CSE374 Winter 2016, Final Examination
March 17, 2016

Please do not turn the page until the bell rings.

Rules:

- The exam is closed-book, closed-note, closed-calculator, and closed-electronics.
- Please stop promptly at 10:20.
- There are 92 points total, distributed unevenly among 10 questions (many with multiple parts):

<table>
<thead>
<tr>
<th>Question</th>
<th>Max</th>
<th>Earned</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>9</td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>11</td>
<td></td>
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<tr>
<td>3</td>
<td>9</td>
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<td>4</td>
<td>14</td>
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<td>5</td>
<td>11</td>
<td></td>
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<td>9</td>
<td>10</td>
<td></td>
</tr>
<tr>
<td>10</td>
<td>1</td>
<td></td>
</tr>
</tbody>
</table>

Advice:

- Read questions carefully. Understand a question before you start writing.
- Write down thoughts and intermediate steps so you can get partial credit. But clearly indicate your final answer.
- The questions are not necessarily in order of difficulty. Skip around. Make sure you get to all the problems.
- There is reference material at the end of the exam.
- If you have questions, ask.
- Relax. You are here to learn.
1. (9 points) (Concurrency) Below we give part of the definition of Queue, a class representing a queue data structure (first-in, first-out) of ints. It uses a linked list internally. The class has a private instance mutex variable that it uses to synchronize concurrent calls to some of its methods.

```cpp
typedef struct node {
    int data;
    struct node *next;
} ListNode;

class Queue{
private:
    ListNode *front;
    ListNode *back;
    mutex mtx;
public:
    int peek() {
        if (front == NULL) {
            throw Exception();
        }
        return front->data;
    }

    int dequeue() {
        // a lock_guard causes the mutex to be held for the duration of the function
        lock_guard<mutex> lock(mtx);
        if (front == NULL) {
            // the lock_guard releases the lock when an exception occurs
            throw Exception();
        }
        int val = front->data;
        front = front->next;
        return val;
        // lock variable goes out of scope, so the mutex is released
    }
};
```

Consider a scenario where we have two threads, T1 and T2, which share a variable q, an instance of the Queue class. For each of the following cases where methods of q are called simultaneously, describe a bad interleaving or briefly explain why no bad interleaving is possible. We define a bad interleaving to be an interleaving of method operations that causes an incorrect result (wrong value returned, segmentation fault, etc.).

(a) T1 calls q.peek() and T2 calls q.peek()
(b) $T_1$ calls `q.peek()` and $T_2$ calls `q.dequeue()`

(c) $T_1$ calls `q.dequeue()` and $T_2$ calls `q.dequeue()`

(d) For each case where you described a bad interleaving, briefly explain how to prevent bad interleavings in that case using mutual exclusion.
2. (11 points) (Concurrency) Below we give part of the definition of \texttt{Str}, a class representing a string.

```cpp
class Str {
public:
    int length() const {
        return strlen(st_);
    }

    void append(const Str &s) {
        mtx.lock();
        char *newst = new char[strlen(st_) + strlen(s.st_) + 1];
        strcpy(newst, st_);
        strcat(newst, s.st_);
        delete [] st_;
        st_ = newst;
        mtx.unlock();
    }

private:
    char *st_;
    mutex mtx;
};
```

We are again interested in the consequences of simultaneous calls to certain methods.

(a) In the case where \texttt{append} and \texttt{length} are called simultaneously, there is a data race. Briefly explain where and why.

(b) Briefly explain how to use mutual exclusion to prevent this data race.
(c) In the case where there are two Str instances, a and b, and a thread T1 calls a.append(b), while another thread T2 calls b.append(a), there is at least one data race. Briefly explain where and why.

(d) Explain why the “fix” to append given below won’t work, and discuss an alternative approach.

```cpp
void append(const Str &s) {
    mtx.lock();
    s.mtx.lock();
    char *newst = new char[strlen(st_) + strlen(s.st_) + 1];
    strcpy(newst, st_);
    strcat(newst, s.st_);
    delete [] st_;
    st_ = newst;
    s.mtx.unlock();
    mtx.unlock();
}
```
3. (9 points) (Integer representation) Please complete the following table of 8-bit integer values (use a two’s complement representation for signed values). Powers of two and multiples of 16 are provided as a reference.

