

CSE373: Data Structures and Algorithms

## Implicit Graphs

Steve Tanimoto  
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### Explicit vs Implicit Graphs

- Explicit Graph:
  - All vertices are identified individually and represented separately.
  - All edges are identified individually and represented separately.
- Implicit Graph:
  - Only a subset, possibly only one, of the vertices is given an explicit representation. (The others are implied.)
  - Only a subset, and possibly zero, of the edges is given an explicit representation.
  - A set of "operators" is provided that can be used to construct "new" edges and vertices.

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### Example Explicit Graph

$G = (V, E), V = \{\text{Seattle, Chicago, Boston}\},$   
 $E = \{(\text{Seattle, Chicago}), (\text{Chicago, Boston}), (\text{Boston, Seattle})\}$



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### Example Implicit Graph

$G = (V, E), V = \{v_0, v_1, \dots, v_9\},$   
 $v_0 = \{0\}$

Operators:  $\{\phi_0, \phi_1\}$

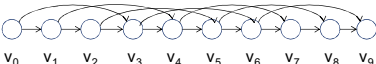
$\phi_0$  precondition:  $i < 9$   
 transition:  $f_0(v_i) = f_0(\{i\}) = \{i+1\}$   
 $v_j = \phi_0(\{i\})$  is a (possibly) new vertex  
 $(v_i, v_j)$  is a (possibly) new edge

$\phi_1$  precondition:  $i < 7$   
 transition:  $f_1(v_i) = f_1(\{i\}) = \{i+3\}$   
 $v_j = \phi_1(\{i\})$  is a (possibly) new vertex  
 $(v_i, v_j)$  is a (possibly) new edge

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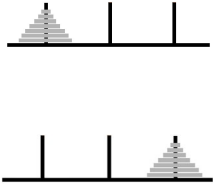
### Some Implicit Graphs Can Be Made Explicit

- $G = (V, E)$
- $V = \{\{0\}, \dots, \{9\}\}$
- $E = \{(\{0\}, \{1\}), (\{1\}, \{2\}), (\{2\}, \{3\}), \dots, (\{8\}, \{9\}), (\{0\}, \{3\}), (\{1\}, \{4\}), (\{2\}, \{5\}), \dots, (\{6\}, \{9\})\}$   $10 + 7 = 17$  edges



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### The Towers of Hanoi (Puzzle)



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### Example: TOH\* Graph

$G = (V, E)$   
 $V_0 = [[4321], [], []]$

$V_k = [p_0, p_1, p_2]; P_m = [d_{m,0}, d_{m,1}, \dots, d_{m,n_m}]$

$\Phi_{i,j}$  precondition :  $p_i \neq []$  and  
 if  $p_j \neq []$  then  $d_{i,n_i} < d_{j,n_j}$

$\Phi_{i,j}$  transition :  
 $p_i: [\alpha, d_{i,n_i}], p_j: [\beta] \Rightarrow p_i: [\alpha], p_j: [\beta, d_{i,n_i}]$

The precondition is that peg i must have at least one disk, and if peg j has any disks, the top (last) disk on peg i must be smaller than the last disk on peg j.

The transition is that the top disk on peg i is removed from peg i and put on the top of the pile on peg j.

\* Towers of Hanoi

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### TOH Graph – properties

$G_n =$  TOH graph for an n-disk TOH puzzle.

$G_n = (V_n, E_n)$   
 What is  $|V_n|$  ?

What is  $|E_n|$  ?

Diameter of  $G_n$  ?

Longest distance between two vertices of the graph.

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### TOH Graph – properties

$G_n =$  TOH graph for an n-disk TOH puzzle. **n=1**

$G_n = (V_n, E_n)$   
 What is  $|V_n|$  ? **3**

What is  $|E_n|$  ? **3**

Diameter of  $G_n$  ? **1**

Longest distance between two vertices of the graph.

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### TOH Graph – properties

$G_n =$  TOH graph for an n-disk TOH puzzle. **n=1, n=2**

$G_n = (V_n, E_n)$   
 What is  $|V_n|$  ? **3 9**

What is  $|E_n|$  ? **3 12**

Diameter of  $G_n$  ? **1 3**

Longest distance between two vertices of the graph.

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### TOH Graph – properties

$G_n =$  TOH graph for an n-disk TOH puzzle. **n=1, n=2, n=3**

$G_n = (V_n, E_n)$   
 What is  $|V_n|$  ? **3 9 27**

What is  $|E_n|$  ? **3 12 39**

Diameter of  $G_n$  ? **1 3 7**

Longest distance between two vertices of the graph.

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### TOH Graph – properties

$G_n =$  TOH graph for an n-disk TOH puzzle. **n=1, n=2, n=3, n=4**

$G_n = (V_n, E_n)$   
 What is  $|V_n|$  ? **3 9 27 81**

What is  $|E_n|$  ? **3 12 39 120**

Diameter of  $G_n$  ? **1 3 7 15**

Longest distance between two vertices of the graph.

