



## CSE332: Data Abstractions

### Lecture 12: Introduction to Sorting

Dan Grossman  
Spring 2010

### *Introduction to sorting*

- Stacks, queues, priority queues, and dictionaries all focused on providing one element at a time
- But often we know we want “all the data items” in some order
  - A huge reason to use computers: an 8-year-old child can sort, but a computer can sort faster
  - Different algorithms have different asymptotic and constant-factor trade-offs
    - Knowing one way to sort just isn't enough

Spring 2010

CSE332: Data Abstractions

2

### *More reasons to sort*

General technique in computing:

*Preprocess data to make subsequent operations faster*

Example: Sort the data so that you can

- Find the  $k^{\text{th}}$  largest in constant time for any  $k$
- Perform binary search to find an element in logarithmic time

Whether the performance of the preprocessing matters depends on

- How often the data will change
- How much data there is

Spring 2010

CSE332: Data Abstractions

3

### *The main problem, stated carefully*

For now we will assume we have  $n$  comparable elements in an array and we want to rearrange them to be in increasing order

Input:

- An array  $\mathbf{A}$  of data records
- A key value in each data record
- A comparison function (consistent and total)

Effect:

- Reorganize the elements of  $\mathbf{A}$  such that for any  $i$  and  $j$ , if  $i < j$  then  $\mathbf{A}[i] \leq \mathbf{A}[j]$
- ( $\mathbf{A}$  must have all the same data it started with)

An algorithm doing this is a [comparison sort](#)

Spring 2010

CSE332: Data Abstractions

4

## Variations on the basic problem

1. Maybe elements are in a linked list (could convert to array and back in linear time, but some algorithms needn't do so)
2. Maybe ties need to be resolved by "original array position"
  - Sorts that do this naturally are called **stable sorts**
  - Others could tag each item with its original position and adjust comparisons accordingly (non-trivial constant factors)
3. Maybe we must not use more than  $O(1)$  "auxiliary space"
  - Sorts meeting this requirement are called **in-place sorts**
4. Maybe we can do more with elements than just compare
  - Sometimes leads to faster algorithms
5. Maybe we have too much data to fit in memory
  - Use an "**external sorting**" algorithm

## The Big Picture

Surprising amount of juicy computer science: 2-3 lectures...

