

# Lecture 2: Implementing ADTs

Data Structures and Algorithms

# Warm Up – Discuss with your neighbors!

#### From last lecture:

- What is an ADT?
- What is a data structure?

#### From CSE 143:

- What is a "linked list" and what operations is it best at?
- What is a "stack" and what operations is it best at?
- What is a "queue" and what operations is it best at?

Socrative: <u>www.socrative.com</u> Room Name: CSE373 Please enter your name as: Last, First

#### Announcements/ Questions

No overloading, wait for drops

Class page to be live tonight

Sections start tomorrow

#### TA Introductions!



Ryan Pham Office Hours: Monday 9:30-11:30 Section: Thursday 1:30



Meredith Wu Office Hours: Friday 1:00 – 3:00pm Section: Thursday 10:30

### Design Decisions

For every ADT there are lots of different ways to implement them Example: List can be implemented with an Array or a LinkedList

Based on your situation you should consider:

- Memory vs Speed
- Generic/Reusability vs Specific/Specialized
- One Function vs Another
- Robustness vs Performance

This class is all about implementing ADTs based on making the right design tradeoffs! > A common topic in interview questions

# Review: "Big Oh"

efficiency: measure of computing resources used by code.

- can be relative to speed (time), memory (space), etc.
- most commonly refers to run time

Assume the following:

- Any single Java statement takes same amount of time to run.
- A method call's runtime is measured by the total of the statements inside the method's body.
- A loop's runtime, if the loop repeats N times, is N times the runtime of the statements in its body.

We measure runtime in proportion to the input data size, N.

- growth rate: Change in runtime as N gets bigger. How does this algorithm perform with larger and larger sets of data?

Say an algorithm runs  $0.4N^3 + 25N^2 + 8N + 17$  statements.

- We ignore constants like 25 because they are tiny next to N.
- The highest-order term (N<sup>3</sup>) dominates the overall runtime.
- We say that this algorithm runs "on the order of" N<sup>3</sup>.
- or O(N<sup>3</sup>) for short ("Big-Oh of N cubed")

# **Review:** Complexity Class

**complexity class:** A category of algorithm efficiency based on the algorithm's relationship to the input size N.

Class	Big-Oh	If you double N,	Example
constant	O(1)	unchanged	Accessing an index of an array
logarithmic	O(log <sub>2</sub> N)	increases slightly	Binary search
linear	O(N)	doubles	Looping over an array
log-linear	O(N log <sub>2</sub> N)	slightly more than doubles	Merge sort algorithm
quadratic	O(N <sup>2</sup> )	quadruples	Nested loops!
exponential	O(2 <sup>N</sup> )	multiplies drastically	Fibonacci with recursion

# **Review:** Case Study: The List ADT

list: stores an ordered sequence of information.

- -Each item is accessible by an index.
- -Lists have a variable size as items can be added and removed

Supported Operations:

- -get(index): returns the item at the given index
- -set(value, index): sets the item at the given index to the given value
- -append(value): adds the given item to the end of the list
- -insert(value, index): insert the given item at the given index maintaining order
- -delete(index): removes the item at the given index maintaining order
- -size(): returns the number of elements in the list

#### List ADT tradeoffs

Time needed to access i-th element:

- <u>Array</u>: O(1) constant time

- LinkedList: O(n) linear time

Time needed to insert at i-th element

- <u>Array</u>: O(n) linear time

- LinkedList: O(n) linear time

Amount of space used overall

- Array: sometimes wasted space

- <u>LinkedList</u>: compact

Amount of space used per element

- <u>Array</u>: minimal

- LinkedList: tiny extra

char[] myArr = new char[5] 0 1 2 3 4 'h' 'e' 'l' 'l' 'o'

LinkedList<Character> myLl = new LinkedList<Character>();



# Thought Experiment

**Discuss with your neighbors:** How would you implement the List ADT for each of the following situations? For each consider the most important functions to optimize.

Situation #1: Write a data structure that implements the List ADT that will be used to store a list of songs in a playlist.

LinkedList

Situation #2: Write a data structure that implements the List ADT that will be used to store the count of students who attend class each day of lecture.

ArrayList

Situation #3: Write a data structure that implements the List ADT that will be used to store the set of operations a user does on a document so another developer can implement the undo function.

# *Review:* What is a Stack?

stack: A collection based on the principle of adding elements and retrieving them in the opposite order.

- Last-In, First-Out ("LIFO")
- Elements are stored in order of insertion.
  - We do not think of them as having indexes.
- Client can only add/remove/examine the last element added (the "top").

#### basic stack operations:

- **push(item)**: Add an element to the top of stack
- **pop()**: Remove the top element and returns it
- **peek()**: Examine the top element without removing it
- **size():** how many items are in the stack?
- isEmpty(): true if there are 1 or more items in stack, false otherwise



stack

### Implementing a Stack with an Array

push(3) push(4) pop() push(5)



numberOfItems = 2

### **Review:** Generics

// a parameterized (generic) class
public class name<TypeParameter> {

- Forces any client that constructs your object to supply a type.
  - Don't write an actual type such as String; the client does that.
  - Instead, write a type variable name such as  ${\rm E}$  (for "element") or  ${\rm T}\,$  (for "type").
  - You can require multiple type parameters separated by commas.
- The rest of your class's code can refer to that type by name.

```
public class Box {
    private Object object;
    public void set(Object object) {
        this.object = object;
    }
    public Object get() {
        return object;
    }
}
```



# Implementing a Generic Stack

# **Review:** What is a Queue?

queue: Retrieves elements in the order they were added.

- First-In, First-Out ("FIFO")
- Elements are stored in order of insertion but don't have indexes.
- Client can only add to the end of the queue, and can only examine/remove the front of the queue.

#### basic queue operations:

- add(item): aka "enqueue" add an element to the back.
- remove(): aka "dequeue" Remove the front element and return.
- **peek()**: Examine the front element without removing it.
- **size():** how many items are stored in the queue?
- isEmpty(): if 1 or more items in the queue returns true, false otherwise



Front of Queue

PAY HERE



Rear of

Queue



#### Implementing a Queue

enqueue(5) enqueue(8) enqueue(9) dequeue()





#### 2 3 4 8 9

3

numberOfItems =

### **Circular Queues**

enqueue(5) enqueue(8) enqueue(9) dequeue()





# TODO list

Fill out survey! -Link on class page

Class webpage to be live tonight: -Skim through full Syllabus on class web page -Sign up for Piazza

-Review 142/143 materials. Materials provided on class webpage