CSE401: Semantic Analysis

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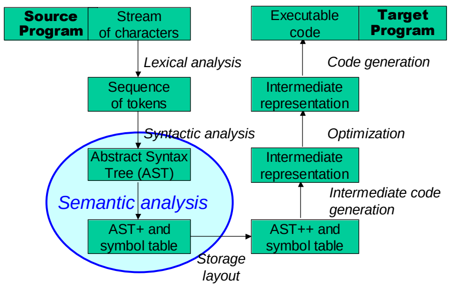
Prototype compiler structure



Semantic analysis

- Perform final legality checking of input program
  - Properties not checked by lexical or syntactic checking
    - Ex: type checking, ensuring break statement is in a loop, etc.
- "Understand" program well enough to do the back-end synthesis activities
  - Ex: relate particular names to particular declarations

Symbol tables

- Key data structure (at *compile* time, not run time)
  - Produced (and used) during semantic analysis
  - Used during code generation
- Stores information about names used in the program
  - Declarations add entries to the symbol table
  - Uses of names look up appropriate symbol table entry

What information about names?

- Kind of declaration
  - `var`, `const`, `proc`, etc.
- Type
- For `const:` keep value
- For `var:` Where allocated in memory?
  - Static, stack, heap? Offset?
  - Not computed initially, but later on
- For formal parameter: passed by-value, by-ref…

Example: a PL/0 `DeclList`

```
var x : int;
var q : array[20] of bool;
procedure foo(a : int); begin … end foo;
const z : int = 10;
```

PL/0 symbol table entries

```
class SymTabEntry {
public:
  char* name();
  Type* type();

  virtual bool isConstant();
  virtual bool isVariable();
  virtual bool isFormal();
  virtual bool isProcedure();

  virtual int value();              // const only
  virtual int offset(SymTabScope* s); // var only
}
```

More soon

SymTab subclasses

```
class VarSTE    : public SymTabEntry { …
};
class FormalSTE : public VarSTE { … };
class ConstSTE  : public SymTabEntry { …
};
class ProcSTE   : public SymTabEntry { …
};
```

Nested scopes: Example

```
procedure foo(x:int, w:int);
  var z:bool;
  const y:bool = true;
  procedure bar(x:array[5] of bool);
    var y:int;
  begin
    x[y] := z;
  end bar;
begin
  while z do
    var z:int, y:int;
    y := z * x;
  end;
  output := x + y;
end foo;
```

Nested scopes: How to handle?

- What happens when the same name is declared in different scopes?
- This is first a question of language design: what is the defined semantics?
- Two standard choices
  - Lexical (static) scoping: use the block structure of the program
  - Do you remember choice #2 from 341?

Nested Scopes: Lexical/static

- The syntactic (block) structure of the program determines how names are resolved
- Given a name in a block
  - The nearest enclosing block with a declaration for that name is the relevant declaration
  - If none, it's an error

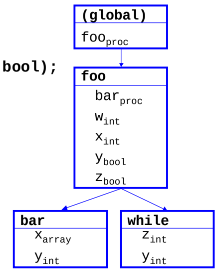
Nested scopes: Dynamic

Lexical scope and symbol tables

- Each scope has its own symbol table
- Logically, for a block-structured program, there is a *tree* of symbol tables
  - Root = outermost block

Tree of symbol tables

```
procedure foo(x:int, w:int);
  var z:bool;
  const y:bool = true;
  procedure bar(x:array[5] of bool);
    var y:int;
  begin
    x[y] := z;
  end bar;
begin
  while z do
    var z:int, y:int;
    y := z * x; end;
  output := x + y;
end foo;
```



Lexical scope and symbol tables

- Each scope has its own symbol table
- Logically, for a block-structured program, there is a tree of symbol tables
  - Root = outermost block
- But at a given point in the program, only part of the tree is relevant
  - Current block == X
  - Nearest enclosing block == parent(X)
  - Next nearest == parent(parent(X))
  - Etc., up to root

Nested scope operations

- When encounter a new scope during semantic analysis
  - Create a new, empty scope
  - Its parent is the current scope (that of enclosing block)
  - New scope becomes "current"
- When encounter a declaration
  - Add entry to the current scope
  - Check for duplicates in the current scope only (why?)
- When encounter a use
  - Search scopes for declaration: current, its parent, grandparent,…
- When exiting a scope
  - Parent becomes current again

