CSE341: Programming Languages
Lecture 23
Multiple Inheritance, Mixins, Interfaces, Abstract Methods

Dan Grossman
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What next?

Have used classes for OOP's essence: inheritance, overriding, dynamic dispatch

Now, what if we want to have more than just 1 superclass

- **Multiple inheritance**: allow > 1 superclasses
  - Useful but has some problems (see C++)

- Ruby-style *mixins*: 1 superclass; > 1 method providers
  - Often a fine substitute for multiple inheritance and has fewer problems (see also Scala *traits*)

- Java/C#-style *interfaces*: allow > 1 types
  - Mostly irrelevant in a dynamically typed language, but fewer problems
Multiple Inheritance

• If inheritance and overriding are so useful, why limit ourselves to one superclass?
  – Because the semantics is often awkward (this topic)
  – Because it makes static type-checking harder (not discussed)
  – Because it makes efficient implementation harder (not discussed)

• Is it useful? Sure!
  – Example: Make a ColorPt3D by inheriting from Pt3D and ColorPt (or maybe just from Color)
  – Example: Make a StudentAthlete by inheriting from Student and Athlete
  – With single inheritance, end up copying code or using non-OOP-style helper methods
Trees, dags, and diamonds

- Note: The phrases subclass, superclass can be ambiguous
  - There are immediate subclasses, superclasses
  - And there are transitive subclasses, superclasses

- Single inheritance: the class hierarchy is a tree
  - Nodes are classes
  - Parent is immediate superclass
  - Any number of children allowed

- Multiple inheritance: the class hierarchy is no longer a tree
  - Cycles still disallowed (a directed-acyclic graph)
  - If multiple paths show that X is a (transitive) superclass of Y, then we have diamonds
What could go wrong?

• If \( V \) and \( Z \) both define a method \( m \), what does \( Y \) inherit? What does `super` mean?
  – *Directed resends* useful (e.g., \( Z::\text{super} \))

• What if \( X \) defines a method \( m \) that \( Z \) but not \( V \) overrides?
  – Can handle like previous case, but sometimes undesirable (e.g., `ColorPt3D` wants `Pt3D`'s overrides to “win”)

• If \( X \) defines fields, should \( Y \) have one copy of them (\( f \)) or two (\( V::f \) and \( Z::f \))?
  – Turns out each behavior can be desirable (next slides)
  – So C++ has (at least) two forms of inheritance
3DColorPoints

If Ruby had multiple inheritance, we would want ColorPt3D to inherit methods that share one @x and one @y

class Pt
  attr_accessor :x, :y
  ...
end
class ColorPt < Pt
  attr_accessor :color
  ...
end
class Pt3D < Pt
  attr_accessor :z
  ... # override some methods
end
class ColorPt3D < Pt3D, ColorPt # not Ruby!
end
This code has Person define a pocket for subclasses to use, but an ArtistCowboy wants two pockets, one for each draw method.

class Person
    attr_accessor :pocket
    ...
end
class Artist < Person # pocket for brush objects
    def draw # access pocket
    ...
end
class Cowboy < Person # pocket for gun objects
    def draw # access pocket
    ...
end
class ArtistCowboy < Artist, Cowboy # not Ruby!
end
Mixins

• A *mixin* is (just) a collection of methods
  – Less than a class: no instances of it

• Languages with mixins (e.g., Ruby modules) typically let a class have one superclass but *include* number of mixins

• Semantics: *Including a mixin makes its methods part of the class*
  – Extending or overriding in the order mixins are included in the class definition
  – More powerful than helper methods because mixin methods can access methods (and instance variables) on *self* not defined in the mixin
Example

```ruby
module Doubler
  def double
    self + self # assume included in classes w/ +
  end
end
class String
  include Doubler
end
class AnotherPt
  attr_accessor :x, :y
  include Doubler
  def + other
    ans = AnotherPt.new
    ans.x = self.x + other.x
    ans.y = self.y + other.y
    ans
  end
end
```
Lookup rules

Mixins change our lookup rules slightly:

- When looking for receiver obj's method m, look in obj's class, then mixins that class includes (later includes shadow), then obj's superclass, then the superclass' mixins, etc.

- As for instance variables, the mixin methods are included in the same object
  - So usually bad style for mixin methods to use instance variables since a name clash would be like our CowboyArtist pocket problem (but sometimes unavoidable?)
The two big ones

The two most popular/useful mixins in Ruby:

• Comparable: Defines `<`, `>`, `==`, `!=`, `>=`, `<=` in terms of `<=>`

• Enumerable: Defines many iterators (e.g., `map`, `find`) in terms of `each`

Great examples of using mixins:
  – Classes including them get a bunch of methods for just a little work
  – Classes do not “spend” their “one superclass” for this
  – Do not need the complexity of multiple inheritance

• See the code for some examples
Replacement for multiple inheritance?

