Database Systems
CSE 414

Section 10: Big Data & Review
Non-Parallel Query Evaluation
Example Schema

Product(\textit{pid}, name, category)
- 10,000 tuples and 1,000 blocks
- 40 different categories

Order(store, pid, price, quantity)
- 1,000,000 tuples and 50,000 blocks
- prices range from $1 to $100
Example Query

Compute the total revenue, for each store, from electronics costing more than $5 each:

SELECT o.store, sum(o.price * o.quantity)
FROM Order o, Product p
WHERE o.pid = p.pid AND o.price > 5 AND
  p.category = 'electronics'
GROUP BY o.store
Problem 1

Give an RA expression that:

- computes the result of the query
- does not benefit from the indexes already present

```
Order    σ_{price > 5}
         ⋈_{pid = pid} π_{store, rev}
```

```
Product   σ_{category = 'electronics'}
```

```
γ_{store, sum(price*quanity) \rightarrow rev}
```
Problem 2

Estimate the cost of the RA expression from Problem 1 after filling in physical implementation details
  – assume grouping / aggregation can be done on the fly

• Details:
  – nested loop join
  – write Products to temp T1
  – grouping / aggregation done with in memory hash table

• Scan Product & writing to T1 costs 50k + 1k + 25
• Nested loop join costs 47.5k * 25 = 1,125k
• Total cost is 1,238,525 blocks (~1M is fine)
Problem 3

Give an RA expression that:
- computes the result of the query
- does benefit from the indexes already present

\[
\text{Order} \quad \sigma_{\text{price} > 5} \quad \bowtie_{\text{pid} = \text{pid}} \quad \sigma_{\text{category} = \text{electronics}}
\]

\[
\forall \text{store}, \sum(\text{price} \times \text{quantity}) \rightarrow \text{rev} \quad \pi_{\text{store}, \text{rev}}
\]
Problem 4

Estimate the cost of the RA expression from Problem 3 after filling in physical implementation details

- assume grouping / aggregation can be done on the fly

• Details:
  - nested loop join using index on Product(pid)
  - grouping / aggregation done with in memory hash table

• Lookup of Product costs 1 block
• Nested loop join costs $50k + 950k \times 1 = 1000k$
• Total cost is $\sim1M$ blocks (everything else on the fly)
Parallel Query Evaluation
Problem 5

Draw a pipeline that computes the same result in a parallel fashion using N nodes

\[ O_{1,P_{1}} \]

select on O.price & P.category

\[ O'_{1,P'_{1}} \]

join

\[ R_{1} \]

aggregate

\[ O_{N,P_{N}} \]

select on O.price & P.category

\[ O'_{N,P'_{N}} \]

join

\[ R_{N} \]

aggregate

shuffle on O.pid & P.pid

shuffle on store
Problem 6

Estimate the cost of executing the pipeline of Problem 5

• Only costs are on disk reads of input
  – (everything should fit in memory)
• Each worker reads $50k/N + 1k/N$ blocks
• Since all workers are reading simultaneous, wait time is time to read $51k/N$ blocks (plus lower order work)
Problem 7

1. Does your analysis predict a linear speedup as more nodes are added?  
   Yes

2. Does your analysis predict a linear scaleup as more nodes are added?  
   Yes

3. How realistic is this?  
   Fair with a small number of machines, but expect stragglers to be noticeable with 1000s
Describe how to achieve a similar speedup with MapReduce

- MapReduce does only one shuffle, so we need 2 jobs
- First job:
  - map Orders to (pid, (‘O’, …)) and Products to (pid, (‘P’, …)) for those rows that satisfy selection criteria
  - reducer adds product info to each order in the list
    - note: only one Product in each list since pid is primary key
- Second job:
  - map Order+Product to (store, (...))
  - reducer sums revenue and outputs (store, revenue)
Problem 9

Would your MapReduce have the same IO cost and speedup as the pipeline from problem 6?

• MapReduce writes intermediate results to disk resulting in more IO
  – Two intermediate results and two outputs written
  – None of these are larger than the input, though, so the total cost is no more than 7x the ideal pipeline
    • really 6x since the final output is small

• Despite a constant factor more IO, it should still have a linear speedup (in principle).