CSE 341 - Programming Languages Midterm - Autumn 2015

Your Name:

```
(for recording grades):
1. (max 6)
2. (max 12)
3. (max 10)
4. (max 5)
5. (max 6)
6. (max 6)
7. (max 6)
8. (max 5)
9. (max 8)
10. (max 8)
```

You can bring a maximum of 2 (paper) pages of notes. No laptops, tablets, or smart phones.

1. (6 points) Suppose the following Racket program has been read in.

```
(define x '(1 2 3))
(define y '(10 11 12))
(define z (append (cdr x) (cdr y)))
```

Draw box-and-arrow diagram of the resulting lists, being careful to show correctly what parts are shared and what parts are separate. (If you run out of room, please draw a fresh copy of the diagram on the back of this page instead.)

2. (12 points) Consider the following Haskell function definitions. (Some of these are from the Prelude; some are written from scratch.)

```
filter p xs = [ x | x <- xs, p x]
repeat x = x : repeat x
memb xs a = or (map (==a) xs)
vowel = memb "aeiou"</pre>
```

Here are some possible types for each of these functions. After each type, write G if the type is correct and the most general, C if it is correct but not the most general, and I if it is incorrect.

```
filter :: [a] -> [a] -> [a]

filter :: (a -> Bool) -> [a] -> [a]

filter :: (Eq a) => (a -> Bool) -> [a] -> [a]

filter :: (Num a) => (a -> Bool) -> [a] -> [a]

filter :: a -> a

repeat :: a -> a

repeat :: [a] -> [[a]]

memb :: [a] -> a -> Bool

memb :: Eq a => [a] -> a -> Bool

memb :: Ord a => [a] -> a -> Bool

vowel :: Char -> Bool

vowel :: String -> Bool
```

- 3. (10 points) Consider a Haskell function remove that takes an item and a list, and returns a new list, dropping elements that are equal to the item. For example, remove 3 [1,2,3,4,3,2,1] evaluates to [1,2,4,2,1], and remove 3 [] evaluates to [].
 - (a) What is the most general type for remove?
 - (b) Write a recursive definition of remove.
 - (c) Write a non-recursive definition of remove using Haskell's filter function. (Hint: there is a definition of filter in Question 2.)

4.	(5 points) What is the value of mystery? (If it's infinite give the first several elements.) mystery = 0 : map $(\x->2*x+1)$ mystery
5.	(6 points) Define a curried plus function curried-plus in Racket. So (curried-plus 1) should return a function that adds 1 to numbers, and ((curried-plus 1) 3) should evaluate to 4.
6.	(6 points) Now define a function <code>extra-spicy-plus</code> in Racket that is also curried but that works with 3 numbers instead of 2. So (((extra-spicy-plus 1) 2) 3) should evaluate to 6.
7.	(6 points) Consider the following Racket expressions. (They are identical except that the first uses let, the second uses let*, and the third uses letrec.) What does each one evaluate to? (let ([y 1]
	<pre>(let ([y 1] [f (lambda (x) (+ x 1))]) (let* ([y 5]</pre>

[f (lambda (x) (if (> x 5) x (f (+ x y))))])

(let ([y 1]

(letrec ([y 5]

(f 3)))

[f (lambda (x) (+ x 1))])

8. (5 points) Consider the following Racket program.

- (a) What is the result of evaluating (test)?
- (b) Suppose Racket used dynamic scoping. What would be the result of evaluating (test)?
- 9. (8 points) The lecture notes for Racket macros included a cosmetic macro for my-if (which just provides a different syntax for if).

Show the code to add my-if to OCTOPUS. For full credit, do this by rewriting the my-if expression as a normal if and then calling eval on the result, in the same way that the Racket my-if macro produces new Racket source code that is then interpreted.

Hint: you should write another case for the eval function. Here are two example cases from the starter file:

```
-- A quoted expression evaluates to that expression.
eval (OctoList [OctoSymbol "quote", x]) env = x
{- An expression starting with (lambda ...) evaluates to a closure,
where a closure consists of a list of variable names (OctoSymbols),
the environment of definition, and the body. -}
eval (OctoList [OctoSymbol "lambda", OctoList vars, body]) env =
    OctoClosure vars env body
```

10. (8 points) Write a slightly silly Racket macro called const that does something like const in Haskell: the result of evaluating (const exp1 exp2) should be the result of evaluating exp1 (just throw exp2 away without evaluating it). Unlike Haskell, though, evaluate exp1 each time the const expression is evaluated. For example, (const (+ 2 3) (/ 1 0)) should evaluate to 5. (There is no divide-by-zero error since we don't evaluate the second expression.)

11. (8 points) Convert the following Haskell action into an equivalent one that doesn't use do.

```
echo = do
    x <- readLn
    y <- readLn
    putStr "the sum is "
    putStrLn (show (x+y))</pre>
```