<table>
<thead>
<tr>
<th>(2^7)</th>
<th>(2^6)</th>
<th>(2^5)</th>
<th>(2^4)</th>
<th>(2^3)</th>
<th>(2^2)</th>
<th>(2^1)</th>
<th>(2^0)</th>
</tr>
</thead>
<tbody>
<tr>
<td>128</td>
<td>64</td>
<td>32</td>
<td>16</td>
<td>8</td>
<td>4</td>
<td>2</td>
<td>1</td>
</tr>
</tbody>
</table>

\[
16 \times 2 \begin{array}{cccccccccccc}
2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 & 13 & 14 & 15 \\
32 & 48 & 64 & 80 & 96 & 112 & 128 & 144 & 160 & 176 & 192 & 208 & 224 & 240
\end{array}
\]

<table>
<thead>
<tr>
<th>binary</th>
<th>unsigned integer</th>
<th>unsigned hexadecimal integer</th>
<th>signed integer</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
<td>-32</td>
</tr>
<tr>
<td>10011100</td>
<td></td>
<td>0x3F</td>
<td></td>
</tr>
</tbody>
</table>

Space for scratch work:
4. (14 points) (x86-64 assembly)

(a) For each of the following C snippets, give a series of assembly instructions equivalent to that snippet. Relevant register values are given in each case. You can omit width specifiers on instructions (e.g., mov instead of movq). Note that register %rax is used as the return value, so your assembly code must ensure the return value is stored in %rax when the ret instruction executes.

i. C code: `return a * b - c`
   a has type int
   b has type int
   c has type int

   Register state:
   %rdi = value of variable a
   %rsi = value of variable b
   %rdx = value of variable c

   x86-64 assembly:

ii. C code: `return pt->x * pt->x + pt->y * pt->y`
    pt has type Point*
    Point is a typedef for a structure representing a 2D point:

    ```
    typedef struct {
        int x;
        int y;
    } Point;
    ```

    Register state:
    %rdi = value of pt

    x86-64 assembly:
(b) For the following series of assembly instructions, give the equivalent snippet of C. Relevant register values and variable names are given.

x86-64 assembly:

\texttt{mov (\%rdi,\%rsi,4), \%rbx}
\texttt{mov 4(\%rdi,\%rsi,4), \%rcx}
\texttt{mov \%rcx, (\%rdi,\%rsi,4)}
\texttt{mov \%rbx, 4(\%rdi,\%rsi,4)}

<table>
<thead>
<tr>
<th>register</th>
<th>value</th>
</tr>
</thead>
<tbody>
<tr>
<td>%rdi</td>
<td>value of variable array that has type \texttt{int*}</td>
</tr>
<tr>
<td>%rsi</td>
<td>value of variable index that has type \texttt{int}</td>
</tr>
</tbody>
</table>

C code:
5. (11 points) (Bash scripting) Write a shell script `projectsize` that prints the total size in bytes of the C source files in a directory. The command

```
projectsize DIR
```

should print the total size in bytes of the C source files in the directory `DIR`.

Your script should meet the following specification:

1. Print an informative error message to stderr if `DIR` does not exist.
2. If no argument is given, print the total size in bytes of the C source files in the current directory.
3. Treat as a C source file any file that ends in `.c` or `.h`.
4. Ignore subdirectories of `DIR` (or the current directory).

You can use `wc` to get the number of bytes in a file. To test if a file ends in `.c` or `.h`, you can use the `extension-test` command. `man extension-test` provides the following documentation:

```
NAME
   extension-test -- test if a file ends in a specific extension

SYNOPSIS
   extension-test [OPTION]... FILE EXT

DESCRIPTION
   The extension-test utility checks if FILE ends in EXT, and sets its exit status accordingly (similar to `[ ]` command).
   EXT should not include the `.` part of the extension.

   The following options are available:

   -i
     Ignores case when comparing FILE and EXT

   --mult
     Multiple EXT are given, exit with success if FILE ends in any of the given EXT

   -x
     In addition to checking that FILE ends in EXT, also checks that FILE is executable
```

