

CSE 351: The Hardware/Software Interface

Section 2

Integer representations, two's complement, and bitwise operators

Introduction

- * CE ugrad (SP14) / 5th year CSE Masters student
 - * Computer architecture, HW/SW Interface, digital design
 - * SAMPA – Approximate Computing / NPU
- * Experience
 - * TA for CSE 351 (WI13) and CSE 352 (AU13)
 - * Amazon
 - * Lockheed Martin Aeronautics
- * OH: Wed 2:30-3:20 in CSE 002, or by appointment
 - * Contact: discussion board or email (wysem@cs)

Integer representations

- * In addition to decimal notation, it's important to be able to understand binary and hexadecimal representations of integers
- * Decimal: 3735928559
 - * No prefix, just the number
- * Binary: 0b11011110101011011011111011101111
 - * "0b" prefix denotes binary notation
- * Hexadecimal: 0xDEADBEEF
 - * "0x" prefix denotes hexadecimal notation
- * Which notation is the most compact of the three?
Why use one over another?

Binary scale

- * Each digit in binary notation is either 0b0 (zero) or 0b1 (one)
- * To convert from (unsigned) binary to decimal notation, take the sum of the n th digit multiplied by 2^{n-1}
- * As an example, $0b1101 = 1 * 2^3 + 1 * 2^2 + 0 * 2^1 + 1 * 2^0 = 8 + 4 + 0 + 1 = 13$

Binary scale

- * To convert from decimal to binary, use a combination of division and modulus to get each digit, tracking the remainder
- * As an example, let's convert 11 to binary
 - * $(11 / 2^0) \% 2 = 1$, so the first digit is 0b1. Remainder is $11 - 1 * 2^0 = 10$
 - * $(10 / 2^1) \% 2 = 5 \% 2 = 1$, so the second digit is 0b1. Remainder is $10 - 1 * 2^1 = 8$
 - * $(8 / 2^2) \% 2 = 4 \% 2 = 0$, so the third digit is 0b0. Remainder is $8 - 0 * 2^2 = 8$
 - * $(8 / 2^3) \% 2 = 1 \% 2 = 1$, so the fourth digit is 0b1
 - * Finally, we have that 11 is 0b1011 in binary

Hexadecimal scale

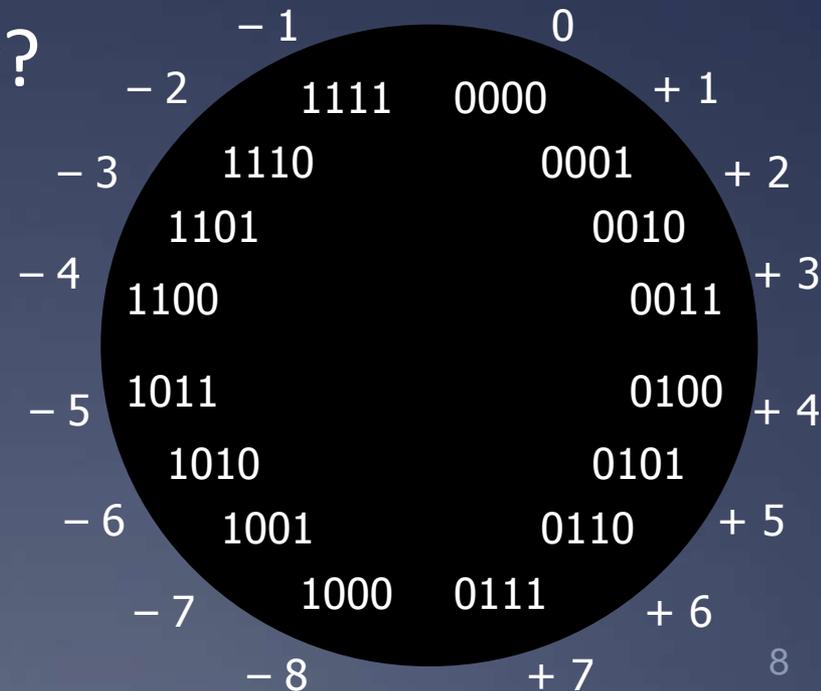
- * Each digit ranges in value from 0x0 (zero) to 0xF (fifteen)
 - * A => ten, B => eleven, C => twelve, D => thirteen, E => fourteen, F => fifteen
- * To convert from (unsigned) hexadecimal to decimal notation, take the sum of the n th digit multiplied by 16^{n-1}
 - * As an example, $0xACE = 0xA * 16^2 + 0xC * 16^1 + 0xE * 16^0 = 10 * 256 + 12 * 16 + 14 = 2766$

Hexadecimal scale

- * The decimal to hexadecimal conversion is the same process as decimal to binary except with 2 instead of 16
- * As an example, let's convert 3254 to hexadecimal
 - * $(3254 / 16^0) \% 16 = 6$, so first digit is 0x6. Remainder is $3254 - 0x6 * 16^0 = 3248$
 - * $(3248 / 16^1) \% 16 = 203 \% 16 = 11 = 0xB$, so second digit is 0xB. Remainder is $3248 - 0xB * 16^1 = 3248 - 176 = 3072$
 - * $(3072 / 16^2) \% 16 = 12 \% 16 = 12 = 0xC$, so third digit is 0xC
 - * Finally, we have that 3254 is 0xCB6 in hexadecimal
- * If we were to write a program to convert from decimal to binary or to hexadecimal, how could we compute the n th digit efficiently using bitwise operators and modulus (%)?

Two's complement review

* In class, we established that two's complement is a nice format for representing signed integers for a couple different reasons. What were they?