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### TOH Graph – properties

$G_n$  = TOH graph for an n-disk TOH puzzle. n=1, n=2, n=3, n=4

$G_n = (V_n, E_n)$

What is  $|V_n|$  ?  
 $|V_n| = 3 \cdot |V_{n-1}|$       3   9   27   81

What is  $|E_n|$  ?      3   12   39   120

Diameter of  $G_n$  ?      1   3   7   15

Longest distance between two vertices of the graph.

### TOH Graph – properties

$G_n$  = TOH graph for an n-disk TOH puzzle. n=1, n=2, n=3, n=4

$G_n = (V_n, E_n)$

What is  $|V_n|$  ?  
 $|V_n| = 3 \cdot |V_{n-1}|$       3   9   27   81

What is  $|E_n|$  ?  
 $|E_n| = 3 \cdot |V_n| + 3$       3   12   39   120

Diameter of  $G_n$  ?      1   3   7   15

Longest distance between two vertices of the graph.

### TOH Graph – properties

$G_n$  = TOH graph for an n-disk TOH puzzle. n=1, n=2, n=3, n=4

$G_n = (V_n, E_n)$

What is  $|V_n|$  ?  
 $|V_n| = 3 \cdot |V_{n-1}|$       3   9   27   81

What is  $|E_n|$  ?  
 $|E_n| = 3 \cdot |V_n| + 3$       3   12   39   120

Diameter of  $G_n$  ?  
 $D(G_n) = 2 \cdot D(G_{n-1}) + 1$       1   3   7   15

Longest distance between two vertices of the graph.

### TOH Graph – properties

$G_n$  = TOH graph for an n-disk TOH puzzle. n=1, n=2, n=3, n=4

$G_n = (V_n, E_n)$

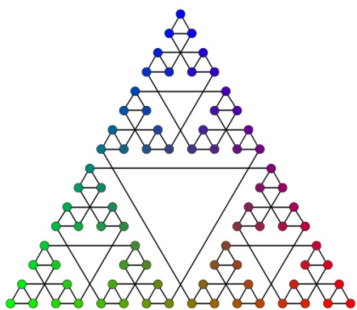
What is  $|V_n|$  ?  
 $|V_n| = 3 \cdot |V_{n-1}|$       3   9   27   81

What is  $|E_n|$  ?  
 $|E_n| = 3 \cdot |V_n| + 3$       3   12   39   120

Diameter of  $G_n$  ?  
 $D(G_n) = 2 \cdot D(G_{n-1}) + 1$       1   3   7   15

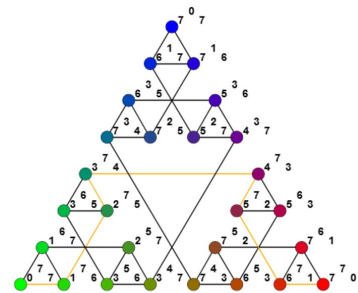
Longest distance between two vertices of the graph.

$n\_disks = 4$



$n\_disks = 3$

With heuristic distances to landmarks.  
 Solution path shown (the golden path)



### Creating Graph Layouts

- Each vertex should be assigned a location on the "page" that helps reflect the structure of the graph.
- One method is based on "landmark vertices."
  - Formulate a distance function  $d(v_0, v_1)$ .
  - Identify 3 or more "landmark" states in Sigma:  $\{L_1, L_2, L_3, \dots, L_n\}$  and assign them (x,y) locations.
  - For vertex  $v$ , use  $d$  to find barycentric coordinates  $(b_1, b_2, \dots, b_n)$  in terms of  $L_1, \dots, L_n$ , and plot  $v$ .

### Heuristic distance $d_h(v_0, v_1)$

Natural method: weight larger disks more heavily. Then add up the absolute values of differences.

$$d_h(v_0, v_1) = \sum_{k=0}^{n-1} 2^k |q_k(v_0, v_1)|$$

Where  $q_k(v_0, v_1) = 1$  if disk  $k$  is on different pegs in  $v_0$  and  $v_1$ , and 0 otherwise.  
 Disk 0 is the smallest disk. ( $n$  is the number of disks.)

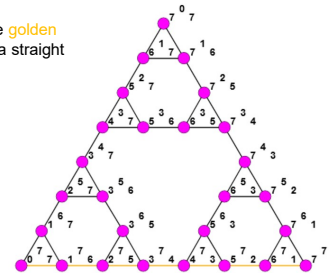
### Layout depends on distance function

Let  $d_m(v_0, v_1)$  = minimum number of moves needed to reach  $v_0$  from  $v_1$ .

Then the solution path can be a straight line.

### TOH Graph Laid Out with $d_m$ and barycentric coordinates.

Note that the golden path is now a straight line.



### Summary

Graphs can be represented explicitly, implicitly, or with a combination of methods.

Some graphs come in families that share properties.  
 e.g., TOH graph  $G_n \in \{G_n \text{ such that } n > 0\}$

Graph layout can reveal structural information.

Problems and puzzles can be represented by graphs.

Path-finding methods can be used to solve problems.