



CSE373: Data Structures and Algorithms

Lecture 2+: Induction Supplemental

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The problem

- Find the sum of the integers from 1 to n
- 1 + 2 + 3 + 4 + ... + (n-1) + n

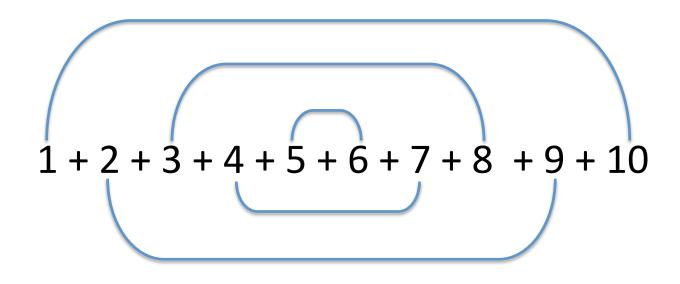
$$\sum_{i=1}^{n} i$$

- For any $n \ge 1$
- Could use brute force, but would be slow
- There's probably a clever shortcut

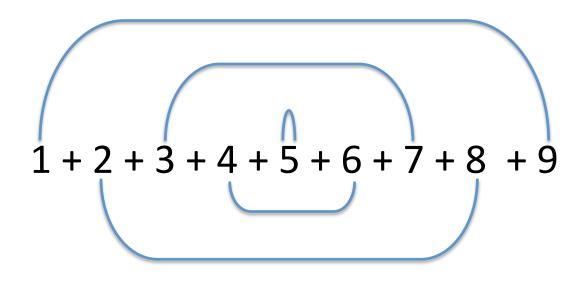
- Shortcut will be some formula involving n
- Compare examples and look for patterns
 - Not something I will ask you to do!
- Start with n = 10:

$$1+2+3+4+5+6+7+8+9+10$$

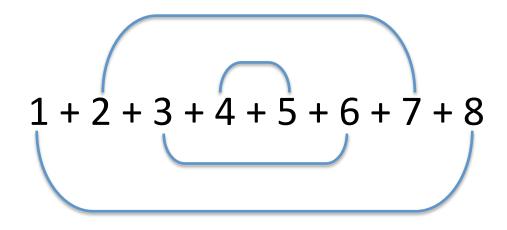
- Large enough to be a pain to add up
- Worthwhile to find shortcut



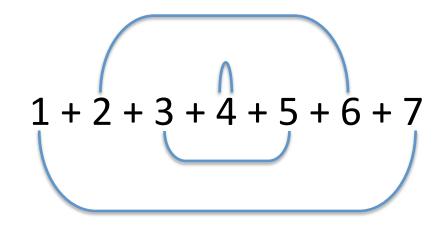
$$= 5 \times 11$$



$$= 4 \times 10 + 5$$



$$= 4 \times 9$$



$$= 3 \times 8 + 4$$

n=7	3×8 + 4
n=8	4×9
n=9	4×10 + 5
n=10	5×11

n=7	3×8 + 4	n is odd
n=8	4×9	n is even
n=9	4×10 + 5	n is odd
n=10	5×11	n is even

When n is even

$$= (n/2) \times (n+1)$$

3×8 + 4	
4×9	n(n+1)/2
4×10 + 5	
5×11	n(n+1)/2

$$= ((n-1)/2) \times (n+1) + (n+1)/2$$

$$= ((n-1)\times(n+1) + (n+1))/2$$

$$= ((n-1+1)\times(n+1))/2$$

$$= (n \times (n+1))/2$$

3×8 + 4	n(n+1)/2
4×9	n(n+1)/2
4×10 + 5	n(n+1)/2
5×11	n(n+1)/2

Are we done?

- The pattern seems pretty clear
 - Is there any reason to think it changes?
- But we want something for any $n \ge 1$
- A mathematical approach is skeptical

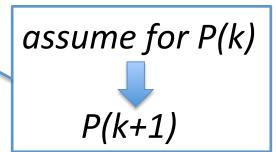
$$n(n+1)$$

2

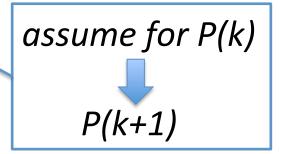
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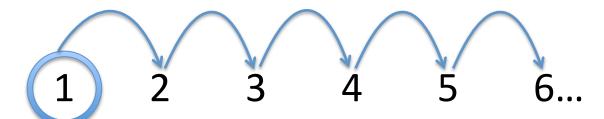
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 - Is there any reason to think it changes?
- But we want something for any $n \ge 1$
- A mathematical approach is skeptical
- All we know is n(n+1)/2 works for 7 to 10
- We must prove the formula works in all cases
 - A rigorous proof

- Type of mathematical proof
 - Sequence of deductive steps
- P(n) = sum of integers from 1 to n
- Two things we need to do
 - Base case ← prove for 1
 - Induction step
- n and k are just variables!



- $P(n) = \text{sum of integers from 1 to } n \text{ (for } n \ge 1)$
- Two things we need to do
 - − Base case prove for 1
 - Induction step





- What we are trying to prove: P(n) = n(n+1)/2
- Base case

$$-P(1)=1$$

$$-1(1+1)/2 = 1(2)/2 = 1(1) = 1$$



- What we are trying to prove: P(n) = n(n+1)/2
- Induction step:
 - Assume true for k
 - -P(k) = k(k+1)/2
 - Now consider P(k+1)

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Why you should care

- Induction turns out to be a useful technique
 - AVL trees
 - Heaps
 - Graph algorithms
 - Can also prove things like $3^n > n^3$ for $n \ge 4$
- Exposure to rigorous thinking