CSE332: Data Abstractions

Section 2

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Section Agenda

- Bugs & Testing
- Induction Review
- Recurrence Relations
- Asymptotic Analysis
- Homework Tips & Questions

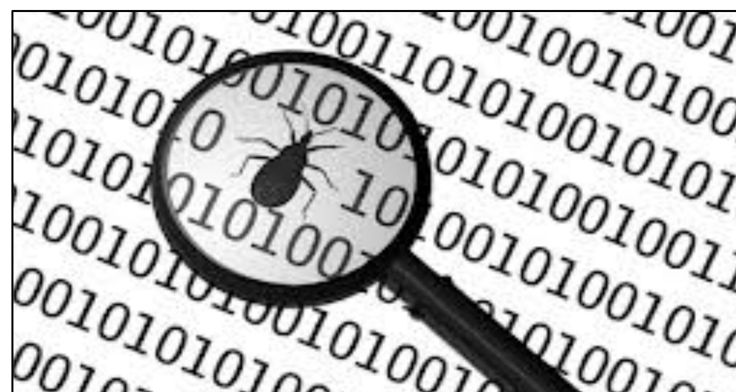
Bugs & Testing

Bugs & Testing

- **Software Bugs**

- Error in a computer program

- Causes program to behave in unexpected ways



Bugs & Testing

- **Why Testing?**

Bugs can be costly

- Cost points in homework
- Can cost $$$ and even life (Therac-25)

Interesting Bug References

- List of bugs http://en.wikipedia.org/wiki/List_of_software_bugs
- History's worst http://www.wired.com/software/coolapps/news/2005/11/69355?currentPage=all
- Bugs of the month http://www.gimpel.com/html/bugs.htm

Bugs & Testing

- **Reverse.java**
  does not test your stack!!

  - Stack can still have lots of bugs
    when working perfectly with Reverse.java

  - Some extreme case in past quarter:
    it only worked with Reverse.java
    (Not a good stack!)

Bugs & Testing

- **Tips for Testing**

  - Make sure program meets the spec

  - Test if each method works independently

  - Test if methods work together

  - Test for edge cases

Bugs & Testing

- **Make sure program meets the spec**
  - What is wrong with this implementation?

```
public class ListStack implements DStack {
    private LinkedList<Double> myStack;

    public ListStack() {
        myStack = new LinkedList<Double>();
    }
    public void push(double d) {
        myStack.add(d);
    }
    …
}
```

Bugs & Testing

- **Test for edge cases**

  - Empty stack

  - Push after resizing

  - Anything else?

Bugs & Testing

- **Testing tools:** JUnit Testing

  - Not required for Project 1

  - Required for Project 2

  - Covered in section later

Induction Review

Induction Review

- **Proof by Induction**

  - Prove that the **first** statement in the infinite sequence of statements is true (Base case)

  - Prove that if **any one** statement in the infinite sequence of statements is true, then so is the **next** one. (Inductive case)

Induction Review

- **Proof by Induction**

  To prove statement P(n),

- Base Case:

  Prove that P(1) is true

- Inductive Case:

  Assuming P(k) is true,
  prove that P(k+1) is true

Recurrence Relations

Recurrence Relations

- **Recursively defines a Sequence**

  - Example: $T(n) = T(n-1) + 3$, $T(1) = 5$

    ^ Has $T(x)$ in definition

- **Solving Recurrence Relation**

  - Eliminate recursive part in definition

    = Find "Closed Form"

