The Hardware/Software Interfacecse351

Autumn20111st Lecture, September 28

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Teaching Assistants:

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Who is Luis?



PhD in architecture, multiprocessors, parallelism, compilers.





Who are you?

- 85+ students (wow!)
- Who has written programs in assembly before?
- Written a threaded program before?
- What is hardware? Software?
- What is an interface?
- Why do we need a hardware/software interface?

C vs. Assembler vs. Machine Programs

if (x != 0) y = (y+z) / x;

cmpl \$0, -4(%ebp)

je .L2

movl -12(%ebp), %eax

movl -8(%ebp), %edx

leal (%edx,%eax), %eax

movl %eax, %edx

sarl \$31, %edx

idivl -4(%ebp)

movl %eax, -8(%ebp)

.L2:

b111010000011000

10001011010001000010010000010100

10001011010001100010010100010100

100011010000010000000010

1000100111000010

1100000111111101000011111

111101110111111000010010000011100

10001001010001000010010000011000

C vs. Assembler vs. Machine Programs

if (x!=0)y=(y+z)/x;

cmpl \$0, -4(%ebp)

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movl -12(%ebp), %eax

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b111010000011000

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1000100111000010

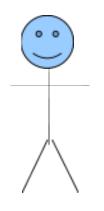
- The three program, frægments are equivalent
- You'd rather write C!
 The hardware likes bit strings!

111101110111111000010010000011100

o The machine in the bits required actually much shorter than the bits required

HW/SW Interface: The Historical Perspective

- Hardware started out quite primitive
 - Design was expensive □ the instruction set was very simple
 - E.g., a single instruction can add two integers
- Software was also very primitive

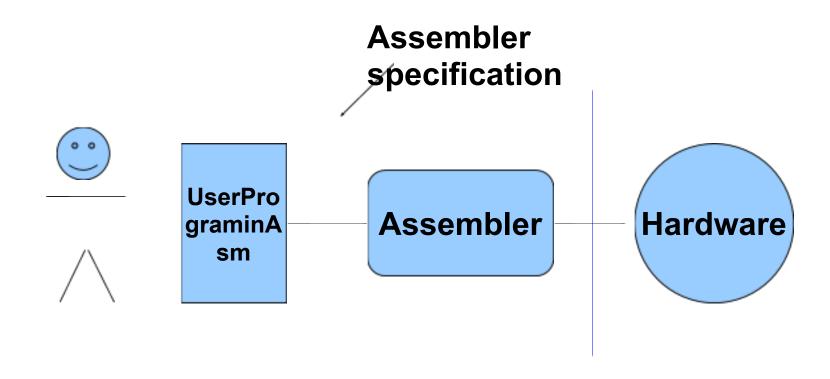


Architecture Specification (Interface)

Hardware

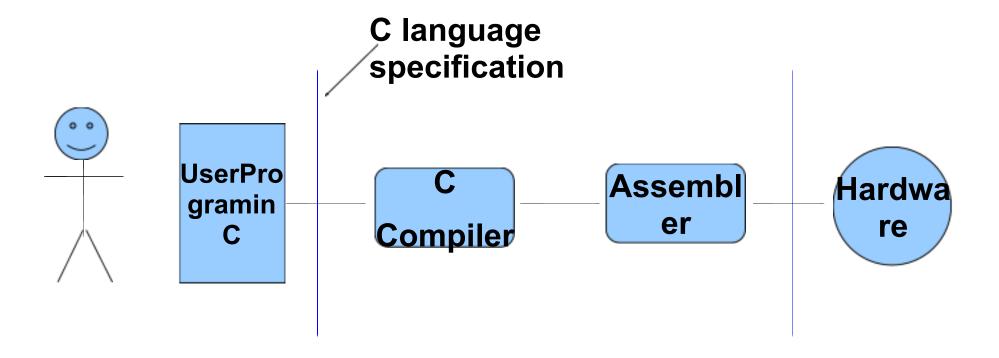
HW/SW Interface: Assemblers

- Life was made a lot better by assemblers
 - 1 assembly instruction = 1 machine instruction, but...
 - different syntax: assembly instructions are character strings, not bit strings