Stack-like

PL/0 symbol table interface

```
class SymTabScope {
public:
  SymTabScope(SymTabScope* enclosingScope);

  void enter(SymTabEntry* newSymbol);
  SymtabEntry* lookup(char* name);
  SymtabEntry* lookup(char* name,
                 SymTabScope*& retScope);
  …
}
```

Implementing nested scopes

- Each scope (instance of `SymTabScope`) keeps a pointer to its enclosing `SymTabScope (_parent)`
- Each scope maintains “down links”, too (`_children`, so we can walk the whole tree)

Symbol tables:Implementation

- Abstractly, it's simple:
  a mapping from names to information, aka key/value pairs
- Concretely, there are lots of choices, each with different performance consequences, e.g.
  - Linked list (or dynamic array)
  - Binary search tree
  - Hash table
- So, we'll take a brief trip down CSE326 memory lane…

Symbol tables: Complexity

	Enter	Lookup	Space cost
A. Linked lists	O(1)		
B. Binary search tree			
C. Hash table			

Symbol tables: Other issues

- Linked lists must have keys that can be compared for equality
- Binary search trees must have keys that can be ordered
- Hash tables must have keys that can be hashed (well)
- Hash table size?

ST:Implementation Summary

- In general
  - Use a hash table for big mappings
  - Use a binary tree or linked list for small mappings
- Ideally, use a self-reorganizing data structure

Types

- Types are abstractions of values that share common properties
  - What operations can be performed on them
  - (Usually) how they are represented in memory
- Types usually guide how compilation proceeds

Taxonomy of types

- Basic/atomic types
  - `int, bool, char, real, string, …`
  - `enum(`$v_1, v_2, \ldots, v_n$`)`
- User-defined types: `Stack, SymTabScope,…`
  - Type constructors
  - Parameterized types
  - Type synonyms

Type constructors

- `ptr(`type`)`
- `array(`index-range, element-type`)`
- `record(`$name_1$:$type_1$, … $name_n$:$type_n$`)`
- `tuple(`$type_1$, …, $type_n$`)` or $type_1 \times \ldots \times type_n$
- `union(`$type_1$, …, $type_n$`)` or $type_1 + \ldots + type_n$
- `function(`arg-types, result-type`)` or
  $type_1 \times \ldots \times type_n \rightarrow$ result-type

Parameterized types

Functions returning types

- `Array<T>`
- `Stack<T>`
- `HashTable<Key,Value>`
- …

Type synonyms

Give alternative name to existing type

- `typedef SymTabScope* SymTabReg`

Type checking

- A key part of language implementation
  - Semantic analysis phase, linking, and/or runtime
- Verifies that operations on values will be legal
  - I.e., they compute values that will be legal in context
- Examples

3 + 4	3 + 4.0
3 + x	3 + `'x'`
3[x]	x[3]
3 + TRUE	`x.y->z*`

Type checking terminology

- Static vs. dynamic typing
  - Static: checked prior to execution (e.g., compile-time)
  - Dynamic: checked during execution
- Strong vs. weak typing
  - Strong: guarantees no illegal operations performed
  - Weak: no such guarantee
- Caveats
  - Hybrids are common
  - Mistaken usages of these terms is common
    - Ex: "untyped", "typeless" could mean "dynamic" or

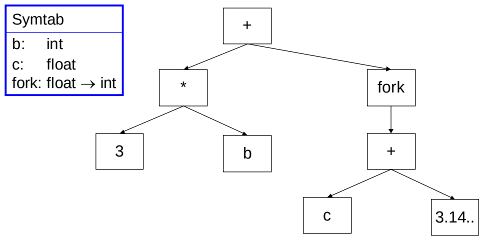
Type checking

- Assume we have an AST for the source program
  - It is syntactically correct
  - The symbol table has been computed
- Does it meet the type constraints of the language?
  - Ex: `a := 3 * b + fork(c + 3.14159)`
    - What are the types of `a`, `b`, and `c`?
    - What type does `fork` return?
    - What type does `fork` accept?
    - What happens when `c` is added to a `float`?
    - What happens when `b` is multiplied by `3`?
    - What happens when `fork`'s result is added to `3 * b`?

