u>Complete?</u> Yes (if tree is finite)
- Optimal? Yes (against an optimal opponent)
- <u>Time complexity?</u> O(b^m)
- <u>Space complexity?</u> O(bm) (depth-first exploration)
- For chess, b ≈ 35, m ≈100 for "reasonable" games
 → exact solution completely infeasible

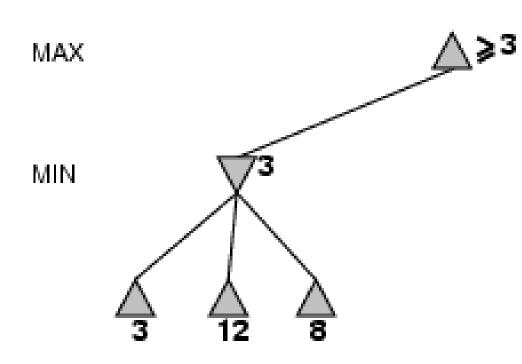
Need to speed it up.

Alpha-Beta Procedure

- The alpha-beta procedure can speed up a depth-first minimax search.
- Alpha: a lower bound on the value that a max node may ultimately be assigned

$v \ge \alpha$

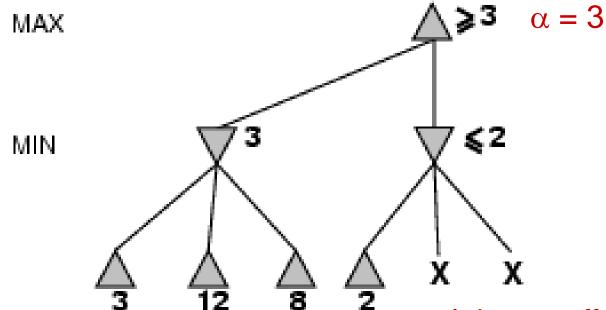
Beta: an upper bound on the value that a minimizing node may ultimately be assigned
 ν ≤ β



 $\alpha = 3$

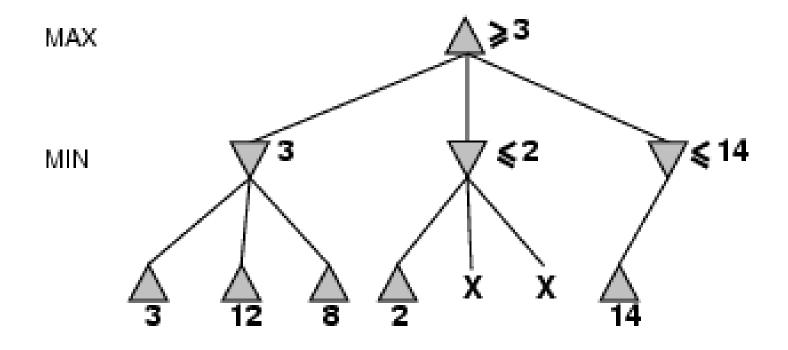
The root node will end up with a value of 3 or more, no matter what the rest of its children do.

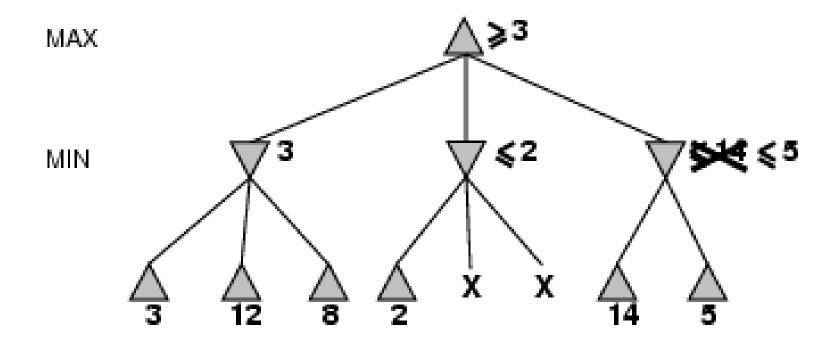
first child of root has been expanded two levels