  - Example: $T(n) = 3n + 2$

Recurrence Relations

- **Expansion Method example**

  - Solve $T(n) = T(n-1) + 2n - 1$, $T(1) = 1$

$$T(n) = T(n-1) + 2n - 1$$

$$T(n-1) = T([n-1]-1) + 2[n-1] - 1$$

$$= T(n-2) + 2(n-1) - 1$$

$$T(n-2) = T([n-2]-1) + 2[n-2] - 1$$

$$= T(n-3) + 2(n-2) - 1$$

Recurrence Relations

- **Expansion Method example**

$$T(n) = T(n-1) + 2n - 1$$

$$T(n-1) = T(n-2) + 2(n-1) - 1$$

$$T(n-2) = T(n-3) + 2(n-2) - 1$$

$$T(n) = [T(n-2) + 2(n-1) - 1] + 2n - 1$$

$$= T(n-2) + 2(n-1) + 2n - 2$$

$$T(n) = [T(n-3) + 2(n-2) - 1] + 2(n-1) + 2n - 2$$

$$= T(n-3) + 2(n-2) + 2(n-1) + 2n - 3$$

Recurrence Relations

- **Expansion Method example**

$$T(n) = T(n-1) + 2n - 1$$

$$T(n) = T(n-2) + 2(n-1) + 2n - 2$$

$$T(n) = T(n-3) + 2(n-2) + 2(n-1) + 2n - 3$$

…

$$T(n) = T(n-k) + [2(n-(k-1)) + \ldots + 2(n-1) + 2n] - k$$

$$= T(n-k) + [2(n-k+1) + \ldots + 2(n-1) + 2n] - k$$

Recurrence Relations

- **Expansion Method example**

$T(n) = T(n-k) + [2(n-k+1) + \ldots + 2(n-1) + 2n] - k$

When expanded all the way down, $T(n-k) = T(1)$

$n-k = 1$, $k = n-1$

$T(n) = T(n-[n-1]) + [2(n-[n-1]+1) + \ldots + 2(n-1)$
$+ 2n] - [n-1]$

$= T(1) + [2(2) + \ldots + 2(n-1) + 2n] - n + 1$

Recurrence Relations

- **Expansion Method example**

$T(n) = T(1) + [2(2) + \ldots + 2(n\text{-}1) + 2n] – n + 1$

$= T(1) + 2[2 + \ldots + (n\text{-}1) + n] – n + 1$

$= T(1) + 2[(n+1)(n/2) -1] – n + 1$

$= T(1) + (n+1)(n) - 2 – n + 1$

$= T(1) + (n^2+n) – n - 1$

$= T(1) + n^2 – 1$

$= 1 + n^2 – 1$

$= n^2$

Recurrence Relations

- **Expansion Method example** **Check it!**

$T(n) = T(n-1) + 2n - 1, \quad T(1) = 1$

$T(n) = n^2$

$T(1) = 1$	same as 1^2
$T(2) = T(1) + 2(2) - 1 = 4$	same as 2^2
$T(3) = T(2) + 2(3) - 1 = 9$	same as 3^2
$T(4) = T(3) + 2(4) - 1 = 16$	same as 4^2

Recurrence Relations

- **For Homework**

  Remember to show steps!!

  - Correct answer with no steps gets no credit

    a) Show at least 2 expansions of T(n)

    b) At least 2 representations of T(n), using a)

    c) Writing T(n) in terms of k, using b)

    d) How you solved for k

    e) Plug in k and get the closed form

Asymptotic Analysis

Asymptotic Analysis

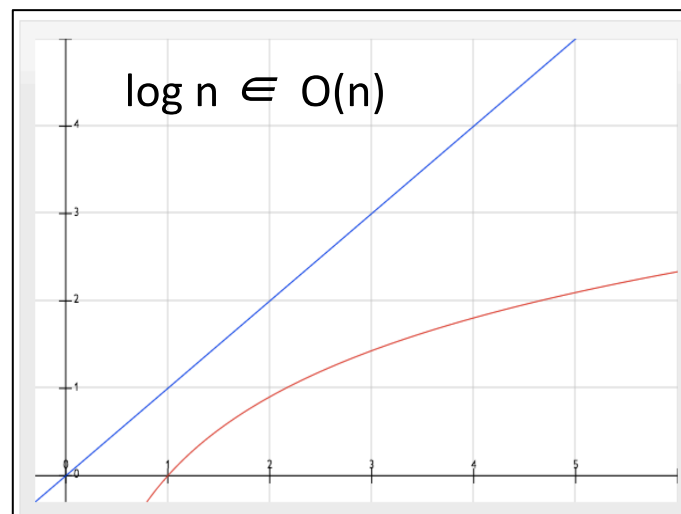
- **Describe Limiting behavior of F(n)**

  - Characterize growth rate of F(n)

  - Use O(g(n)), *Ω*(g(n)), *Θ*(g(n)) for **set** of functions with asymptotic behavior ≤, ≥, ≤ & ≥ to g(n)

