CSE373: Data Structures and Algorithms

Bucket Sort and Radix Sort

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This lecture material represents the work of multiple instructors at the University of Washington. Thank you to all who have contributed!

Sorting: The Big Picture

Surprising amount of neat stuff to say about sorting:

Simple algorithms: $O(n^2)$	Fancier algorithms: $O(n \log n)$	Comparison lower bound: $\Omega(n \log n)$	Specialized algorithms: $O(n)$	Handling huge data sets
Insertion sort Selection sort Shell sort …	Heap sort Merge sort Quick sort …		Bucket sort Radix sort	External sorting

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Radix sort

- Origins go back to the 1890 U.S. census
- Radix = "the base of a number system"
  - Examples will use 10 because we are used to that
  - In implementations use larger numbers
    - For example, for ASCII strings, might use 128
- Idea:
  - Bucket sort on one digit at a time
    - Number of buckets = radix
    - Starting with *least* significant digit
    - Keeping sort *stable*
  - Do one pass per digit
  - Invariant: After *k* passes (digits), the last *k* digits are sorted

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Example

Radix = 10

0	1	2	3	4	5	6	7	8	9
	721		3 143				537 67	478 38	9

Input: 478 537 9 721 3 38 143 67

First pass:
bucket sort by ones digit

Order now: 721 3 143 537 67 478 38 9

Example

Radix = 10

0	1	2	3	4	5	6	7	8	9
	721		3 143				537 67	478 38	9

0	1	2	3	4	5	6	7	8	9
3 9		721	537 38	143		67	478		

Order was: 721 3 143 537 67 478 38 9

Second pass:
stable bucket sort by tens digit

Order now: 3 9 721 537 38 143 67 478

Example

Radix = 10

0	1	2	3	4	5	6	7	8	9
3 9		721	537 38	143		67	478		

0	1	2	3	4	5	6	7	8	9
3 9 38 67	143			478	537		721		

Order was: 3 9 721 537 38 143 67 478

Third pass:
stable bucket sort by 100s digit

Order now: 3 9 38 67 143 478 537 721

Analysis

Input size: n
Number of buckets = Radix: B
Number of passes = “Digits”: P

Work per pass is 1 bucket sort: $O(B+n)$

Total work is $O(P(B+n))$

Compared to comparison sorts, sometimes a win, but often not
- Example: Strings of English letters up to length 15
  - Run-time proportional to: 15*(52 + n)
  - This is less than n log n only if n > 33,000
  - Of course, cross-over point depends on constant factors of the implementations
    - And radix sort can have poor locality properties

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Last Slide on Sorting

- Simple $O(n^2)$ sorts can be fastest for small n
  - Selection sort, Insertion sort (latter linear for mostly-sorted)
  - Good for “below a cut-off” to help divide-and-conquer sorts
- $O(n \log n)$ sorts
  - Heap sort, in-place but not stable nor parallelizable
  - Merge sort, not in place but stable and works as external sort
  - Quick sort, in place but not stable and $O(n^2)$ in worst-case
    - Often fastest, but depends on costs of comparisons/copies
- $\Omega(n \log n)$ is worst-case and average lower-bound for sorting by comparisons
- Non-comparison sorts
  - Bucket sort good for small number of possible key values
  - Radix sort uses fewer buckets and more phases
- Best way to sort? It depends!

Done with sorting! (phew..)

- Moving on….
- There are many many algorithm techniques in the world
  - We’ve learned a few
- What are a few other “classic” algorithm techniques you should at least have heard of?
  - And what are the main ideas behind how they work?