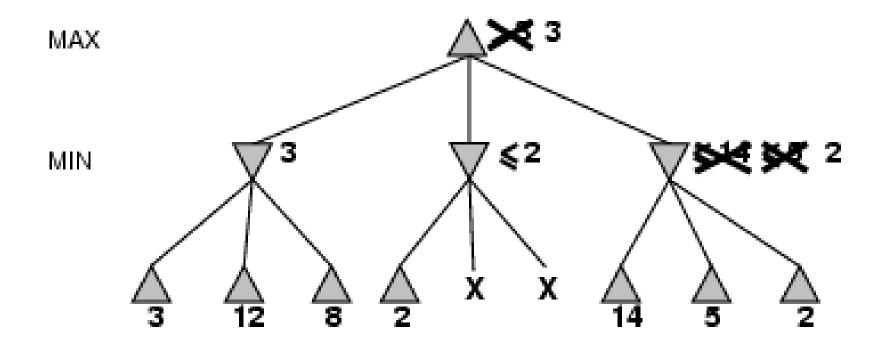


alpha cutoff

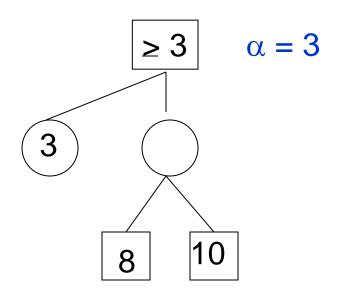
As soon as we know that the min node will return a value less than alpha, no need to look at more of its children 16



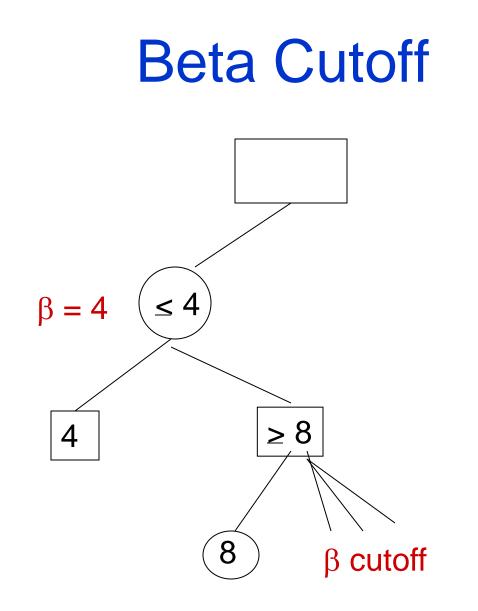




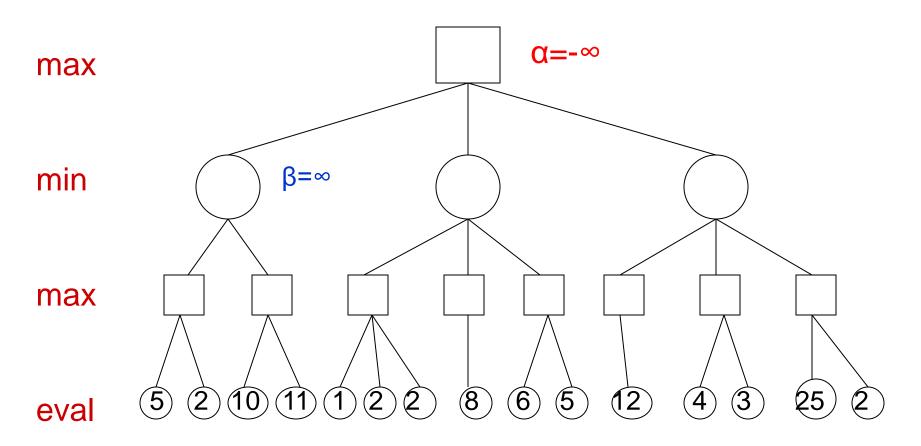
Alpha Cutoff



What happens here? Is there an alpha cutoff?



Alpha-Beta Pruning



Properties of α - β

- Pruning does not affect final result. This means that it gets the exact same result as does full minimax.
- Good move ordering improves effectiveness of pruning
- With "perfect ordering," time complexity = O(b^{m/2})
 → doubles depth of search
- A simple example of the value of reasoning about which computations are relevant (a form of metareasoning)

The α - β algorithm

function ALPHA-BETA-SEARCH(state) returns an action inputs: state, current state in game