- **Upper Bound: O(n)**

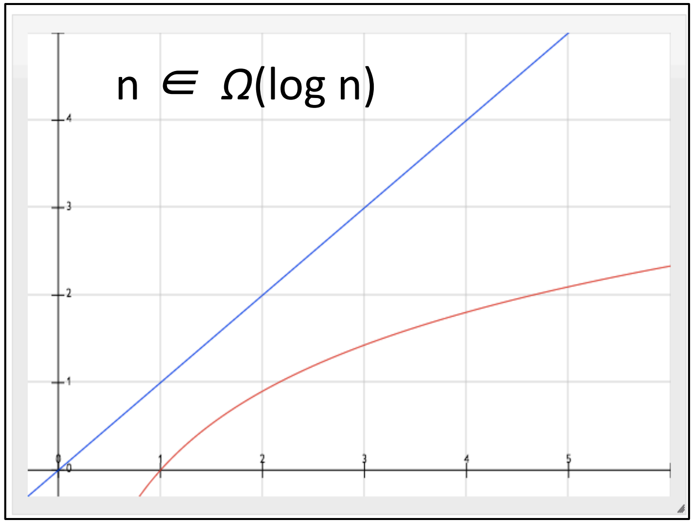
  f(n) $\in$ O(g(n)) if and only if there exist positive constants c and n_0 such that $f(n) \leq c*g(n)$ for all $n_0 \leq n$



Asymptotic Analysis

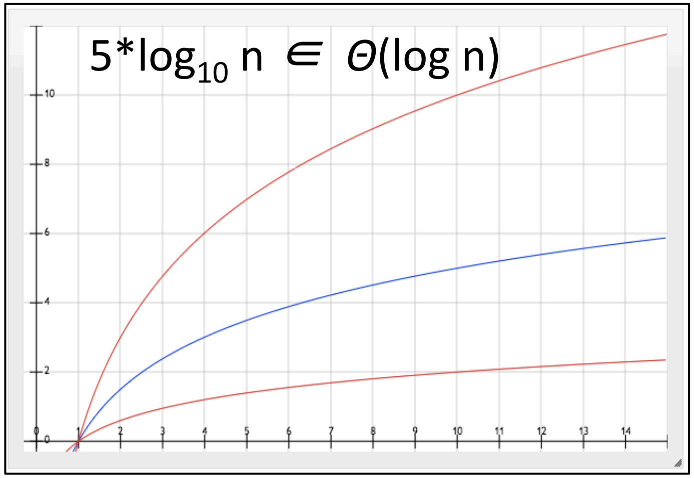
- **Lower Bound: Ω(n)**

  f(n) ∈ *Ω*(g(n)) if and only if
  if there exist positive constants
  c and n_0 such that
  c*g(n) ≤ f(n) for all n_0 ≤ n



- **Tight Bound: Θ(n)**

  f(n) ∈ *Θ*(g(n)) if and only if
  f(n) ∈ *Ω*(g(n)) and
  f(n) ∈ O(g(n))



Asymptotic Analysis

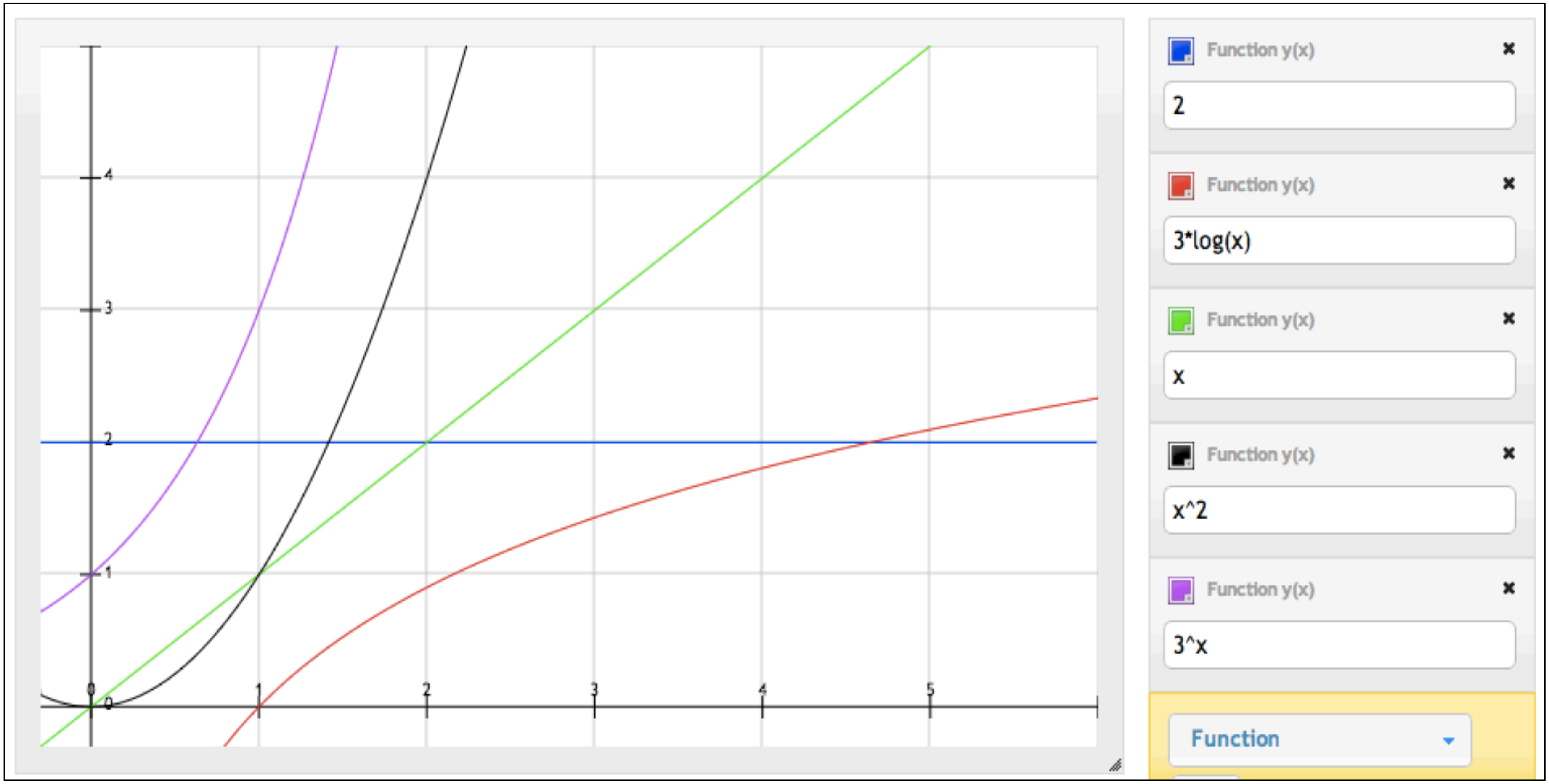
- **Ordering Growth rates (k = constant)**
  - Ignore Low-Order terms & Coefficients