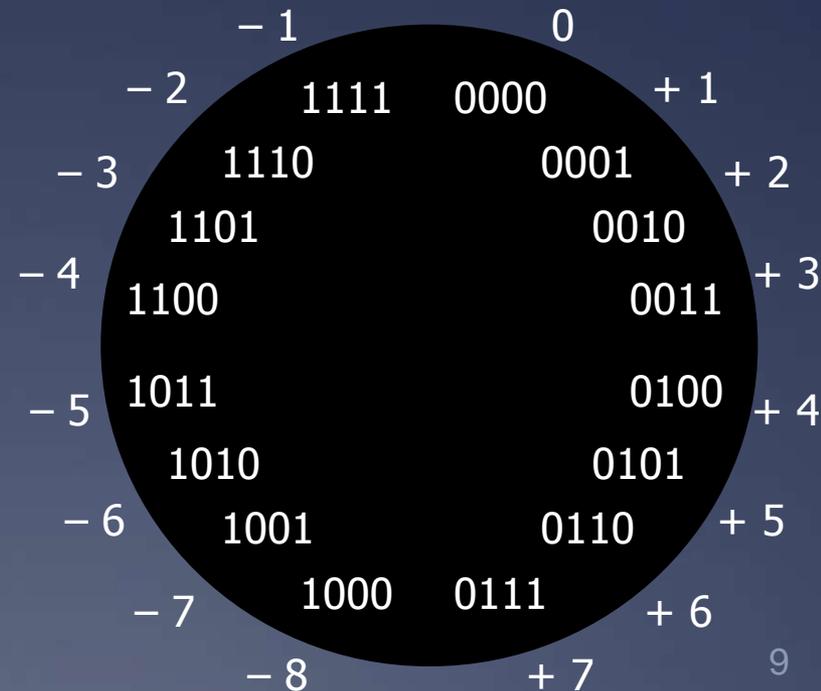


Two's complement review

- * Let's say that we want to encode -5 in binary using two's complement form and four bits
 - * With four bits, the highest bit has a negative weight of 2^3 , so 0b1000 = -8

- * $-5 = -8 + 2 + 1$
$$= 1 * -2^3 + 0 * 2^2 + 1 * 2^1 + 1 * 2^0$$
$$= 10b1011$$

- * $5 = 4 + 1$
$$= 0 * -2^3 + 1 * 2^2 + 0 * 2^1 + 1 * 2^0$$
$$= 0b0101$$



Operator review

- * \sim is arithmetic not (flip all bits)
 - * Example: $\sim 0b1010 = 0b0101$
- * $!$ is logical not (1 if 0b0, else 0)
 - * Example: $!0b100 = 0$, $!0b0 = 1$
- * $\&$ is bitwise and
 - * Example: $0b101 \& 0b110 = 0b100$
- * $|$ is bitwise or
 - * Example: $0b101 | 0b100 = 0b101$
- * \gg is bitwise right shift
 - * Example: $0b1010 \gg 1 = 0b1101$, $0b0101 \gg 1 = 0b0010$
- * \ll is bitwise left shift
 - * Example: $0b1010 \ll 1 = 0b0100$, $0b1000 \ll 1 = 0b0000$

Operator uses

- * Can express negation in terms of arithmetic not and addition
 - * For example, $\sim 4 + 1 = \sim 0b0100 + 1 = 0b1011 + 1 = -5 + 1 = -4$
- * Can use shifting, bitwise and, and logical not to detect if a particular bit is set
 - * As a simple example, $!(x \& (0x1 \ll 1))$ evaluates to 1 if the second bit is set in x and 0 otherwise
 - * Useful for checking if a value is negative
- * Can implement ternaries ($x = _ ? _ : _$) using bitwise and, bitwise or, and arithmetic not
 - * This has wide-ranging applications in lab 1

Bitwise operators in practice

- * Is what we're learning ever useful in practice?
 - * Thankfully (or not, depending on how you look at it), it is
 - * Setting bits in permission strings
 - * For example, to choose the permissions for chmod using octal codes
 - * `chmod 744 <file> = chmod u+rwx,g+r,o+r`

Packing and unpacking

- * Let's say that you have values x , y , and z that take 3, 4, and 1 bit to represent, respectively
- * Is there a way to store these three values using only eight bits?
- * In C, we can define a struct that specifies the width in bits of each value
 - * ...though the compiler will add padding to make the struct a certain size if you don't do so yourself
- * In Java, there are no structs, and we have to use bitwise operators

Packing and unpacking (C)

```
#include <stdio.h>

typedef struct {
    int x : 3;
    int y : 4;
    int z : 1;
    int padding : 24;
} Flags;

int main(int argc, char* argv[]) {
    Flags flags = {3, 8, 1, 0x8fffffff};
    printf("sizeof(flags) is %ju and it stores 0x%x\n",
           sizeof(flags), *(int*) &flags);

    return 0;
}
```

Packing and unpacking (Java)

```
// Pack some values into a byte
byte bitValue = 0;
bitValue |= 3;
bitValue |= 8 << 3;
bitValue |= 1 << 7;
```

```
// Unpack the values from the byte
byte x = bitValue & 0x7;
byte y = bitValue & 0x78;
byte z = bitValue & 0x80;
```

```
// Alternatively, we could have shifted a particular
// mask instead, e.g. (0x1 << 7) instead of 0x80
```

Lab 1 hints

- * Decompose each problem into smaller problems
- * If you are stuck on how to solve something, write it as a combination of functions and boolean logic
 - * Over time, replace each function or boolean operator with a combination of permitted operators
- * Hint for detecting overflow: what is the sign of the integer produced by adding TMax to a positive value? What about when adding negative numbers?
- * Hint for counting bits: consider multiple bits at once. 40 operations isn't enough to check each individually