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## 1 Conceptual Design, Constraints, Views

1. (25 points)
(a) (10 points) Consider a relation $R(A, B, C, D, E)$ that satisfies the following functional dependencies:

$$
\begin{aligned}
A & \rightarrow B \\
C D & \rightarrow E
\end{aligned}
$$

Decompose the schema in BCNF. Show all your steps. A relation R is in BCNF if and only if: whenever there is a nontrivial functional dependency $A_{1}, A_{2}, \ldots, A_{n} \rightarrow$ $B_{1}, B_{2}, \ldots, B_{n}$ for R, then $\left\{A_{1}, A_{2}, \ldots, A_{n}\right\}$ is a superkey for R.
Answer (Show the steps leading to the BCNF decomposition and show the keys in the decomposed relations):

## BCNF Decomposition Algorithm

## BCNF_Decompose(R)

find X s.t.: $\mathrm{X} \neq \mathrm{X}^{+} \neq$[all attributes]
if (not found) then " $R$ is in BCNF"
let $\mathrm{Y}=\mathrm{X}^{+}-\mathrm{X}$
let $Z=$ [all attributes] $-X^{+}$
decompose R into R1 $(\mathrm{X} \cup \mathrm{Y})$ and R2 $(\mathrm{X} \cup \mathrm{Z})$ continue to decompose recursively R1 and R2

Solution: Straightforward application of BCNF Decomp Algo from lecture 16, slide 44.

Iteration 1: R
$\mathrm{A}+=\mathrm{AB}$ Decompose into $\mathrm{R} 1=\underline{\mathrm{AB}}, \mathrm{R} 2=\mathrm{ACDE}$.
Continue to decompose R2 since R1 is in BCNF form already.
Iteration 2: R2
$\mathrm{CD}+=\mathrm{CDE}$
Decompose R 2 into $\mathrm{R} 3=\underline{\mathrm{CDE}}$ and $\mathrm{R} 4=\underline{\mathrm{CDA}}$

## 2 Transactions

(b) (25 points)

Consider a database consisting of a single relation R :
R:

| A | B |
| :--- | :--- |
| 1 | 10 |
| 2 | 20 |

Two transactions run concurrently on this database, resulting in the following schedule:

| Line | T1 | T2 |
| :---: | :---: | :---: |
| 1 | begin; |  |
| 2 |  | begin; |
| 3 | update $R$ set $B=($ select $\operatorname{sum}(B)$ from $R)$ where $A=1$; |  |
| 4 |  | update $R$ set $B=($ select $\operatorname{sum}(B)$ from $R)$ where $A=2$; |
| 5 | select * from R; |  |
| 6 |  | select * from R; |
| 7 | insert into r values ( 3,300 ); |  |
| 8 |  | insert into r values ( 4,400 ; |
| 9 | select * from R; |  |
| 10 |  | select * from R; |
| 11 | update $R$ set $B=($ select $\operatorname{sum}(B)$ from $R)$ where $A=1$; |  |
| 12 |  | update $R$ set $B=($ select $\operatorname{sum}(B)$ from $R)$ where $A=2$; |
| 13 | select * from R; |  |
| 14 |  | select * from R; |
| 15 | commit; |  |
| 16 |  | commit; |

(a) (5 points) Is this schedule possible in SQL Lite? If not, then indicate the first line where SQL Lite will change the schedule.
(a) No: line 4

Yes ? Or No (and indicate line number) ?
(b) (5 points) Is this schedule possible in SQL Server ? If not, then indicate the first line where SQL Server will change the schedule.
(b) No: line 4

Yes ? Or No (and indicate line number) ?
(c) (10 points) Consider running these two transactions in SQL Server, using isolation level SERIALIZABLE. Indicate the result of each of the six select $*$ statements, as well as the content of the table after both transactions commit.

| Line Number | Result of select * from r; |
| :--- | :--- |
| 5 |  |
| 9 |  |
| 13 |  |
| after T1 commits |  |
| 6 |  |
| 10 |  |
| 14 |  |
| after T2 commits |  |


| Solution: | Line Number | Result of select * from r; |
| :---: | :---: | :---: |
|  | 5 | (1,30), (2,20) |
|  | 9 | (1,30), $(2,20),(3,300)$ |
|  | 13 | $(1,350),(2,20),(3,300)$ |
|  | after T1 commits | (1,350), (2,20), (3,300) |
|  | 6 | (1,350), (2,670), (3,300) |
|  | 10 | (1,350), (2,670), (3,300), (4,400) |
|  | 14 | (1,350), (2,1720), (3,300), (4,400) |
|  | after T2 commits | (1,350), (2,1720), (3,300), (4,400) |

(d) (5 points) Is the schedule for these transactions serializable?
(d) Yes. T1; T2

Yes or No?