 $v \leftarrow \text{MAX-VALUE}(state, -\infty, +\infty)$

return the action in SUCCESSORS(state) with value v

function MAX-VALUE(*state*, α , β) returns a utility value inputs: *state*, current state in game

lpha, the value of the best alternative for ${
m MAX}$ along the path to state

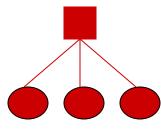
 $\beta,$ the value of the best alternative for $\ _{\rm MIN}$ along the path to state

```
if TERMINAL-TEST(state) then return UTILITY(state)
```

```
v \leftarrow -\infty
```

for a, s in SUCCESSORS(state) do $v \leftarrow MAX(v, MIN-VALUE(s, \alpha, \beta))$ if $v \ge \beta$ then return v CUtOff $\alpha \leftarrow MAX(\alpha, v)$

return v



The α - β algorithm

```
function MIN-VALUE(state, \alpha, \beta) returns a utility value

inputs: state, current state in game

\alpha, the value of the best alternative for MAX along the path to state

\beta, the value of the best alternative for MIN along the path to state

if TERMINAL-TEST(state) then return UTILITY(state)

v \leftarrow +\infty

for a, s in SUCCESSORS(state) do

v \leftarrow MIN(v, MAX-VALUE(s, \alpha, \beta))

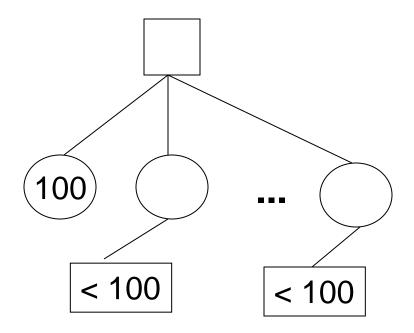
if v \le \alpha then return v CUtOff

\beta \leftarrow MIN(\beta, v)

return v
```

Should α and β be passed by value or reference? ie. Should a lower α affect an upper one?

When do we get alpha cutoffs?



Shallow Search Techniques

1. limited search for a few levels

2. reorder the level-1 sucessors

3. proceed with α - β minimax search

Additional Refinements

- Waiting for Quiescence: continue the search until no drastic change occurs from one level to the next.
- Secondary Search: after choosing a move, search a few more levels beneath it to be sure it still looks good.
- Book Moves: for some parts of the game (especially initial and end moves), keep a catalog of best moves to make.

Evaluation functions

 For chess/checkers, typically linear weighted sum of features

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + ... + w_n f_n(s)$$

e.g., $w_1 = 9$ with

 $f_1(s) = (number of white queens) - (number of black queens), etc.$

Example: Samuel's Checker-Playing Program

• It uses a linear evaluation function $f(n) = a_1 x_1(n) + a_2 x_2(n) + \dots + a_m x_m(n)$

- For example: f = 6K + 4M + U
- K = King Advantage
- M = Man Advantage
- U = Undenied Mobility Advantage (number of moves that Max has that Min can't jump after)

Samuel's Checker Player

• In learning mode

- Computer acts as 2 players: A and B
- A adjusts its coefficients after every move
- B uses the static utility function
- If A wins, its function is given to B

Samuel's Checker Player

- How does A change its function?
 - 1. Coefficent replacement

(node) = backed-up value(node) - initial value(node)

if $\triangle > 0$ then terms that contributed **positively** are given more weight and terms that contributed negatively get less weight

if \triangle < 0 then terms that contributed **negatively** are given more weight and terms that contributed positively get less weight

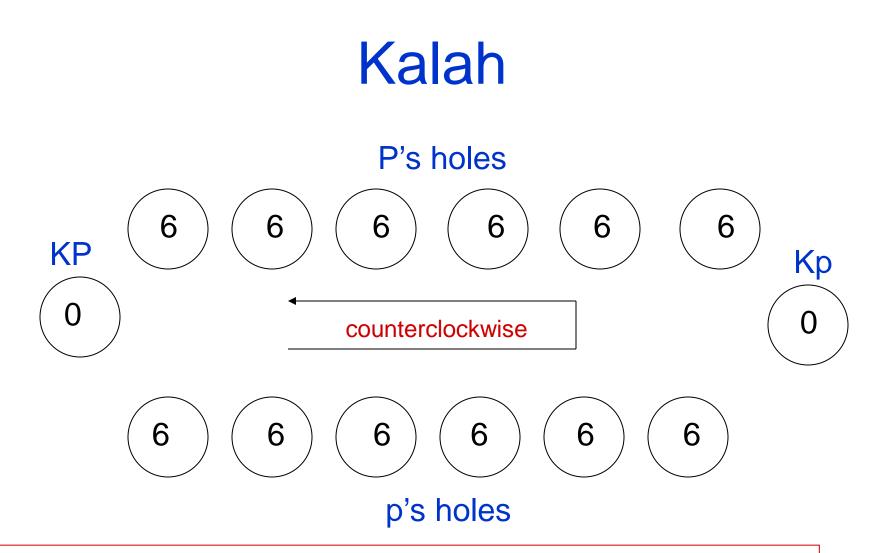
Samuel's Checker Player

How does A change its function?
 2. Term Replacement

 38 terms altogether
 16 used in the utility function at any one time

Terms that **consistently correlate low** with the function value are removed and added to the end of the term queue.