Type checking strategy

- Traverse AST recursively, starting at root node
  - Most work is on the bottom-up pass
- At each node
  - Recursively type check any subtrees
  - Check legality of current node, given children's types
  - Compute and return result type (if any) of current node

Ex: 3 * b + fork(c + 3.14159)

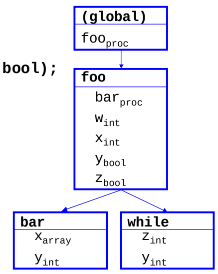


Top-down information also:
From enclosing context

- Need to know types of variables referenced
  - Must pass down symbol table during traversal
- Legality of (e.g.) `break` and `return` statements depends on context: pass down
  - whether in loop,
  - what the result type of the function must be,
  - etc.

PL/0 Type Checking

```
procedure foo(x:int, w:int);
  var z:bool;
  const y:bool = true;
  procedure bar(x:array[5] of bool);
    var y:int;
  begin
    x[y] := z;
  end bar;
begin
  while z do
    var z:int, y:int;
    y := z * x; end;
  output := x + y;
end foo;
```



Representing types in PL/0

```
class Type {
  virtual bool same(Type* t);
  …
};

class IntegerType   : public Type {…};
class BooleanType   : public Type {…};
class ProcedureType : public Type {
   …
   TypeArray* _formalTypes;
};

IntegerType* integerType; // predefined instances
BooleanType* booleanType;
```

PL/0 type checking: overview

```
Type* Expr::typecheck(SymTabScope* s);
void  Stmt::typecheck(SymTabScope* s);
void  Decl::typecheck(SymTabScope* s);

Type* LValue::
         typecheck_lvalue(SymTabScope* s);

int   Expr::resolve_constant(SymTabScope* s);

Type* TypeAST::typecheck(SymTabScope* s);
```

Type checking PL/0 expressions

A simple case: integer literals (like "0" or "-17")

```
Type* IntegerLiteral::typecheck(SymTabScope* s) {
    return integerType;
}
```

Type checking var references

```
Type* VarRef::typecheck(SymTabScope* s) {
    SymTabEntry* ste = s->lookup(_ident);
    if (ste == NULL) {
      char* errormsg = new char[errormsgbuffsize];
      sprintf(errormsg,
              "undeclared var \"%s\" referenced", _ident);
      Plzero->typeError(errormsg, line);
    }
    if (! ste->isConstant() &&
        ! ste->isVariable()) {
      char* errormsg = new char[errormsgbuffsize];
      sprintf(errormsg,"\"%s\" not const or var",_ident);
      Plzero->typeError(errormsg, line);
    }
    return ste->type();
  }
```

Type checking operators

```
Type* BinOp::typecheck(SymTabScope* s) {
    Type* left  = _left->typecheck(s);
    Type* right = _right->typecheck(s);
        switch(_op) {
      case PLUS:case MINUS:case MUL: case LEQ: …
        if (left->different(integerType) ||
            right->different(integerType)) {
          Plzero->typeError("args not ints");
      }
      break;
      case EQL: case NEQ:
        if (left->different(right)) {
          Plzero->typeError("args not same type");
      }
      break;
      default:
        Plzero->fatal("unexpected BINOP");
    }
```

Continued

```
      switch (_op) {
        case PLUS:case MINUS:case MUL:case DIVIDE:
          return integerType;

        case EQL:case NEQ:case LSS:
        case LEQ:case GTR:case GEQ:
          return booleanType;

        default:
          Plzero->fatal("unexpected BINOP");
          return NULL;  // not actually executed
      }
  }
```

Type checking assignments

```
void AssignStmt::typecheck(SymTabScope* s) {
     Type* lhs = _lvalue->typecheck_lvalue(s);
     Type* rhs = _expr->typecheck(s);
     if (lhs->different(rhs)) {
       Plzero->typeError("lhs type differs from rhs");
     }
}
```

Type checking if statements

```
void IfStmt::typecheck(SymTabScope* s) {
  Type* testType = _test->typecheck(s);
  if (testType->different(booleanType)) {
      Plzero->typeError("test not Boolean");
  }

  for (int i = 0;
        i < _then_stmts->length(); i++) {
         _then_stmts->fetch(i)->typecheck(s);
        }
  }
```

