Threads

CSE 410, Spring 2004
Computer Systems

http://www.cs.washington.edu/education/courses/410/04sp/
Reading and References

• Reading
  » Chapter 5, *Operating System Concepts*, Silberschatz, Galvin, and Gagne

• Other References
  » *Pthreads Programming*, Nichols, Buttlar and Farrell
A Process

- A complete process includes numerous things
  - address space (all the code and data pages)
  - OS resources and accounting information
  - a “thread of control”, which defines where the process is currently executing
    - the Program Counter
    - CPU registers
Processes are heavyweight objects

• Creating a new process is costly
  » lots of data must be allocated and initialized
  » operating system control data structures
  » memory allocation for the process

• Communicating between processes is costly
  » most communication goes through the OS
  » need a context switch for each process
Parallelism

• Why build a parallel program?
  » responsiveness to user
  » web server handling simultaneous web requests
  » execute faster on a multiprocessor

• One approach using heavyweight processes
  » create several processes to execute in parallel
  » map each process to same address space
  » specify starting address and initial parameters
Parallelism

• With multiple paths of execution, we can implement (or simulate) simultaneous actions

• Why build a parallel program?
  » responsiveness to user
    • user interface always responds quickly
  » web server handling simultaneous web requests
    • each request is handled independently
  » execute faster on a multiprocessor
    • two CPUs can run two programs at once
Parallel processes are expensive

- There’s a lot of performance cost
  - creating these processes
  - coordinating them through the OS
- There’s a lot of duplication
  - same program code, protection, etc…
- It may be time for a little refinement and complexity …
Process definition

• What is fundamental in a process?
  » Code and data
  » Access and control privileges
  » Operating system management
    • scheduling, memory map, ...

• What else is there?
  » Program Counter, registers, and stack

• Separate the idea of “process” from the idea of a “thread of control” (PC, SP, registers)
Threads are “Lightweight Processes”

- Most operating systems now support two entities
  - the process, which defines the address space and general process attributes
  - the thread, which defines one or more execution paths within a process
- Threads are the unit of scheduling
- Processes are the “containers” in which threads execute
Multi-threaded design benefits

• Separating execution path from address space simplifies design of parallel applications

• Some benefits of threaded designs
  » improved responsiveness to user actions
  » handling concurrent events (e.g., web requests)
  » simplified program structure (code, data)
  » more efficient and so less impact on system
  » map easily to multi-processor systems
One thread

Three threads
Cookbook Analogy

- Think of a busy kitchen
  - 3 cooks and 1 cookbook
- Each cook maintains a pointer to where they are in the cookbook (the Program Counter)
- Two cooks could both be making the same thing (threads running the same procedure)
- The cooks must coordinate access to the kitchen appliances (resource access control)
Implementation

- A thread is bound to the process that provides its address space
- Each process has one or more threads
- How are threads actually implemented?
  - Kernel threads
    - In the kernel (OS) and user mode libraries combined
  - User threads
    - In user mode libraries alone
Kernel Threads

- The operating system knows about and manages the threads in every program.
- Thread operations (create, yield, ...) all require kernel involvement.
- Major benefit is that threads in a process are scheduled independently:
  - one blocked thread does not block the others
  - threads in a process can run on different CPUs
Kernel Thread Performance

• Kernel threads have performance issues
• Even though threads avoid process overhead, operations on kernel threads are still slow
  » a thread operation requires a kernel call
  » kernel threads may be overly general, in order to support needs of different users, languages, etc.
  » the kernel can’t trust the user, so there must be lots of checking on kernel calls
User Threads

- To make thread operations faster, they can be implemented at the user level
  » Each thread is managed by the run-time system
  » user-mode libraries are linked with your program
- Each thread is represented simply by a PC, registers, stack and a control block, managed in the user’s address space
User Thread Performance

• All activities happen in user address space so thread operations can be faster

• But OS scheduling takes place at process level
  » block entire process if a single thread is I/O blocked
  » may run a process that is just running an idle thread

• Win2K provides “fibers” as user mode threads
  » application can schedule its own “lightweight threads” in user mode code
Simplified Thread Interface

- \( t = \text{thread}_\text{create}(), \text{thread}_\text{start}(t) \)
  - create a new thread of control and start it
- \( \text{thread}_\text{yield}() \)
  - voluntarily give up the processor for awhile
- \( \text{thread}_\text{exit}() \)
  - terminate the calling thread