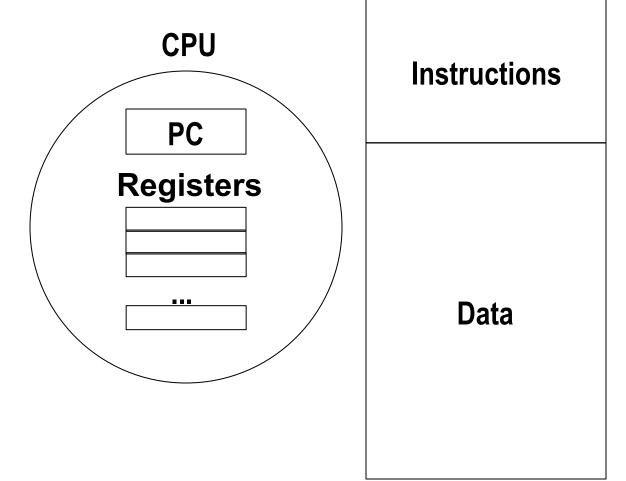
#### Instruction Set Architectures

- ISAs
- Brief history of processors and architectures
- C, assembly, machine code
- Assembly basics: registers, operands, move instructions

# What should the HW/SW interface contain?

### The General ISA



Memory

# **General ISA Design Decisions**

#### Instructions

- What instructions are available? What do they do?
- How are then encoded?

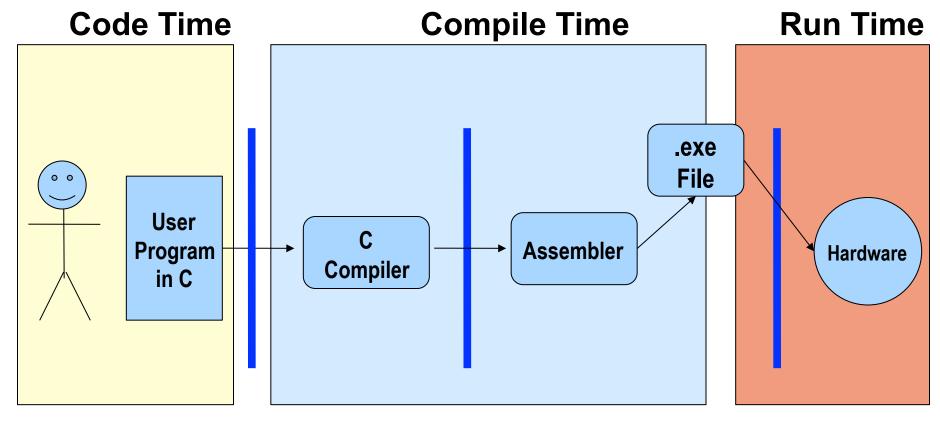
#### Registers

- How many registers are there?
- How wide are they?

#### Memory

How do you specify a memory location?

# HW/SW Interface: Code / Compile / Run Times



What makes programs run fast?

# **Executing Programs Fast!**

- The time required to execute a program depends on:
  - The program (as written in C, for instance)
  - The compiler: what set of assembler instructions it translates the C program into
  - The ISA: what set of instructions it made available to the compiler
  - The hardware implementation: how much time it takes to execute an instruction
- There is a complicated interaction among these

#### Intel x86 Processors

#### Totally dominate the server/laptop market

#### Evolutionary design

- Backwards compatible up until 8086, introduced in 1978
- Added more features as time goes on

#### Complex instruction set computer (CISC)

- Many different instructions with many different formats
  - But, only small subset encountered with Linux programs
- Hard to match performance of Reduced Instruction Set Computers (RISC)
- But, Intel has done just that!

#### Intel x86 Evolution: Milestones

Name Date Transistors MHz

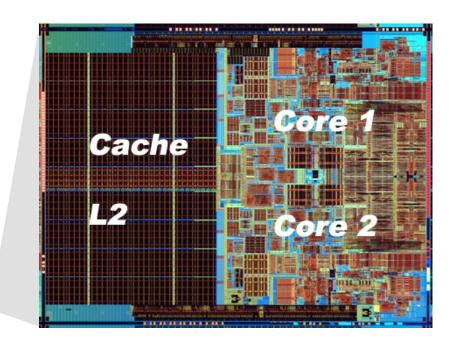
■ 8086 1978 29K 5-10

- First 16-bit processor. Basis for IBM PC & DOS
- 1MB address space
- 386 1985 275K 16-33
  - First 32 bit processor, referred to as IA32
  - Added "flat addressing"
  - Capable of running Unix
  - 32-bit Linux/gcc uses no instructions introduced in later models
- Pentium 4F 2005 230M 2800-3800
  - First 64-bit processor
  - Meanwhile, Pentium 4s (Netburst arch.) phased out in favor of "Core" line