Type checking call statements

```
void CallStmt::typecheck(SymTabScope* s) {
  int i;
  TypeArray* argTypes = new TypeArray;
  for (i = 0; i < _args->length(); i++) {
    Type* argType = _args->fetch(i)->typecheck(s);
    argTypes->add(argType);
  }

  SymTabEntry* ste = s->lookup(_ident);
  if (ste == NULL) {
   Plzero->typeError("undeclared procedure");
  }
```

Continued

```
  Type* procType = ste->type();
  if (! procType->isProcedure()) {
    Plzero->typeError("not a procedure");
  }
  TypeArray* formalTypes = procType->formalTypes();
  if (formalTypes->length() != argTypes->length()) {
    Plzero->typeError("call doesn't match proto");
  }
  for (i = 0; i < formalTypes->length(); i++) {
    if (formalTypes->fetch(i)->
               different(argTypes->fetch(i))) {
      Plzero->typeError(…);
    }
  }
  return;         // whew! passed all checks!
```

Type checking declarations

```
void VarDecl::typecheck(SymTabScope* s) {
  for (int i = 0; i < _items->length(); i++) {
    _items->fetch(i)->typecheck(s);
  }
}

void VarDeclItem::typecheck(SymTabScope* s) {
  Type* t = _type->typecheck(s);

  VarSTE* varSTE = new VarSTE(_name, t);
  s->enter(varSTE, line);
}
```

Continued

```
void ConstDecl::typecheck(SymTabScope* s) {
  for (int i = 0; i < _items->length(); i++) {
    _items->fetch(i)->typecheck(s);
  }
}

void ConstDeclItem::typecheck(SymTabScope* s) {
  Type* t    = _type->typecheck(s);
  Type* type = _expr->typecheck(s);
  Value* constant_value = _expr->resolve_constant(s);
  if (t->different(type)) {
    Plzero->typeError(…);
  }

  ConstSTE* constSTE =
      new ConstSTE(_name, t, constant_value);
  s->enter(constSTE, line);
}
```

Continued

```
void ProcDecl::typecheck(SymTabScope* s) {
  SymTabScope* body_scope = new SymTabScope(s);

  TypeArray* formalTypes = new TypeArray;
  for (int i = 0; i < _formals->length(); i++) {
    FormalDecl* formal = _formals->fetch(i);
    Type* t = formal->typecheck(s, body_scope);
    formalTypes->add(t);
  }

  ProcedureType* procType =
     new ProcedureType(formalTypes);

  ProcSTE* procSTE = new ProcSTE(_name, procType);
  s->enter(procSTE, line);  // add to enclosing scope

  _block->typecheck(body_scope); // check in new scope
}
```

Continued

```
void Block::typecheck(SymTabScope* s) {

  for (int i = 0; i < _decls->length(); i++) {
      _decls->fetch(i)->typecheck(s);
  }

  for (int j = 0; j < _stmts->length(); j++) {
    _stmts->fetch(j)->typecheck(s);
  }
}
```

Type checking

- We've covered the basic issues in how to check semantic, type-oriented, properties for the data types and constructs in PL/0 (and some more)

Records

Records (aka structs) group heterogeneous types into a single, usually named, unit

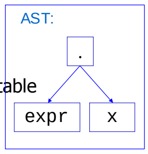
```
record R = begin
  x : int;
  a : array[10] of bool;
  m : char;
end record;

var t : R;
…
r.x
```

Type checking records

- Need to represent record type, including fields of record
- Need to name user-defined record types
- Need to access fields of record values
- May need to handle unambiguous but not fully qualified names (depending on language definition)

An implementation

- Representing record type using a symbol table for fields
  - `class RecordType: public Type {..};`
  - Create `RecordTypeSTE`
- To typecheck `expr.x`
  - Typecheck `expr`
    - Error if not record type
  - Lookup `x` in record type's symbol table
    - Error if not found
  - Extract and return type of `x`



Type checking classes & modules

- A class/module is just like a record, except that it contains procedures in addition to simple variables
- So they are already supported by using a symbol table to store record/class/module fields
- Procedures in the class/module can access other fields of the class/module
  - Already supported: nest procs in record symbol table
- Inheritance?

