378: Machine Organization and Assembly Language

Autumn 2008

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Slides adapted from: Larry Snyder, Luis Gare, Linnee Torrellas, Craig Zilles, and Howard Huang

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What is computer architecture about?

- **Computer architecture** is the study of building computer systems.
- CSE 378 is roughly split into three parts.
  - The first third discusses instruction set architectures—the bridge between hardware and software.
  - Next, we introduce more advanced processor implementations. The focus is on pipelining, which is one of the most important ways to improve performance.
  - Finally, we talk about memory systems, I/O, and how to connect it all together.

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Why should you care?

- It is interesting.
  - You will learn how a processor actually works!

- It will help you be a better programmer.
  - Understanding how your program is translated to assembly code lets you reason about correctness and performance.
  - Demystify the seemingly arbitrary (e.g., bus errors, segmentation faults)

- Many cool jobs require an understanding of computer architecture.
  - The cutting edge is often pushing computers to their limits.
  - Supercomputing, games, portable devices, etc.

- Computer architecture illustrates many fundamental ideas in computer science.
  - Abstraction, caching, and indirection are CS staples

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CSE 370 vs. CSE 378

- This class expands upon the computer architecture material from the last few weeks of CSE 370, and we rely on many other ideas from CSE 370.
  - Understanding binary, hexadecimal and two's-complement numbers is still important.
  - Devices like multiplexers, registers and ALUs appear frequently. You should know what they do, but not necessarily how they work.
  - Finite state machines and sequential circuits will appear again.

  - We do not spend time with logic design topics like Karnaugh maps, Boolean algebra, latches and flip-flops.

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Who we are

- Instructor:
  - Ruth Anderson, rea@cs, Office: CSE 360

- Teaching Assistants:
  - Jacob Nelson, nelson@cs
  - Allan Carroll, allanc@cs
  - Guy Bordelon, bordelon@cs

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Administrivia

- The textbook provides the most comprehensive coverage
  - Computer Organization and Design, Patterson and Hennessy, 3rd Edition

- Lectures will present course material

- Sections will clarify course material and homeworks

- Grading:
  - Lab assignments: 30%
  - Homeworks: 25%
  - Midterm: 20%
  - Final: 25%
  - Participation: 5%

- Getting in touch with us: cse378-tas@cs

- Course webpage:

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Major
2) Information Sheet
Homework for Today!!
- Attend Sections Tomorrow: (Sept 25th)
  - Section AA, Th, 9:30-10:20, MEB 238
  - Section AB, Th, 12:30-1:20, EEB 003
1) Reading in P & H: [Ch 1, Ch 2.1-2.3]
2) Information Sheet: bring to lecture on Friday Sept 26th

Bring to Class on Friday:
- Name
- Email address
- Year (1,2,3,4)
- Major
- Hometown
- Interesting Fact or what I did over summer/break.

Instruction set architectures
- Interface between hardware and software
  - abstraction: hide HW complexity from the software through a set of simple operations and devices
  - add, mul, and, lw, ...

MIPS
- In this class, we’ll use the MIPS instruction set architecture (ISA) to illustrate concepts in assembly language and machine organization
  - Of course, the concepts are not MIPS-specific
  - MIPS is just convenient because it is real, yet simple (unlike x86)
- The MIPS ISA is still used in many places today. Primarily in embedded systems, like:
  - Various routers from Cisco
  - Game machines like the Nintendo 64 and Sony PlayStation 2

From C to Machine Language
High-level language (C)
\( a = b + c; \)
Compiler
Assembly Language (MIPS)
\( \text{add } $16, $17, $18 \)
Assembler
Binary Machine Language (MIPS)
01010111101001101101101...
What you will need to learn soon

- You must become fluent in MIPS assembly:
  - Translate from C to MIPS and MIPS to C
- Example problem: Write a recursive function

```c
int pow(int n, int m) {
    if (m == 1)
        return n;
    return n * pow(n, m-1);
}
```

Translate this into a MIPS assembly language function.