They are replaced by terms from the front of the term queue.



To move, pick up all the stones in one of your holes, and put one stone in each hole, starting at the next one, including your Kalah and skipping the opponent's Kalah.

34

Kalah

- If the last stone lands in your Kalah, you get another turn.
- If the last stone lands in your empty hole, take all the stones from your opponent's hole directly across from it and put them in your Kalah.
- If all of your holes become empty, the opponent keeps the rest of the stones.
- The winner is the player who has the most stones in his Kalah at the end of the game.

Cutting off Search

MinimaxCutoff is identical to *MinimaxValue* except

- 1. Terminal? is replaced by Cutoff?
- 2. Utility is replaced by Eval

Does it work in practice? $b^m = 10^6$, $b=35 \rightarrow m=4$

4-ply lookahead is a hopeless chess player!

- 4-ply ≈ human novice
- 8-ply ≈ typical PC, human master
- 12-ply ≈ Deep Blue, Kasparov

Deterministic Games in Practice

- Checkers: Chinook ended 40-year-reign of human world champion Marion Tinsley in 1994. Used a precomputed endgame database defining perfect play for all positions involving 8 or fewer pieces on the board, a total of 444 billion positions.
- Chess: Deep Blue defeated human world champion Garry Kasparov in a six-game match in 1997. Deep Blue searches 200 million positions per second, uses very sophisticated evaluation, and undisclosed methods for extending some lines of search up to 40 ply.
- Othello: human champions refuse to compete against computers, who are too good.
- Go: human champions refuse to compete against computers, who are too bad. In Go, b > 300, so most programs use pattern knowledge bases to suggest plausible moves.

Games of Chance

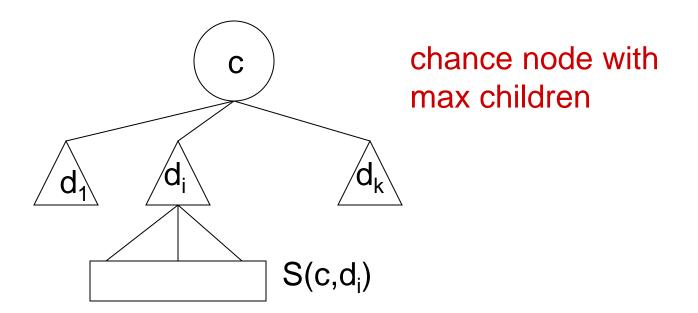
- What about games that involve chance, such as
 - rolling dice
 - picking a card
- Use three kinds of nodes:
 - max nodes
 - min nodes
 - chance nodes

min

chance

max

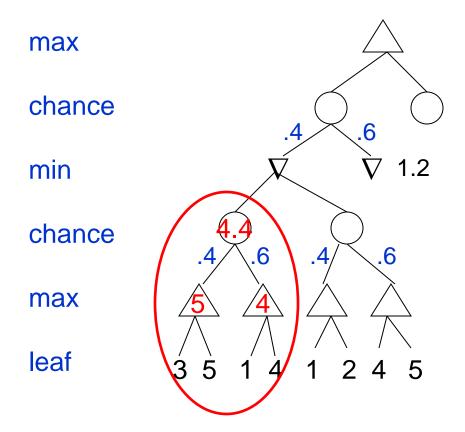
Games of Chance



expectimax(c) = $\sum P(d_i) \max(backed-up-value(s))$ i s in S(c,d_i)

expectimin(c') = $\sum P(d_i) \min(backed-up-value(s))$ i s in S(c,d_i)

Example Tree with Chance



5(.4)+4(.6)=4.4

Complexity

 Instead of O(b^m), it is O(b^mn^m) where n is the number of chance outcomes.

• Since the complexity is higher (both time and space), we cannot search as deeply.

Pruning algorithms may be applied.

Summary

- Games are fun to work on!
- They illustrate several important points about AI.
- Perfection is unattainable \rightarrow must approximate.
- Game playing programs have shown the world what AI can do.