## Intel x86 Processors, contd.

#### Machine Evolution

<b>486</b>	1989	1.9M
<ul><li>Pentium</li></ul>	1993	3.1M
Pentium/MMX	1997	4.5M
<ul><li>PentiumPro</li></ul>	1995	6.5M
Pentium III	1999	8.2M
Pentium 4	2001	42M
Core 2 Duo	2006	291M



#### Added Features

- Instructions to support multimedia operations
  - Parallel operations on 1, 2, and 4-byte data, both integer & FP
- Instructions to enable more efficient conditional operations

#### Linux/GCC Evolution

Very limited impact on performance --- mostly came from HW.

# x86 Clones: Advanced Micro Devices (AMD)

#### Historically

- AMD has followed just behind Intel
- A little bit slower, a lot cheaper

#### Then

- Recruited top circuit designers from Digital Equipment and other downward trending companies
- Built Opteron: tough competitor to Pentium 4
- Developed x86-64, their own extension to 64 bits

#### Recently

- Intel much quicker with dual core design
- Intel currently far ahead in performance
- em64t backwards compatible to x86-64

#### Intel's 64-Bit

- Intel Attempted Radical Shift from IA32 to IA64
  - Totally different architecture (Itanium)
  - Executes IA32 code only as legacy
  - Performance disappointing
- AMD Stepped in with Evolutionary Solution
  - x86-64 (now called "AMD64")
- Intel Felt Obligated to Focus on IA64
  - Hard to admit mistake or that AMD is better
- 2004: Intel Announces EM64T extension to IA32
  - Extended Memory 64-bit Technology
  - Almost identical to x86-64!
- Meanwhile: EM64t well introduced, however, still often not used by OS, programs

# Our Coverage in 351

#### ■ IA32

The traditional x86

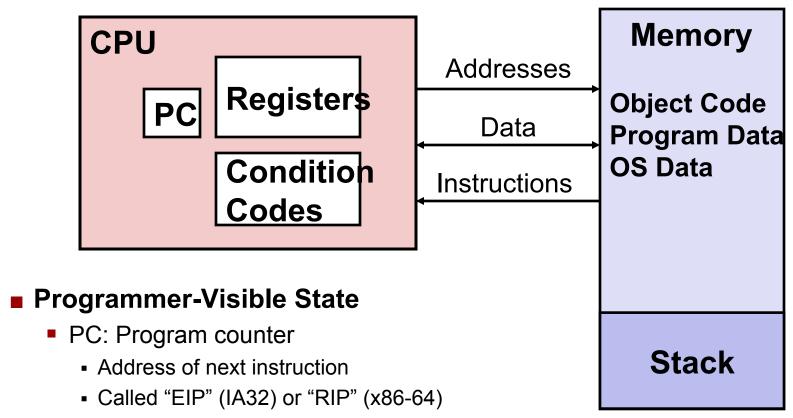
#### x86-64/EM64T

The emerging standard – we'll just touch on its major additions

#### **Definitions**

- Architecture: (also instruction set architecture or ISA) The parts of a processor design that one needs to understand to write assembly code ("what is directly visible to SW")
- Microarchitecture: Implementation of the architecture
- Is cache size "architecture"?
- How about core frequency?
- And number of registers?

# **Assembly Programmer's View**



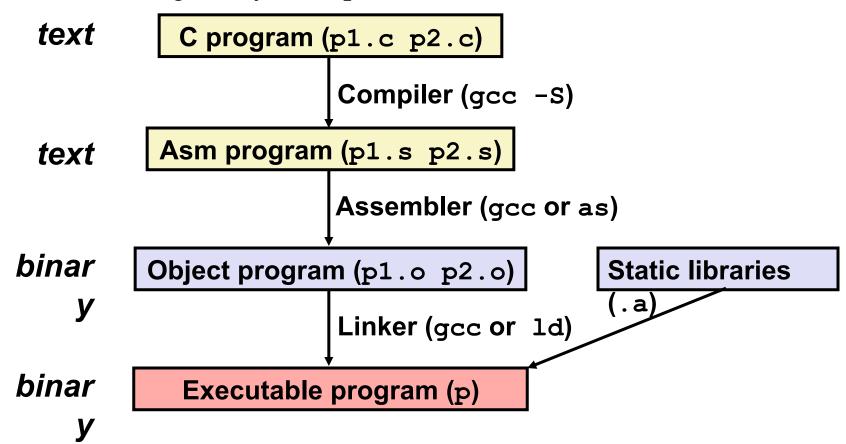
- Register file
  - Heavily used program data
- Condition codes
  - Store status information about most recent arithmetic operation
  - Used for conditional branching

