
Procedures

CSE 413, Autumn 2005
Programming Languages

<http://www.cs.washington.edu/education/courses/413/05au/>

References

- Sections 1.1.6-1.1.8, *Structure and Interpretation of Computer Programs*
- Section 4.1, *Revised⁵ Report on the Algorithmic Language Scheme (R5RS)*

Combinations

- (operator operand operand)
- There are numerous pre-defined operators
- We can define our own, arbitrarily complex operators (functions, procedures) as well
- This is a key capability by which we can operate at higher levels of abstraction

Recall the *define* special form

- Special forms have unique evaluation rules
- **(define x 3)** is an example of a special form; it is not a combination
 - » the evaluation rule for a simple define is "associate the given name with the given value"

Define and name a variable

- **(define** *<name>* *<expr>*)
 - » **define** - special form
 - » *name* - name that the value of *expr* is bound to
 - » *expr* - expression that is evaluated to give the value for *name*
- **define** is valid only at the top level of a *<program>* and at the beginning of a *<body>*

Define and name a procedure

- **(define** (*<name>* *<formal params>*) *<body>*)
 - » **define** - special form
 - » *name* - the name that the procedure is bound to
 - » *formal params* - names used within the body of procedure
 - » *body* - expression (or sequence of expressions) that will be evaluated when the procedure is called.
 - » The result of the last expression in the body will be returned as the result of the procedure call

Example definitions

```
(define pi 3.1415926535)

(define (area-of-disk r)
  (* pi (* r r)))

(define (area-of-ring outer inner)
  (- (area-of-disk outer)
     (area-of-disk inner)))
```

Defined procedures are "first class"

- Compound procedures that we define are used exactly the same way the primitive procedures provided in Scheme are used
 - » names of built-in procedures are not treated specially; they are simply names that have been pre-defined
 - » you can't tell whether a name stands for a primitive (built-in) procedure or a compound (defined) procedure by looking at the name or how it is used

Evaluation example

- **(area-of-ring 4 1)**
 - » evaluate operator **area-of-ring** => procedure definition
 - » evaluate 4 => 4
 - » evaluate 1 => 1
 - » apply the procedure to the arguments

Booleans

- Recall that one type of data object is boolean
 - » **#t** (true) or **#f** (false)
- We can use these explicitly or by calculating them in expressions that yield boolean values
- An expression that yields a true or false value is called a predicate
 - » **#t => #t**
 - » **(< 5 5) => #f**
 - » **(> pi 0) => #t**

Conditional expressions

- As in all languages, we need to be able to make decisions based on inputs and do something depending on the result
- A predicate expression is evaluated
 - » true or false
- The consequent expression is evaluated if the predicate is true

Special form: **cond**

- **(cond <clause₁> <clause₂> ... <clause_n>)**
- each clause is of the form
 - » **(<predicate> <expression>)**
 - » where *<predicate>* is a boolean expression and *<expression>* is the consequent expression to execute if *<predicate>* is true
- the last clause can be of the form
 - » **(else <expression>)**
 - » in which case *<expression>* is executed if none of the preceding *<predicates>* were true

Example: sign.scm

```
; return the sign of x as -1, 0, or 1
```

```
(define (sign x)
  (cond
    ((< x 0) -1)
    ((= x 0) 0)
    (> x 0) +1)))
```

also interest-rate.scm

Special form: **if**

- **(if** *<predicate>* *<consequent>* *<alternate>*)
- **(if** *<predicate>* *<consequent>*)

- *<predicate>* is a boolean expression
- *<consequent>* is the expression to execute if *<predicate>* is true
- *<alternate>* is the expression to execute if *<predicate>* is false

Examples : abs.scm, true-false.scm

```
; absolute value function
```

```
(define (abs a)
  (if (< a 0)
      (- a)
      a))
```

```
; return 1 if arg is true, 0 if arg is false
```

```
(define (true-false arg)
  (if arg 1 0))
```

Logical composition

- **(and** *<e₁>* *<e₂>*... *<e_n>*)
- **(or** *<e₁>* *<e₂>*... *<e_n>*)
- **(not** *<e>*)

- Scheme interprets the expressions *e_i* one at a time in left-to-right order until it can tell the correct answer
» ie, these are short-circuit operators

in-range.scm

```
; true if val is lo <= val <= hi

(define (in-range lo val hi)
  (and (<= lo val)
       (<= val hi)))
```

Newton's method for square root

- Guess a value y for the square root of x
- Is it close enough to the desired value \sqrt{x} ?
» ie, is y^2 close to x ?
- If yes, then done. Return recent guess.
- If no, then new guess is average of current *guess* and $\frac{x}{guess}$
- Repeat with new guess

sqrta.scm

```
; Square root using Newton's method

(define (average a b)
  (/ (+ a b) 2.0))

(define (good-enough? guess x)
  (< (abs (- (* guess guess) x)) 0.001))

(define (improve guess x)
  (average guess (/ x guess)))

(define (sqrt-iter guess x)
  (if (good-enough? guess x)
      guess
      (sqrt-iter (improve guess x) x)))

(define (sqrta x)
  (sqrt-iter 1.0 x))
```

auxiliary functions

```
; Square root using Newton's method

(define (average a b)
  (/ (+ a b) 2.0))

(define (good-enough? guess x)
  (< (abs (- (* guess guess) x)) 0.001))

(define (improve guess x)
  (average guess (/ x guess)))
```

iterator and main functions

```
(define (sqrt-iter guess x)
  (if (good-enough? guess x)
      guess
      (sqrt-iter (improve guess x) x )))

(define (sqrta x)
  (sqrt-iter 1.0 x))
```

sqrt-iter

- Our first example of recursion
- Note that this recursion is used to implement a loop (an iteration)
 - » We will see this over and over in Scheme
- Iteration is calling the same block of code with a changing set of parameters
- The syntax of the procedure is recursive but the resulting process is iterative
 - » more on this next lecture