Please write your script on the following page.
Write your answer to problem 5 here.
6. (10 points) (C programming) Your task is to implement the C standard library function \texttt{strcat}. You may implement any helper functions you find useful. Your \texttt{strcat} function must meet the following specification:

1. Matches the declaration \texttt{char \* strcat (char \* destination, char \* source)}
2. Assumes \texttt{destination} and \texttt{source} are valid C strings (i.e., non-NULL, end in null terminator).
3. Appends a copy of the \texttt{source} string to the \texttt{destination} string. The terminating null character in \texttt{destination} is overwritten by the first character of \texttt{source}, and a null-character is included at the end of the new string formed by the concatenation of both in \texttt{destination}.
4. Assumes \texttt{destination} is large enough to contain the concatenated resulting string.
5. Returns \texttt{destination}
6. Does not use any standard library functions (e.g., \texttt{strlen})
7. (8 points) (Toolchain) Recall that there are several steps needed to build an executable program from source files and libraries. Below is a list of several possible errors that can occur when a program is compiled, linked, or executed. For each error, indicate the earliest stage in the process of building and executing the program where it is **always possible** to discover the error and produce some sort of error message or failure. (Note, for example, that some errors can be detected early, say division by 0 if the program contains x/0 in the source code, but in general division by 0 can't be detected until the program is executed if it is dividing x/y and the value of y is not known until runtime.)

Identify where (when) each possible error can definitely be detected. Fill in one of the following codes in the space provided:

- `cpp` - C preprocessor
- `comp` - C compiler
- `ld` - linking/loading step
- `exe` - during program execution
- `can't` - cannot be detected always (including illegal programs that might not actually fail during execution)

_____ calling `free` on a pointer to stack-allocated data
_____ dereferencing a pointer `ptr`, where `ptr == 0x0`
_____ using a variable before it has been declared
_____ making a typo when including a standard library header (e.g., `#include <sgtio.h>`
_____ a function returns a pointer to one of its local variables
_____ when the following code is compiled and/or run (the code below is all the code there is)

```c
int Foo();

int main() {
    int a = Foo();
    return 0;
}
```

_____ the function `char* lookup(ListNode *list, char *word)` is called with the parameters in the wrong order
_____ `strncpy` is used to copy a `char*`, but the null terminator is not copied
8. (9 points) (C programming) You are working on a C data structure for a doubly-linked list (i.e., a linked list that has both next and previous links for each node). Each node stores a struct representing a user. Here are the structs you’re using:

```c
typedef struct {
    int id;
    char *name;
} User;

typedef struct node {
    User *user;
    struct node *prev;
    struct node *next;
} ListNode;
```

The nodes and user data they store are allocated on the heap. Your task is to implement a function to remove a node from the list. Your function must meet the following specification:

1. Matches the declaration `int remove(ListNode *root, int target_id)`
2. Removes the first node after `root` that has a `user` with an `id` equal to `target_id` (root itself will never be removed)
3. Deallocates all memory used by the removed node and returns true.
4. If there are no nodes in the list, or no node has an `id` equal to `target_id`, returns false.
9. (10 points) (C++) Consider the following program, which compiles and executes without errors.

class Rational {
public:
    Rational();
    Rational(int n);
    Rational(int n, int d);
    Rational(const Rational &other);
    ~Rational();
    Rational &operator=(const Rational &other);
private:
    int num;
    int denom;
};

// implementations
Rational::Rational(): num(0), denom(0)
    { cout << "default constructor" << endl; }
Rational::Rational(int n): num(n), denom(0)
    { cout << "one int constructor" << endl; }
Rational::Rational(int n, int d): num(n), denom(d)
    { cout << "two int constructor" << endl; }
Rational::Rational(const Rational &other): num(other.num), denom(other.denom)
    { cout << "copy constructor" << endl; }
Rational::Rational(): num(0), denom(0)
    { cout << "destructor" << endl; }
Rational & Rational::operator=(const Rational &other) {
    cout << "assignment" << endl;
    this->num = other.num;
    this->denom = other.denom;
    return *this;
}

int main() {
Rational *r2 = new Rational(1, 2);
Rational r1();
Rational *r3 = &r1;
r1 = *r2;
Rational r4 = 5;
*r2 = r4;
delete r2;
Rational r5 = r1;
return 0;
}

What output is produced when this program is executed?
10. (1 point) (Haiku) Please compose a haiku about something you learned this quarter. Here’s mine:

inheritance can
cause so many bugs, you are
better off without
Name: ________________________________________

Reference

This is an incomplete list. Just because a command or option is documented here doesn’t mean there is a question that uses it.