#### Memory

- Byte addressable array
- Code, user data, (some) OS data
- Includes stack used to support procedures (we'll come back to that)

# **Turning C into Object Code**

- Code in files p1.c p2.c
- Compile with command: gcc -0 p1.c p2.c -o p
  - Use optimizations (-o)
  - Put resulting binary in file p



## **Compiling Into Assembly**

#### C Code

```
int sum(int x, int y)
{
  int t = x+y;
  return t;
}
```

#### **Generated IA32 Assembly**

```
sum:
   pushl %ebp
   movl %esp,%ebp
   movl 12(%ebp),%eax
   addl 8(%ebp),%eax
   movl %ebp,%esp
   popl %ebp
   ret
```

#### **Obtain with command**

gcc -0 -S code.c

Produces file code.s

#### **Three Kinds of Instructions**

- Perform arithmetic function on register or memory data
- Transfer data between memory and register
  - Load data from memory into register
  - Store register data into memory
- Transfer control (control flow)
  - Unconditional jumps to/from procedures
  - Conditional branches

# **Assembly Characteristics: Data Types**

- "Integer" data of 1, 2, or 4 bytes
  - Data values
  - Addresses (untyped pointers)
- Floating point data of 4, 8, or 10 bytes
- No aggregate types such as arrays or structures
  - Just contiguously allocated bytes in memory

# **Object Code**

#### Code for sum

```
0x401040 <sum>:
    0x55
    0x89
    0xe5
    0x8b
    0x45
    0x0c
    0x03
    0x45
    0x08
```

0x89

0xec

0x5d

0xc3

- Total of 13 bytes
- Each instruction 1,2, or 3 bytes
- Starts at address 0x401040

#### Assembler

- Translates .s into .o
- Binary encoding of each instruction
- Nearly-complete image of executable code
- Missing linkages between code in different files

#### Linker

- Resolves references between files
- Combines with static run-time libraries
  - E.g., code for malloc, printf
- Some libraries are dynamically linked
  - Linking occurs when program begins execution

## **Example**

```
int t = x+y;
```

```
addl 8(%ebp),%eax
```

#### Similar to expression:

$$x += y$$

#### More precisely:

```
int eax;
int *ebp;
eax += ebp[2]
```

0x401046: 03 45 08

#### C Code

Add two signed integers

#### Assembly

- Add 2 4-byte integers
  - "Long" words in GCC speak
  - Same instruction whether signed or unsigned
- Operands:

x: Register %eax

y: Memory **M**[%ebp+8]

t: Register %**eax** 

-Return function value in %eax

#### Object Code

- 3-byte instruction
- Stored at address 0x401046

## Disassembling Object Code

#### **Disassembled**

```
00401040 < sum>:
  0:
         55
                              %ebp
                        push
  1:
      89 e5
                              %esp, %ebp
                        mov
  3: 8b 45 0c
                              0xc(%ebp),%eax
                        mov
  6: 03 45 08
                        add
                              0x8(%ebp), %eax
  9:
     89 ec
                        mov
                              %ebp,%esp
     5d
  b:
                              %ebp
                        pop
        c3
  c:
                        ret
  d:
         8d 76 00
                        lea
                              0x0 (%esi), %esi
```

#### Disassembler

objdump -d p

- Useful tool for examining object code
- Analyzes bit pattern of series of instructions
- Produces approximate rendition of assembly code
- Can be run on either a .out (complete executable) or .o file

# **Alternate Disassembly**

#### **Object**

# 0x401040: 0x55 0x89 0xe5 0x8b 0x45 0x0c 0x03 0x45 0x08 0x89 0xec 0x5d 0xc3

#### **Disassembled**

```
0x401040 <sum>:
                    push
                            %ebp
0x401041 < sum + 1>:
                            %esp,%ebp
                    mov
0x401043 < sum + 3>:
                            0xc(%ebp),%eax
                    mov
0x401046 < sum + 6>:
                    add
                            0x8(%ebp), %eax
0x401049 < sum + 9>:
                            %ebp,%esp
                    mov
0x40104b <sum+11>:
                    pop
                            %ebp
0x40104c < sum + 12>:
                    ret
0x40104d <sum+13>:
                            0x0(%esi),%esi
                    lea
```

