Caching

Caching is a general technique for exploiting the locality of data reference to make it appear as if there is a large amount of fast memory.

Memory Technology	Typical Access Time	\$/MB ('97)
SRAM	5 - 25ns	\$100 - \$250
DRAM	60 - 120ns	\$5 - \$10
Magnetic Disk	10 ⁷ - 2x10 ⁷ ns	\$0.10 - \$0.20

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Terminology

Locality -- the property of memory references to cluster

Temporal locality -- the tendency of the time intervals
between references to a given address to be small

Spatial locality -- the tendency of the distances between
consecutive memory references to be small

Memory hierarchy -- a characteristic of computer design in which a series of storage technologies are used such that the access time is faster as the memory is closer to the processor and the capacity is larger as the memory is further from the processor add \$4.\$5.\$6

add \$4,\$5,\$6 lcw1 \$f2,0(\$4) lcw1 \$f3,4(\$4) sw \$7,0(\$29) sw \$8,4(\$29) bne \$9,\$0,loop

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Terminology, continued

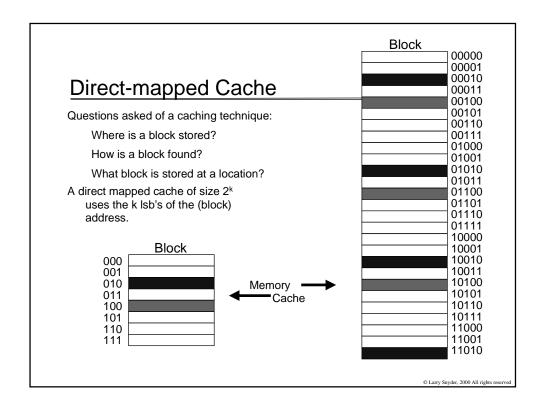
Cache -- memory closest to the processor in a memory hierarchy Caching -- any storage management technique exploiting locality Upper/lower level -- memory closer/further from the processor Block -- unit of memory transfer between two levels in a memory hierarchy. Also called a cache line

Hit/Miss -- accessing data present/not present in a hierarchy level Hit rate -- ratio of hits to total references. Miss rate = 1 - Hit rate Hit time -- time to hit in the cache

Miss penalty -- time to move a block from a lower level in the hierarchy and satisfy the processor's request

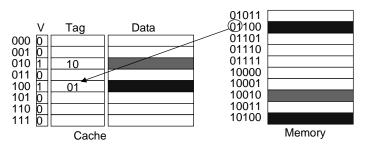
Processor
L1 ____
L2 ____

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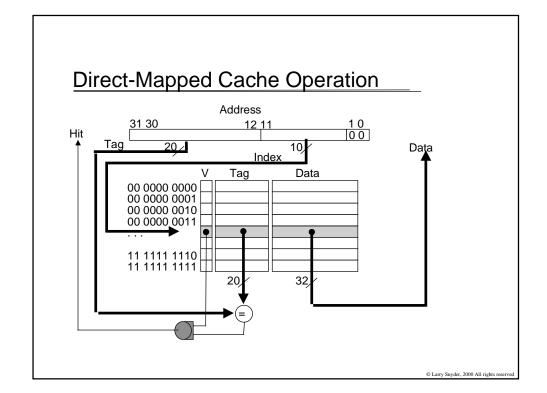


Direct Mapped Cache Fields

- The Tag field stores the msbs of the address.
- The Valid Bit indicates whether the data in the cache block is correct and available.
- The Data field stores the contents of the block.



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Handling Misses

The processor has to stall the instruction that missed; instruction misses stall the pipeline at IF while data misses stall in MEM.

Operations by controller on a miss:

- 1. Compute PC-4
- Access address in main memory and wait for completion.
- 3. Move data to cache, write tag bits, set Valid.
- 4. Restart execution pipeline at the fetch for instruction misses, or MEM for data misses.

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