Bash

`wc [OPTION]... [FILE]...`
Print newline, word, and byte counts for each FILE, and a total line if more than one FILE is specified.
-`c`, `--bytes`
  print the byte counts
-`l`, `--lines`
  print the newline counts
-`w`, `--words`
  print the words counts

shell scripting

%$(cmd) substitute with the stdout from running cmd
$((expr)) substitute with the result of evaluating expr -- useful for math
$n n$th argument ($0$ is the command itself)
$# number of arguments (does not include $0$)
$@ a list of all the arguments (does not include $0$)
$? the exit status of the most recent command

`test`
discard the first argument ($1$) and move the remaining arguments down ($2$ becomes $1$ and so on, this affects $# and $@).

for item in list_of_things
do
  ...
done

if TEST
then
fi

tests:
[ -d file ] true if file exists and is a directory
[ -e file ] true if file exists, regardless of type
[ -f file ] true if file exists, and is a regular file
[ -n string ] true if length of string is nonzero
[ -z string ] true if length of string is zero
[ s1 = s2 ] true if the strings $s1$ and $s2$ are identical.
[ s1 != s2 ] true if the strings $s1$ and $s2$ are not identical.
[ n1 -eq n2 ] true if integer n1 is equal to integer n2.
  similarly for not equal (-ne), greater than (-gt),
  and less than (-lt)
x86-64 assembly

instructions:

<table>
<thead>
<tr>
<th>instruction</th>
<th>effect</th>
<th>description</th>
</tr>
</thead>
<tbody>
<tr>
<td>mov S, D</td>
<td>( D \leftarrow S )</td>
<td>move</td>
</tr>
<tr>
<td>add S, D</td>
<td>( D \leftarrow D + S )</td>
<td>add</td>
</tr>
<tr>
<td>sub S, D</td>
<td>( D \leftarrow D - S )</td>
<td>subtract</td>
</tr>
<tr>
<td>imul S, D</td>
<td>( D \leftarrow D \times S )</td>
<td>multiply</td>
</tr>
<tr>
<td>ret</td>
<td>returns, uses value in %rax as return value</td>
<td>return</td>
</tr>
</tbody>
</table>

operand forms:

<table>
<thead>
<tr>
<th>type</th>
<th>form</th>
<th>operand value</th>
</tr>
</thead>
<tbody>
<tr>
<td>immediate</td>
<td>$D$</td>
<td>( D )</td>
</tr>
<tr>
<td>register</td>
<td>%R</td>
<td>value stored in register ( R )</td>
</tr>
<tr>
<td>memory</td>
<td>( D )</td>
<td>value at memory location ( D )</td>
</tr>
<tr>
<td>memory</td>
<td>( (%R) )</td>
<td>value at memory location given by value stored in register ( R )</td>
</tr>
<tr>
<td>memory</td>
<td>( D(%R_b, %R_i) )</td>
<td>value at memory location ( D + R_b + R_i )</td>
</tr>
<tr>
<td>memory</td>
<td>( (%R_i, s) )</td>
<td>value at memory location ( R_i \cdot s )</td>
</tr>
<tr>
<td>memory</td>
<td>( D(%R_i, s) )</td>
<td>value at memory location ( D + R_i \cdot s )</td>
</tr>
<tr>
<td>memory</td>
<td>( (%R_b, %R_i, s) )</td>
<td>value at memory location ( R_b + R_i \cdot s )</td>
</tr>
<tr>
<td>memory</td>
<td>( D(%R_b, %R_i, s) )</td>
<td>value at memory location ( D + R_b + R_i \cdot s )</td>
</tr>
</tbody>
</table>

The scaling factor \( s \) must be either 1, 2, 4, or 8.