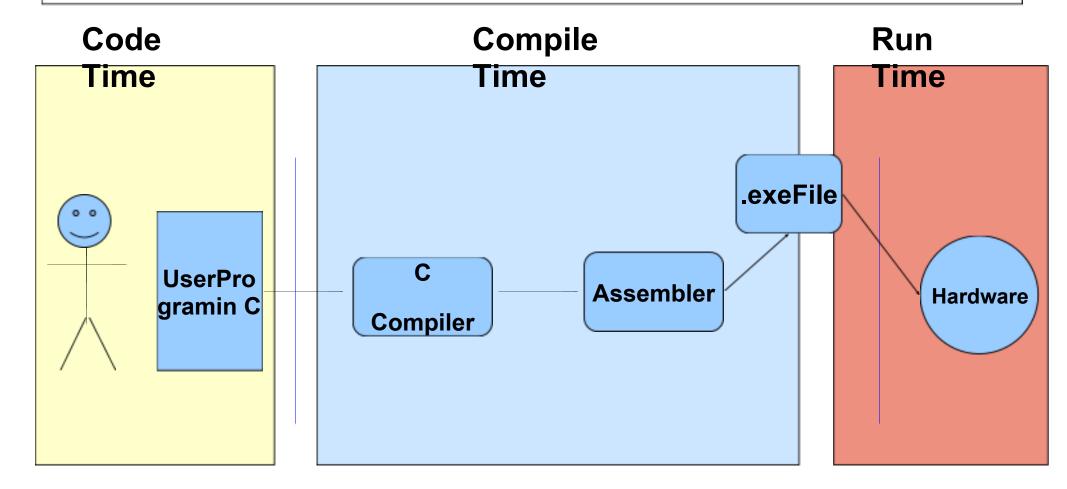


HW/SW Interface: Higher Level Languages (HLL's)

- Highér level of abstraction:
 - 1 HLL line is compiled into many (many) assembler lines



HW/SW Interface: Code / Compile / Run Times



Note: The compiler and assembler are just programs, developed using this same process.

Overview

- Course themes: big and little
- Four important realities
- How the course fits into the CSE curriculum
- Logistics
- HW0 released! Have fun!
- (ready? □)

The Big Theme

- THE HARDWARE/SOFTWARE INTERFACE
- How does the hardware (0s and 1s, processor executing instructions) relate to the software (Java programs)?
- Computing is about abstractions (but don't forget reality)
- What are the abstractions that we use?
- What do YOU need to know about them?
 - When do they break down and you have to peek under the hood?
 - What bugs can they cause and how do you find them?
- Become a better programmer and begin to understand the thought processes that go into building computer systems

Little Theme 1: Representation

- All digital systems represent everything as 0s and 1s
- Everything includes:
 - Numbers integers and floating point
 - Characters the building blocks of strings
 - Instructions the directives to the CPU that make up a program
 - Pointers addresses of data objects in memory
- These encodings are stored in registers, caches, memories, disks, etc.
- They all need addresses
 - A way to find them
 - Find a new place to put a new item
 - Reclaim the place in memory when data no longer needed

Little Theme 2: Translation

- There is a big gap between how we think about programs and data and the 0s and 1s of computers
- Need languages to describe what we mean
- Languages need to be translated one step at a time
 - Word-by-word
 - Phrase structures
 - Grammar
- We know Java as a programming language
 - Have to work our way down to the 0s and 1s of computers
 - Try not to lose anything in translation!
 - We'll encounter Java byte-codes, C language, assembly language, and machine code (for the X86 family of CPU architectures)

Little Theme 3: Control Flow

- How do computers orchestrate the many things they are doing – seemingly in parallel
- What do we have to keep track of when we call a method, and then another, and then another, and so on
- How do we know what to do upon "return"
- User programs and operating systems
 - Multiple user programs
 - Operating system has to orchestrate them all
 - Each gets a share of computing cycles
 - They may need to share system resources (memory, I/O, disks)
 - Yielding and taking control of the processor
 - Voluntary or by force?