Type equivalence

- When is one type equal to another?
  - Implemented in PL/0 with `Type::same` function
- It's generally "obvious" for atomic types like `int`, `string`, user-defined types (e.g., point2d vs complex)
- What about type constructors like arrays?

```
var a1    : array[10]  of int;
var a2,a3 : array[10]  of int;
var a4    : array[20]  of int;
var a5    : array[10]  of bool
var a6    : array[0:9] of int;
```

Equivalence, def I: Structural Eq.

- Two types are *structurally equivalent* if they have the same structure
  - If atomic types, then obvious
  - If type constructors
    - Same constructor
    - Recursively, equivalent arguments to constructor
- Implement with recursive `same`

Equivalence, def II: Name Eq.

- Two types are *name equivalent* if they came from the same textual occurrence of a type constructor
- Implement with pointer equality of `Type` instances
- Special case: type synonyms don't define new types

`same` & `different`

```
class Type {
public:
  …
  virtual bool same(Type* t) = 0;
  bool different(Type* t) { return !same(t); }
  …
};
class IntegerType : public Type {
public:
    …
    bool same(Type* t) { return t->isInteger();
}
    …
};
```

Implementing structural equivalence *(details)*

- Problem: want to dispatch on two arguments, not just receiver
  - That is, choose what method to execute based on more than the class of the receiver
- Why? There's a symmetry that the OO dispatch approach skews
  - `if (lhs->different(rhs)) {…error…}`
- Why not: `if (different(lhs,rhs)) {…error…}`

Multi-methods

- Languages that support dispatching on more than one argument provide *multi-methods*
- For example, they might look like
  - `virtual bool same(type* t1, type* t2) {return false;}`
  - `virtual bool same(IntType* t1, IntType* t2) {return true;}`
  - `virtual bool same(ProcType* t1, ProcType* t2) {return same(t1->args,t2->args);}`
- Different from static overloading in C++

Overloading: quick reminder

- Overloading arises when the same operator or function is used to represent distinct operations
  - `3 + 4`
  - `3.14159 + 2.71828`
  - `"mork" + "mindy"`
- The compiler statically decides which "+" to compile to based on the (type) context

Polymorphism: quick reminder

- Polymorphism is different from overloading
- In overloading the same operator means different things in different contexts
- In polymorphism, the same operator works on different types of data
  - (length '(a b c)) vs. (length '((a) (b c) 3 4))
  - (sort '(4 1 2)) vs. (sort '(c g a))
- In polymorphism, the compiler compiles the same code regardless

But C++ has no multi-methods:
So we use double dispatching

```
class Type {
  virtual bool same(Type* t) = 0;
  virtual bool isInteger() {return
false;}
  virtual bool isProc()    {return
false;}
};

class IntegerType : public Type {
  bool same(Type* t){return t-
>isInteger();}
  bool isInteger()  {return true;}
};
```

Type conversions and coercions

- In C, can explicitly convert data of type `float` to data of type `int` (and some other examples)
  - Represent it explicitly as a unary operator
  - Type checking and code generation work as normal
- In C, can also implicitly coerce
  - System must insert unary conversion operators as part of type checking
  - Code generation works as normal

Type casts

- In C, Java (and some others) can explicitly cast an object of one type to another
  - Sometimes a cast means a conversion
    - E.g., casts between numeric types
    - Type-safe, but sometimes entails loss of accuracy
  - Sometimes a cast means just a change of static type without any computation
    - E.g., casts between pointer types
    - Generally NOT type-safe

Safety of casting

- In C, the safety of casts is not checked
  - That is, it's possible to convert into a representation that is illegal for the new type of data
  - Allows writing of low-level code that's type-unsafe
  - More often used to work around limitations in C's static type system
- In Java, downcasts from superclass to subclass include a run-time type check to preserve type safety
  - This is the primary place where Java uses dynamic type checking

Where are we?

- We now know, in principle, how to
  1. take a string of characters
  2. convert it into an AST with associated symbol table
  3. and know that it represents a legal source program (including semantic checks)
- That is the complete set of responsibilities (at a high-level) of the front-end of a compiler

Next…

- …what to do now that we have this wonderful AST representation
- We'll look mostly at interpreting it or compiling it
  - But you could also analyze it for program properties
  - Or you could "unparse" it to display aspects of the program on the screen for users
  - …

Next lecture

- We'll start looking at the implementation issues in symbol tables
  - For instance, how to efficiently manage references to outer scopes
- With a particular focus on how PL/0 does it