

# CSE 341: Programming Languages

Autumn 2005

Lecture 7 — Motivation; BNF; First-Class Functions

# Today

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- (course motivation)
- Summarize what we've learned with a concise and well-known notation for recursively-defined language constructs
- Begin first-class functions

# Why these 3 Languages?

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*Functional programming* (ML, Scheme) encourages recursion, discourages mutation, provides elegant, lightweight support for first-class code. Support for extensibility complements OO.

- ML has a polymorphic type system (vindication imminent!) complementary to OO-style subtyping, a rich module system for abstract types, and rich pattern-matching.
- Scheme has dynamic typing, “good” macros, fascinating control operators, and a minimalist design.
- Smalltalk has classes but not types, an unconventional environment, and a complete commitment to OO.

# Runners-Up

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Runners-up: Miranda (laziness, simplicity and purity), Haskell (laziness, nonproprietary implementation, but difficult type system), Prolog (unification and backtracking), CPR(R) (everything that Prolog has plus constraints).

There are thousands of other languages as well . . .

## Why not some popular ones?

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- Java: you've already studied it in 142/143. We'll look at some additional features at the end of course and compare it with the other languages we've been studying (e.g., interfaces, anonymous inner classes, container types)
- C: lots of "implementation-dependent" behavior (a bad property), and we have CSE303
- C++: an enormous language, and unsafe like C
- Perl: advantages (strings, files, ...) not foci of this course. Python or Ruby would be closer.

## Are these useful?

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We focus on *interesting language concepts* in ML/Scheme/Smalltalk.

“Real” programming needs file I/O, strings, floating-point, graphics libraries, project managers, unit testers, threads, foreign-function interfaces, . . .

- These languages have all that and more!
- Just not course focus

# Summary and Some Notation

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Learned the syntax, typing rules, and semantics for (a big) part of ML

Can summarize *abstract syntax* with *(E)BNF*. Informally:

$$\begin{aligned} t ::= & \text{int} \mid \text{bool} \mid \text{unit} \mid \text{dtname} \\ & \mid t_1 \rightarrow t_2 \mid t_1 * t_2 \mid \{x_1=t_1, \dots, x_n=t_n\} \\ e ::= & 34 \mid x \mid (e_1, e_2) \mid \text{if } e_1 \text{ then } e_2 \text{ else } e_3 \\ & \mid \text{let } b_1 \dots b_n \text{ in } e \text{ end} \mid e_1 e_2 \\ & \mid \text{case } e \text{ of } p_1 \Rightarrow e_1 \mid \dots \mid p_n \Rightarrow e_n \\ & \mid e_1 + e_2 \mid \{x_1=e_1, \dots, x_n=e_n\} \mid C e \\ b ::= & \text{val } p = e \mid \text{fun } f p = e \\ & \mid \text{datatype dtname} = C_1 \text{ of } t_1 \mid \dots \mid C_n \text{ of } t_n \\ p ::= & 34 \mid x \mid \_ \mid C p \mid (p_1, p_2) \mid \{x_1=p_1, \dots, x_n=p_n\} \end{aligned}$$

Things left out of this *grammar*:  $n$ -tuples, field-accessors, floating-point, boolean constants, andalso/orelse, lists, ...

# First-Class Functions

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- Functions are values. (Variables in the environment are bound to them.)
- We can pass functions to other functions.
  - *Factor* common parts and *abstract* different parts.
- We can return functions as values from other functions.



# Type Inference and Polymorphism

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ML can infer function types based on function bodies. Possibilities:

- The argument/result must be one specific type.
- The argument/result can be *any* type, but may have to be the *same type* as other parts of argument/result.
- Some hand-waving about “equality types”

We will study this *parametric polymorphism* more later.

Without it, ML would be a pain (e.g., a different list library for every list-element type).

Curious fact: If  $f: \text{int} \rightarrow \text{int}$ , there are lots of values  $f$  could return. If  $f: 'a \rightarrow 'a$ , whenever  $f$  returns, it returns its argument!

# Anonymous Functions

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As usual, we can write functions anywhere we write expressions.

- We already could:

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(let fun f x = e in f end)
```

- Here is a more concise way (better style when possible):

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(fn x => e)
```

- Cannot do this for recursive functions (why?)

# Returning Functions

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Syntax note:  $\rightarrow$  “associates to the right”

- $t_1 \rightarrow t_2 \rightarrow t_3$  means  $t_1 \rightarrow (t_2 \rightarrow t_3)$

Again, there is nothing new here.

The key question: What about *free variables* in a function value?  
What *environment* do we use to *evaluate* them?

Are such free variables useful?

You must understand the answers to move beyond being a novice programmer.