Course Outcomes

- Foundation: basics of high-level programming (Java)
- Understanding of some of the abstractions that exist between programs and the hardware they run on, why they exist, and how they build upon each other
- Knowledge of some of the details of underlying implementations
- Become more effective programmers
 - More efficient at finding and eliminating bugs
 - Understand the many factors that influence program performance
 - Facility with some of the many languages that we use to describe programs and data
- Prepare for later classes in CSE

Reality 1: Ints ≠ Integers & Floats ≠ Reals

- Representations are finite
- Example 1: Is x2 ≥ 0?
 - o Floats: Yes!
 - o Ints:
 - **4**0000 * 40000 --> 1600000000
 - **■** 50000 * 50000 --> ??
- Example 2: Is (x + y) + z = x + (y + z)?
 - Unsigned & Signed Ints: Yes!
 - Floats:
 - **■** (1e20 + -1e20) + 3.14 --> 3.14
 - 1e20 + (-1e20 + 3.14) --> ??

Code Security Example

```
/* Kernel memory region holding user-accessible data */
#define KSIZE 1024
char kbuf[KSIZE]; int len = KSIZE;
/* Copy at most maxlen bytes from kernel region to user buffer */
int copy_from_kernel(void *user_dest, int maxlen) {
/* Byte count len is minimum of buffer size and maxlen */
if (KSIZE > maxlen) len = maxlen;
memcpy(user_dest, kbuf, len);
return len;
}
```

- Similar to code found in FreeBSD's implementation of getpeername
- There are legions of smart people trying to find vulnerabilities in programs

Typical Usage

```
/* Kernel memory region holding user-accessible data */
#define KSIZE 1024
char kbuf[KSIZE]; int len = KSIZE;
/* Copy at most maxlen bytes from kernel region to user buffer */
int copy_from_kernel(void *user_dest, int maxlen) {
/* Byte count len is minimum of buffer size and maxlen */
if (KSIZE > maxlen) len = maxlen;
memcpy(user_dest, kbuf, len);
return len;
}
```

```
#define MSIZE 528
void getstuff() {
char mybuf[MSIZE];
copy_from_kernel(mybuf, MSIZE);
printf("%s\n", mybuf);
}
```

Malicious Usage

```
/* Kernel memory region holding user-accessible data */
#define KSIZE 1024
char kbuf[KSIZE]; int len = KSIZE;
/* Copy at most maxlen bytes from kernel region to user buffer */
int copy_from_kernel(void *user_dest, int maxlen) {
/* Byte count len is minimum of buffer size and maxlen */
if (KSIZE > maxlen) len = maxlen;
memcpy(user_dest, kbuf, len);
return len;
}
```

```
#define MSIZE 528
void getstuff() {
char mybuf[MSIZE];
copy_from_kernel(mybuf, -MSIZE);
. . .
}
```

Reality #2: You've Got to Know Assembly

- Chances are, you'll never write a program in assembly code
 - Compilers are much better and more patient than you are
- But: Understanding assembly is the key to the machinelevel execution model
 - Behavior of programs in presence of bugs
 - High-level language model breaks down
 - Tuning program performance
 - Understand optimizations done/not done by the compiler
 - Understanding sources of program inefficiency
 - Implementing system software
 - Operating systems must manage process state
 - Creating / fighting malware
 - x86 assembly is the language of choice

Assembly Code Example

Time Stamp Counter

- Special 64-bit register in Intel-compatible machines
- Incremented every clock cycle
- Read with rdtsc instruction

Application

Measure time (in clock cycles) required by procedure

```
double t;
start_counter();
P();
t = get_counter();
printf("P required %f clock cycles\n", t);
```

Code to Read Counter

- Write small amount of assembly code using GCC's asm facility
- Inserts assembly code into machine code generated by

```
/* Set *hi and *lo to the high and low order bits of the cycle counter.

*/
void access_counter(unsigned *hi, unsigned *lo)
{
    asm("rdtsc; mov! %%edx,%0; mov! %%eax,%1"
    : "=r" (*hi), "=r" (*lo) /* output */
    : /* input */
    : "%edx", "%eax"); /* clobbered */
}
```