- A mixin works pretty well for ColorPt3D:
  - Color a reasonable mixin except for using an instance variable

```ruby
module Color
  attr_accessor :color
end
```

- A mixin works awkwardly-at-best for ArtistCowboy:
  - Natural for Artist and Cowboy to be Person subclasses
  - Could move methods of one to a mixin, but it is odd style and still does not get you two pockets

```ruby
module ArtistM ...
class Artist < Person
  include ArtistM
class ArtistCowboy < Cowboy
  include ArtistM
```
**Statically-Typed OOP**

- Now contrast multiple inheritance and mixins with Java/C#-style interfaces

- Important distinction, but interfaces are about static typing, which Ruby does not have

- So will use Java code after quick introduction to static typing for class-based OOP…
  - Sound typing for OOP prevents “method missing” errors
Classes as Types

- In Java/C#/etc. each class is also a type

- Methods have types for arguments and result

```java
class A {
    Object m1(Example e, String s) {...}
    Integer m2(A foo, Boolean b, Integer i) {...}
}
```

- If $C$ is a (transitive) subclass of $D$, then $C$ is a subtype of $D$
  - Type-checking allows subtype anywhere supertype allowed
  - So can pass instance of $C$ to a method expecting instance of $D$
Interfaces are Types

interface Example {
    void m1(int x, int y);
    Object m2(Example x, String y);
}

• An interface is not a class; it is only a type
  – Does not contain method definitions, only their signatures (types)
    • Unlike mixins
  – Cannot use `new` on an interface
    • Like mixins
Implementing Interfaces

• A class can explicitly implement any number of interfaces
  – For class to type-check, it must implement every method in the interface with the right type
    • More on allowing subtypes later!
  – Multiple interfaces no problem; just implement everything

• If class type-checks, it is a subtype of the interface

```java
class A implements Example {
    public void m1(int x, int y) {...}
    public Object m2(Example e, String s) {...}
}
class B implements Example {
    public void m1(int pizza, int beer) {...}
    public Object m2(Example e, String s) {...}
}
```
Multiple interfaces

• Interfaces provide no methods or fields
  – So no questions of method/field duplication when implementing multiple interfaces, unlike multiple inheritance

• What interfaces are for:
  – “Caller can give any instance of any class implementing I”
    • So callee can call methods in I regardless of class
  – So much more flexible type system

• Interfaces have little use in a dynamically typed language
  – Dynamic typing already much more flexible, with trade-offs we studied
Connections

Let’s now answer these questions:

• What does a statically typed OOP language need to support “required overriding”?

• How is this similar to higher-order functions?

• Why does a language with multiple inheritance (e.g., C++) not need Java/C#-style interfaces?

[Explaining Java’s abstract methods / C++’s pure virtual methods]
Required overriding

Often a class expects all subclasses to override some method(s)
  – The purpose of the superclass is to abstract common functionality, but some non-common parts have no default

A Ruby approach:
  – Do not define must-override methods in superclass
  – Subclasses can add it
  – Creating instance of superclass can cause method-missing errors

```ruby
# do not use A.new
# all subclasses should define m2
class A
  def m1 v
    ... self.m2 e ...
  end
end
```
Static typing

- In Java/C#/C++, prior approach fails type-checking
  - No method \texttt{m2} defined in superclass
  - One solution: provide error-causing implementation

```ruby
class A
  def m1 v
    ... self.m2 e ... 
  end
  def m2 v
    raise "must be overridden"
  end
end
```

- Better: Use static checking to prevent this error…
Abstract methods

• Java/C#/C++ let superclass give signature (type) of method subclasses should provide
  – Called abstract methods or pure virtual methods
  – Cannot creates instances of classes with such methods
    • Catches error at compile-time
    • Indicates intent to code-reader
    • Does not make language more powerful

```java
abstract class A {
    T1 m1(T2 x) { ... m2(e); ... }
    abstract T3 m2(T4 x);
}
class B extends A {
    T3 m2(T4 x) { ... }
}
```
Passing code to other code

- Abstract methods and dynamic dispatch: An OOP way to have subclass “pass code” to other code in superclass

```java
abstract class A {
    T1 m1(T2 x) { ... m2(e); ... }
    abstract T3 m2(T4 x);
}
class B extends A {
    T3 m2(T4 x) { ... }
}
```

- Higher-order functions: An FP way to have caller “pass code” to callee

```scala
fun f (g,x) = ... g e ...
fun h x = ... f((fn y => ...) ,...)
```
No interfaces in C++

• If you have multiple inheritance and abstract methods, you do not also need interfaces

• Replace each interface with a class with all abstract methods

• Replace each “implements interface” with another superclass

So: Expect to see interfaces only in statically typed OOP without multiple inheritance
  – Not Ruby
  – Not C++