CSE 373 Review Session

String Hash Functions
Splay Trees

```
public static int hash( String key, int tableSize )
    int hashVal = 0;
    for( int i = 0; i < key.length( ); i++ )
        hashVal += key.charAt( i );
    return hashVal % tableSize;
```

```
public static int hash( String key, int tableSize )

return ( key.charAt( 0 ) + 27 * key.charAt( 1 ) +

return ( xey.charAt( 2 ) ) % tableSize;

}
```

```
public static int hash( String key, int tableSize )
8
             int hashVal = 0;
10
11
             for( int i = 0; i < key.length( ); i++ )
12
                 hashVal = 37 * hashVal + key.charAt(i);
13
14
             hashVal %= tableSize;
             if(hashVal < 0)
15
                 hashVal += tableSize;
16
17
18
             return hashVal;
19
```

For a string $s = s_1 s_2 \cdots s_n$

$$h_3(s) = \sum_{i=1}^{n} s_i \cdot 37^{n-i}$$

The *i*th probe for a key k and a hash table of size s:

Linear Probing:

The *i*th probe for a key k and a hash table of size s:

Linear Probing: $h_l(k, i, s) = [h(k) + (i - 1)] \% s$

Quadratic Probing:

The *i*th probe for a key k and a hash table of size s:

Linear Probing:
$$h_l(k, i, s) = [h(k) + (i - 1)] \% s$$

Quadratic Probing:
$$h_q(k, i, s) = [h(k) + (i - 1)^2] \% s$$

Double Hashing:

The *i*th probe for a key k and a hash table of size s:

Linear Probing:
$$h_l(k, i, s) = [h(k) + (i - 1)] \% s$$

Quadratic Probing:
$$h_q(k, i, s) = [h(k) + (i - 1)^2] \% s$$

Double Hashing:
$$h_d(k, i, s) = [h(k) + (i - 1) \cdot g(k, s)] \% s$$

The *i*th probe for a key k and a hash table of size s:

Linear Probing:
$$h_l(k, i, s) = [h(k) + (i - 1)] \% s$$

Quadratic Probing:
$$h_q(k, i, s) = [h(k) + (i - 1)^2] \% s$$

Double Hashing:
$$h_d(k, i, s) = [h(k) + (i - 1) \cdot g(k, s)] \% s$$

Where h and g are hash functions and g is never 0. e.g. g(k,s) = R - (R%s) for a prime number R < s.

Splay Trees

Another type of binary search trees:

Structure property: every node has at most 2 children.

Order property: for every node r, every node in the *left subtree* of r is *smaller* than r, and every node in the *right subtree is bigger* than r.

Splay Trees

Runtime guarantee: every tree operation has an **amortized** runtime of $O(\log n)$ where n is the maximal size of the tree.

That is, starting from an empty tree, every sequence of M consecutive operations takes $O(M \log n)$ time.

Nice properties:

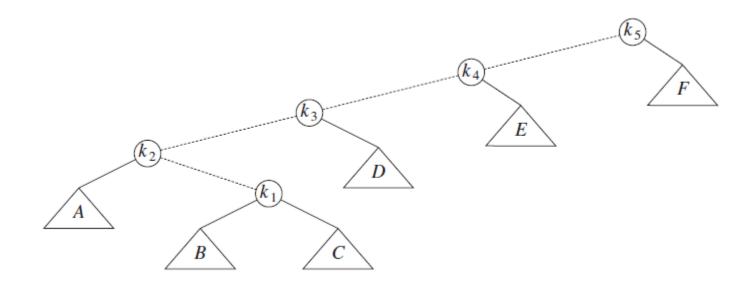
Automatically optimized, such that frequently accessed elements take less time.

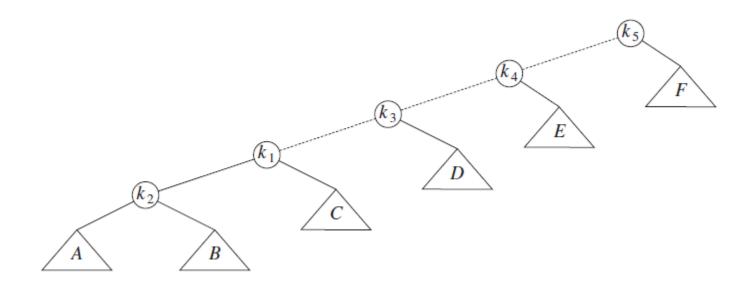
Supports efficient (amortized O(log n)) execution of additional operations, such as merging and splitting around a pivot.