Reality #3: Memory Matters

Memory is not unbounded

- It must be allocated and managed
- Many applications are memory-dominated

Memory referencing bugs are especially pernicious

Effects are distant in both time and space

Memory performance is not uniform

- Cache and virtual memory effects can greatly affect program performance
- Adapting program to characteristics of memory system can lead to major speed improvements

Memory Referencing Bug Example

```
double fun(int i)
{

volatile double d[1] = {3.14};

volatile long int a[2];
a[i] = 1073741824; /* Possibly out of bounds */
return d[0];
}
```

```
fun(0) -> 3.14
fun(1) -> 3.14
fun(2) -> 3.1399998664856
fun(3) -> 2.00000061035156
fun(4) -> 3.14, then segmentation fault
```

Memory Referencing Bug Example

```
double fun(int i)
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volatile double d[1] = {3.14};

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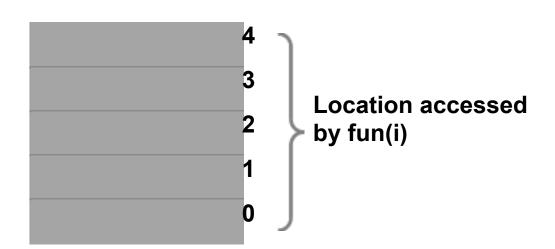
fun(0) -> 3.14

fun(1) -> 3.14

fun(2) -> 3.1399998664856
```

fun(1) -> 3.14 fun(2) -> 3.1399998664856 fun(3) -> 2.00000061035156 fun(4) -> 3.14, then segmentation fault

Explanation:



Memory Referencing Errors

• C (and C++) do not provide any memory protection

- Out of bounds array references
- Invalid pointer values
- Abuses of malloc/free

Can lead to nasty bugs

- Whether or not bug has any effect depends on system and compiler
- Action at a distance
 - Corrupted object logically unrelated to one being accessed
 - Effect of bug may be first observed long after it is generated

How can I deal with this?

- Program in Java (or C#, or ML, or ...)
- Understand what possible interactions may occur
- Use or develop tools to detect referencing errors

Memory System Performance Example

- Hierarchical memory organization
- Performance depends on access patterns
 - Including how program steps through multi-dimensional array

```
void copyij(int src[2048][2048],
int dst[2048][2048])
{
int i,j;
for (i = 0; i < 2048; i++)
for (j = 0; j < 2048; j++)
dst[i][j] = src[i][j];
}</pre>
void copyji(int src[2048][2048],
int dst[2048][2048])
{
int i,j;
for (j = 0; j < 2048; j++)
for (j = 0; i < 2048; i++)
for (j = 0; i < 2048; i++)
for (j = 0; i < 2048; i++)
for (j = 0; j < 2
```

21 times slower (Pentium 4)

Reality #4: Performance isn't counting ops

Exact op count does not predict performance

- Easily see 10:1 performance range depending on how code written
- Must optimize at multiple levels: algorithm, data representations, procedures, and loops

Must understand system to optimize performance

- How programs compiled and executed
- How memory system is organized
- How to measure program performance and identify bottlenecks
- How to improve performance without destroying code modularity and generality

Example Matrix Multiplication

- Standard desktop computer, vendor compiler, using optimization flags
- Both implementations have exactly the same operations count (2n3)



MMM Plot: Analysis

Multiple threads: 4x

Vector instructions: 4x

Memory hierarchy and other optimizations: 20x

- Reason for 20x: blocking or tiling, loop unrolling, array scalarization, instruction scheduling, search to find best choice
- Effect: less register spills, less L1/L2 cache misses, less TLB misses

