Due: **Wednesday, April 27** at the BEGINNING of lecture. Your work should be readable as well as correct. You should refer to the written homework guidelines on the course website for a reminder about what is acceptable pseudocode.

**Problem 1: B-Tree Insertion**
Omit. This was a duplicate of problem 4 from Homework 3.

**Problem 2: B Tree Predecessor**
In this problem, assume every B tree node has a reference to its parent (except for the root node) as well as its children (except for leaves), so that it is easy to move up or down" the tree.

Suppose you have a direct reference to the leaf holding a data item with a key k.

a) Describe in English an algorithm for finding the predecessor of k in the B tree (or determining that k is the minimum element and therefore has no predecessor).

b) Give the worst-case number of total nodes accessed by your algorithm in terms of the tree height h. Describe briefly a worst-case situation.

c) Give the best-case number of total nodes accessed by your algorithm. Explain briefly why most keys would exhibit the best-case behavior.

**Problem 3: Textbook exercise**

**Problem 4: Deletion in Hashing**
In this problem you will think about how lazy deletion is handled in open addressing hash tables.

(a) Suppose a hash table is accessed by open addressing and contains a cell X marked as “deleted”. Suppose that the next successful find hits and moves past cell X and finds the key in cell Y. Suppose we move the found key to cell X, mark cell X as “active” and mark cell Y as “open”. Suppose this policy is used for every find. Would you expect this to work better or worse compared to not modifying the table? Explain your answer.

(b) Suppose that instead of marking cell Y as “open” in the previous question, you mark it as “deleted” (it contains no value, but we treat it as a collision). Suppose this policy is used for every find. Would you expect this to work better or worse compared to not modifying the table? Explain your answer.