#### Within gdb Debugger

```
gdb p
disassemble sum
```

Disassemble procedure

```
x/13b sum
```

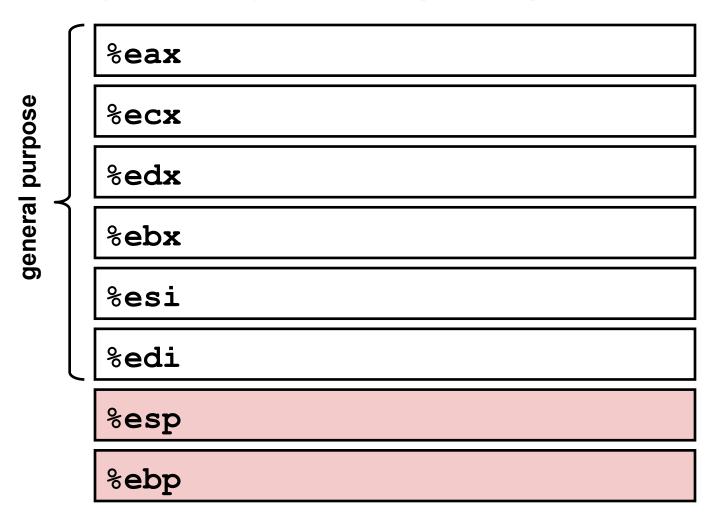
Examine the 13 bytes starting at sum

#### What Can be Disassembled?

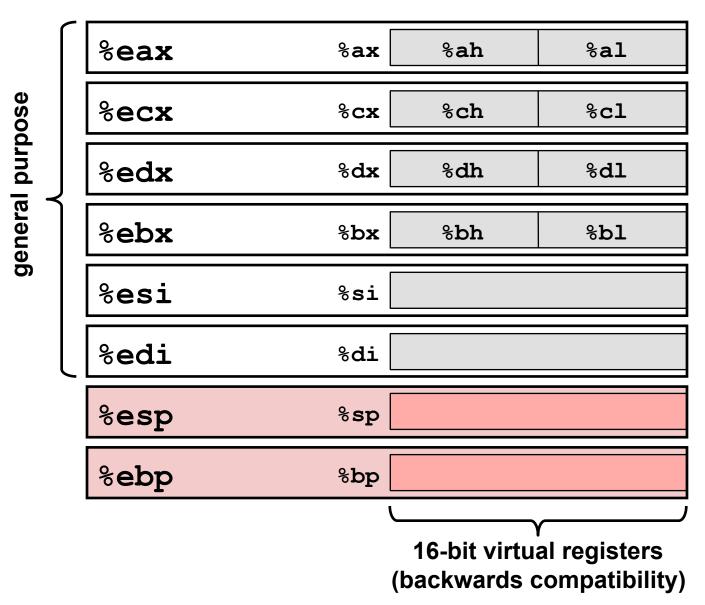
```
% objdump -d WINWORD.EXE
WINWORD.EXE: file format pei-i386
No symbols in "WINWORD.EXE".
Disassembly of section .text:
30001000 <.text>:
30001000: 55
                         push
                               %ebp
30001001: 8b ec
                               %esp,%ebp
                         mov
30001003: 6a ff
                               $0xffffffff
                      push
30001005: 68 90 10 00 30 push
                               $0x30001090
3000100a: 68 91 dc 4c 30 push
                                $0x304cdc91
```

- Anything that can be interpreted as executable code
- Disassembler examines bytes and reconstructs assembly source

# Integer Registers (IA32)



# Integer Registers (IA32)



# Origin (mostly obsolete)

accumulate

counter

data

base

source index

destination index

stack pointer base pointer

# x86-64 Integer Registers

%rax	%eax	%r8	% <b>r8d</b>
%rbx	%ebx	%r9	% <b>r9d</b>
%rcx	%ecx	%r10	%r10d
%rdx	%edx	%r11	%r11d
%rsi	%esi	%r12	%r12d
%rdi	%edi	%r13	%r13d
%rsp	%esp	%r14	%r14d
%rbp	%ebp	%r15	%r15d

- Twice the number of registers
- Accessible as 8, 16, 32, 64 bits

# x86-64 Integer Registers: Usage Conventions

%rax	Return value
%rbx	Callee saved
%rcx	Argument #4
%rdx	Argument #3
%rsi	Argument #2
%rdi	Argument #1
%rsp	Stack pointer
%rbp	Callee saved

% <b>r8</b>	Argument #5
%r9	Argument #6
%r10	Caller saved
%r11	Caller Saved
%r12	Callee saved
%r13	Callee saved
%r14	Callee saved
%r15	Callee saved