CSE351's role in new CSE Curriculum

Pre-requisites

142 and 143: Intro Programming I and II

One of 6 core courses

311: Foundations I

312: Foundations II

331: SW Design and Implementation

332: Data Abstractions

351: HW/SW Interface

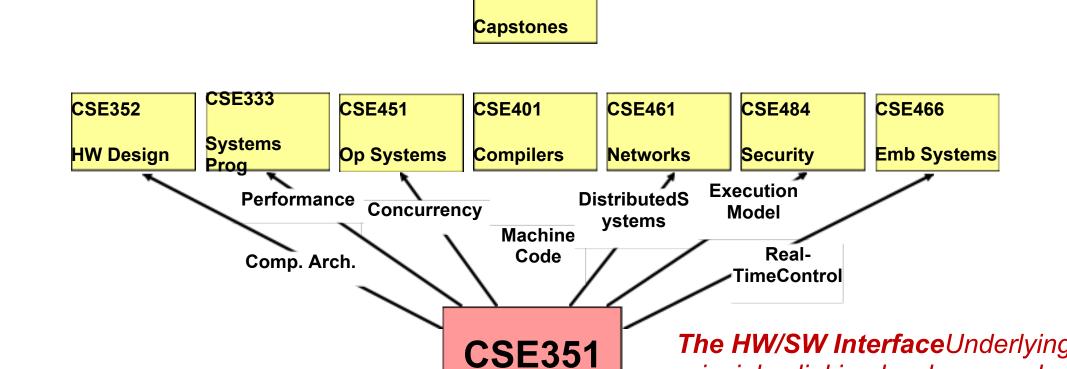
352: HW Design and Implementation

• 351 sets the context for many follow-on courses

principles linking hardware and

software

CSE351's place in new CSE Curriculum



CSE477/481

CS 143 Intro Prog II

Course Perspective

- Most systems courses are Builder-Centric
 - Computer Architecture
 - Design pipelined processor in Verilog
 - Operating Systems
 - Implement large portions of operating system
 - Compilers
 - Write compiler for simple language
 - Networking
 - Implement and simulate network protocols

Course Perspective (Cont.)

This course is Programmer-Centric

- Purpose is to show how software really works
- By understanding the underlying system, one can be more effective as a programmer
 - Better debugging
 - Better basis for evaluating performance
 - How multiple activities work in concert (e.g., OS and user programs)
- Not just a course for dedicated hackers
 - What every CSE major needs to know
- Provide a context in which to place the other CSE courses you'll take

Textbooks

 Computer Systems: A Programmer's Perspective, 2nd Edition

- Randal E. Bryant and David R. O'Hallaron
- o Prentice-Hall, 2010
- http://csapp.cs.cmu.edu
- This book really matters for the course!
 - How to solve labs
 - Practice problems typical of exam problems

A good C book.

- C: A Reference Manual (Harbison and Steele)
- The C Programming Language (Kernighan and Ritchie)



Course Components

Lectures (~30)

Higher-level concepts – I'll assume you've done the reading in the text

Sections (~10)

 Applied concepts, important tools and skills for labs, clarification of lectures, exam review and preparation

Written assignments (4)

Problems from text to solidify understanding

• Labs (4)

Provide in-depth understanding (via practice) of an aspect of systems

Exams (midterm + final)

Test your understanding of concepts and principles

Resources

Course Web Page

- http://www.cse.washington.edu/351
- Copies of lectures, assignments, exams

Course Discussion Board

- Keep in touch outside of class help each other
- Staff will monitor and contribute

Course Mailing List

Low traffic – mostly announcements; you are already subscribed

Staff email

- Things that are not appropriate for discussion board or better offline
- Anonymous Feedback (will be linked from homepage)
 - Any comments about anything related to the coursewhere you would feel better not attaching your name

Policies: Grading

- Exams: weighted 1/3 (midterm), 2/3 (final)
- Written assignments: weighted according to effort
 - We'll try to make these about the same
- Labs assignments: weighted according to effort
 - These will likely increase in weight as the quarter progresses
- Grading:
 - 25% written assignments
 - 35% lab assignments
 - 40% exams

Welcome to CSE351!

- Let's have fun
- Let's learn together
- Let's communicate
- Let's set the bar for a useful and interesting class
- Many thanks to the many instructors who have shared their lecture notes – I will be borrowing liberally through the qtr – they deserve all the credit, the errors are all mine
 - UW: Gaetano Borriello (Inaugural edition of CSE 351, Spring 2010)
 - o CMU: Randy Bryant, David O'Halloran, Gregory Kesden, Markus Püschel
 - Harvard: Matt Welsh
 - UW: Tom Anderson, Luis Ceze, John Zahorjan