$O(k)$	constant
$O(\log n)$	logarithmic
$O(n)$	linear
$O(n^k)$	polynomial
$O(k^n)$	exponential ($k > 1$)

Increasing Growth rate

Asymptotic Analysis

- **Ordering Growth rates**



Asymptotic Analysis

- **Ordering Growth rates (k, b = constant)**
  - $\log^k n \in O(n^b)$ if $1 < k$ & $0 < b$
  - $n^k \in O(b^n)$ if $0 < k$ & $1 < b$

- **Ordering Example**

$2n^{100} + 10n$	n^{100}	4
$2^{n/100} + 2^{n/270}$	$2^{n/100}$	5
$1000n + \log^8 n$	n	3
$23785n^{1/2}$	$n^{1/2}$	2
$1000 \log^{10} n + 1^{n/300}$	$\log^{10} n$	1

Asymptotic Analysis

- **Proof Example: f(n) $\in$ O(g(n))**

  - Prove or disprove nlog n $\in$ O(3n)

nlog n $\in$ O(3n), then by definition of Big-O

nlog n $\leq$ c*(3n), for 0 < c && 0 < $n_0 \leq n$

(1/3)log n $\leq$ c

but as $n \to \infty$, $\log n \to \infty$

Finite constant c always greater than log n cannot exist, no matter what n_0 we choose

nlog n $\notin$ O(3n)

Homework Tips

Homework Tips

- **Problem #1**

  - Use formula in the book
  (You don't have to derive it by yourself)

- **Problem #2**

  - Use following rules:

  1. $\left\lfloor \frac{\lfloor \frac{x}{m} \rfloor}{n} \right\rfloor = \lfloor \frac{x}{mn} \rfloor$ which means $\left\lfloor \frac{\lfloor \frac{n}{2} \rfloor}{2} \right\rfloor = \lfloor \frac{n}{2^2} \rfloor$

  2. $\lfloor x \rfloor = m$ if and only if $m \leq x < m + 1$

Homework Tips

- **Problem #3**

  - $f(n) \times 10^{-6}$ sec $\leq$ t sec, solve for n

- **Problem #4 <= Not in this HW**

  - Remember that when you are proving P(k+1), you are assuming P(k) no matter how silly it is!
  - Find flaw in inductive reasoning

Homework Tips

- **Problem #5**

  - Use definitions and show you can/cannot find the constant c

- **Problem #6**

  - Analyze runtime of each loop & merge when appropriate
  - Practice finding exact runtime when you can
  - Think about maximum iteration of each loop