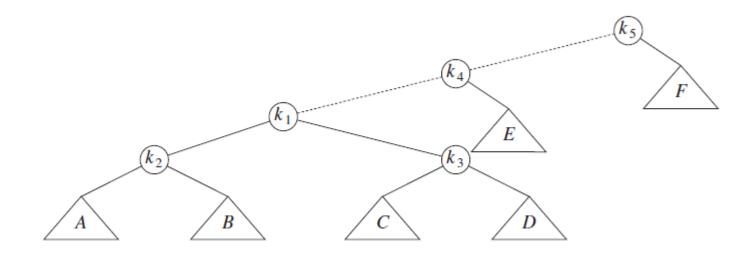
Splay Trees - Operations

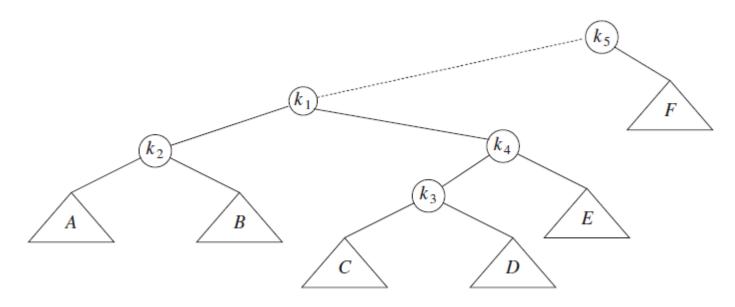
Find: same as a regular BST.

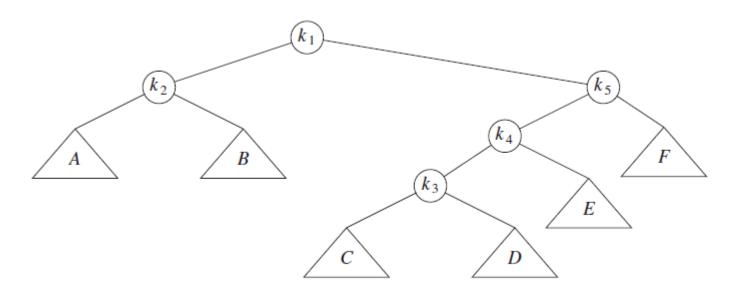
But(!!), in order to maintain runtime guarantees, once the element is found propagate ("splay") it up to the root.

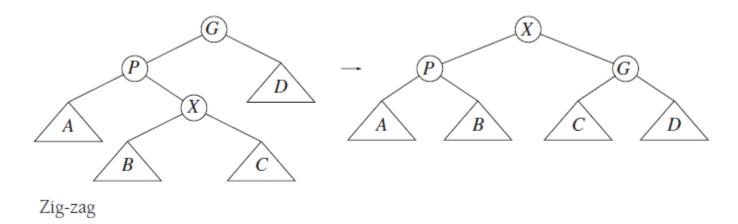


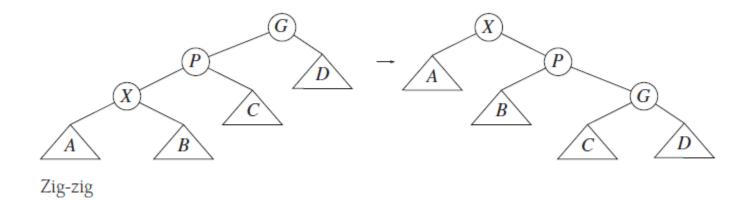


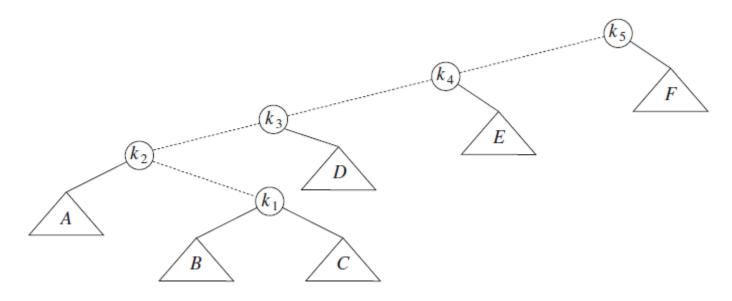


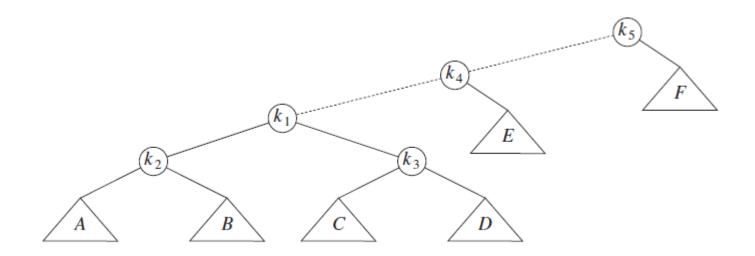


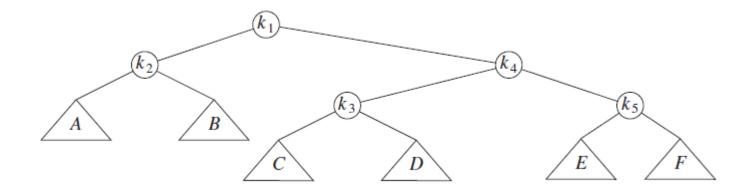












Find - Another Example