MIPS: register-to-register, three address

- MIPS is a register-to-register, or load/store, architecture.
  - The destination and sources must all be registers.
  - Special instructions, which we'll see soon, are needed to access main memory.
- MIPS uses three-address instructions for data manipulation.
  - Each ALU instruction contains a destination and two sources.
  - For example, an addition instruction $(a + b + c)$ has the form:

```
operation    operands
  add        $a, $b, $c
  destination $d
```

Instruction lengths may be affected, as we'll see in the future.

MIPS register file

- MIPS processors have 32 registers, each of which holds a 32-bit value.
  - Register addresses are 5 bits long.
  - The data inputs and outputs are 32 bits wide.
- More registers might seem better, but there is a limit to the goodness.
  - It's more expensive, because of both the registers themselves as well as the decoders and muxes needed to select individual registers.
  - Instruction lengths may be affected, as we'll see in the future.

```
<table>
<thead>
<tr>
<th>Register</th>
<th>Name</th>
<th>Bits</th>
</tr>
</thead>
<tbody>
<tr>
<td>0-15</td>
<td>$a0-$a15</td>
<td>64</td>
</tr>
<tr>
<td>16-31</td>
<td>$a16-$a31</td>
<td>64</td>
</tr>
</tbody>
</table>
```

Basic arithmetic and logic operations

- The basic integer arithmetic operations include the following:
  - `add` sub `mul` `div`

- And here are a few logical operations:
  - `and` `or` `xor`

- Remember that these all require three register operands; for example:
  - `add $t0, $t1, $t2` # $t0 = $t1 + $t2
  - `mul $s1, $s1, $s0` # $s1 = $s1 * $s0

Larger expressions

- More complex arithmetic expressions may require multiple operations at the instruction set level.
  - $10 = (11 + 12) * (13 - 14)$

```
add $t0, $t1, $t2  # $t0 contains $t1 + $t2
sub $s0, $s1, $s4  # Temporary value $s0 = $t3 - $s4
mul $t0, $t0, $s0  # $t0 contains the final product
```

- Temporary registers may be necessary, since each MIPS instructions can access only two source registers and one destination.
  - In this example, we could re-use `$t3` instead of introducing `$s0`.
  - But be careful not to modify registers that are needed again later.
The ALU instructions we've seen so far expect register operands. How do you get data into registers in the first place?

- Some MIPS instructions allow you to specify a signed constant, or “immediate” value, for the second source instead of a register. For example, here is the immediate add instruction, addi:

  ```
  addi $t0, $t1, 4 # $t0 = $t1 + 4
  ```

  Immediate operands can be used in conjunction with the `and` register to write constants into registers:

  ```
  addi $t0, $0, 4 # $t0 = 4
  ```

- MIPS is still considered a load-store architecture, because arithmetic operands cannot be from arbitrary memory locations. They must either be registers or constants that are embedded in the instruction.

We need more space!

- Registers are fast and convenient, but we have only 32 of them, and each one is just 32-bits wide.
  - That’s not enough to hold data structures like large arrays.
  - We also can’t access data elements that are wider than 32 bits.
- We need to add some main memory to the system!
  - RAM is cheaper and denser than registers, so we can add lots of it.
  - But memory is also significantly slower, so registers should be used whenever possible.
- In the past, using registers wisely was the programmer’s job.
  - For example, C has a keyword “register” that marks commonly-used variables which should be kept in the register file if possible.
  - However, modern compilers do a pretty good job of using registers intelligently and minimizing RAM accesses.

Help us help you.

- Remember the big picture.
  - What are we trying to accomplish, and why?
- Read the textbook.
  - It’s clear, well-organized, and well-written. The diagrams can be complex, but are worth studying. Work through the examples and try some exercises on your own. Read the “Real Stuff” and “Historical Perspective” sections.
- Talk to each other.
  - You can learn a lot from other CSE378 students, both by asking and answering questions. Find some good partners for the homeworks/labs (but make sure you all understand what’s going on).
  - Help us help you.
  - Come to lectures, sections and office hours. Send email or post on the mailing list/Wiki. Ask lots of questions! Check out the web page.

How to Succeed in CSE 378